

Compacting the Uncompactable!

MESH

Automatically Compacting
Memory for C/C++ Applications

Emery Berger
UMass Amherst

with Bobby Powers, David Tench, & Andrew McGregor
University of Massachusetts Amherst

<http://libmesh.org>

[PLDI 2019]



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN

Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN

Music by MITCH LEIGH



Save
all of Spain!



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



Save all of Spain!



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN

Music by MITCH LEIGH



“Reallocate”

Save Allocate
all of Spain!



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



(malloc)



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



DARLON MUSIC COMPANY

PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH

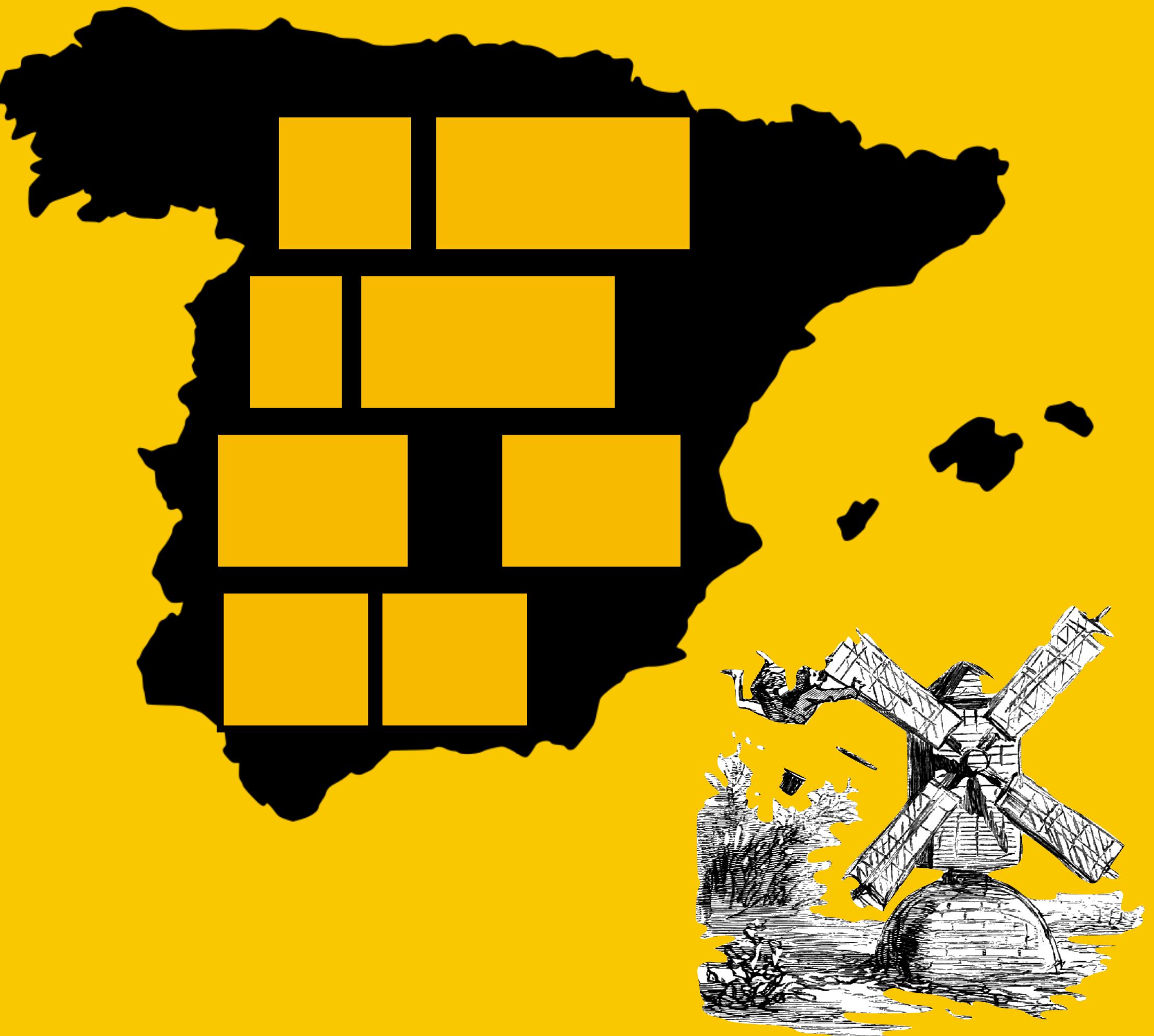


PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH

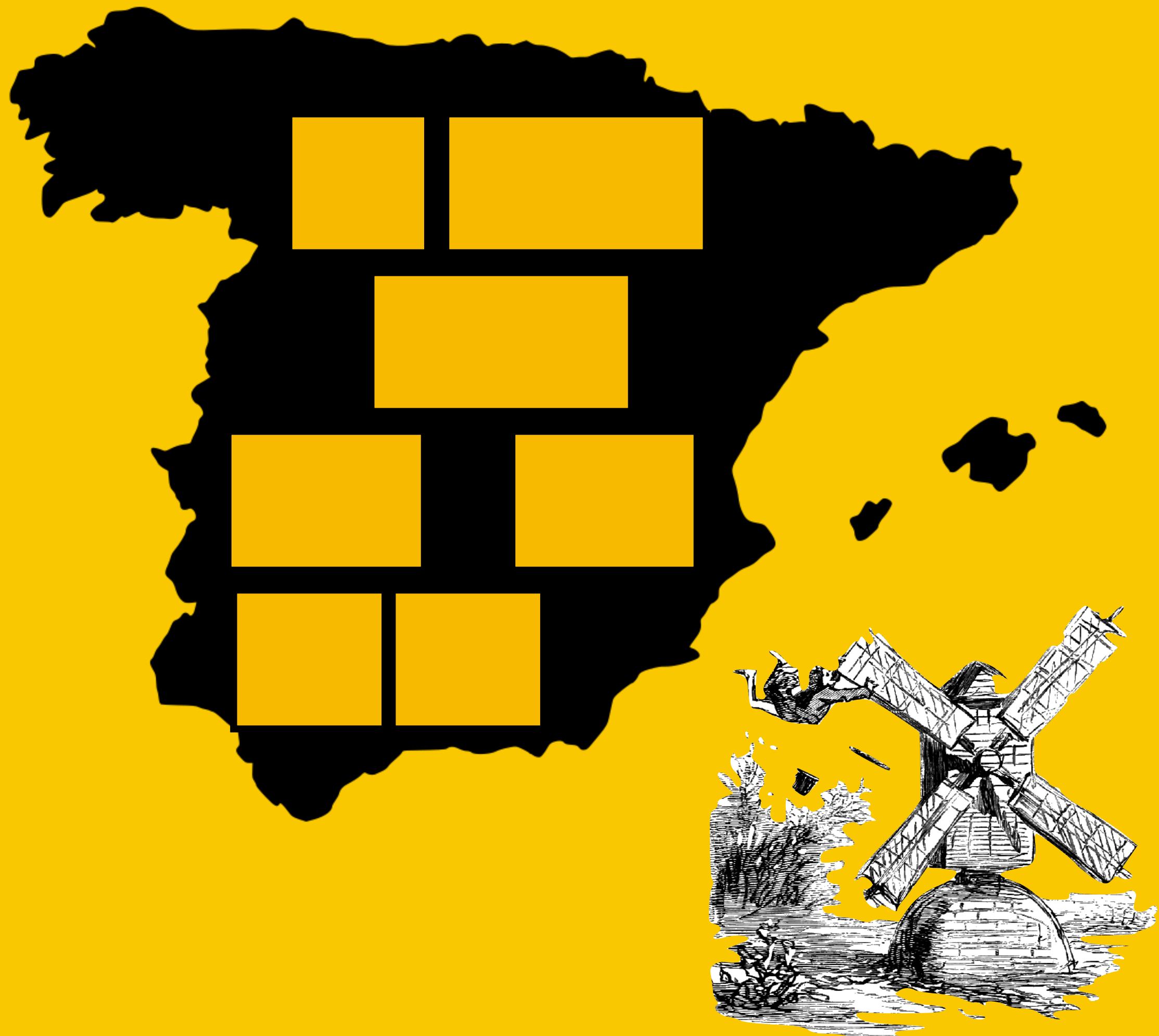


PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH

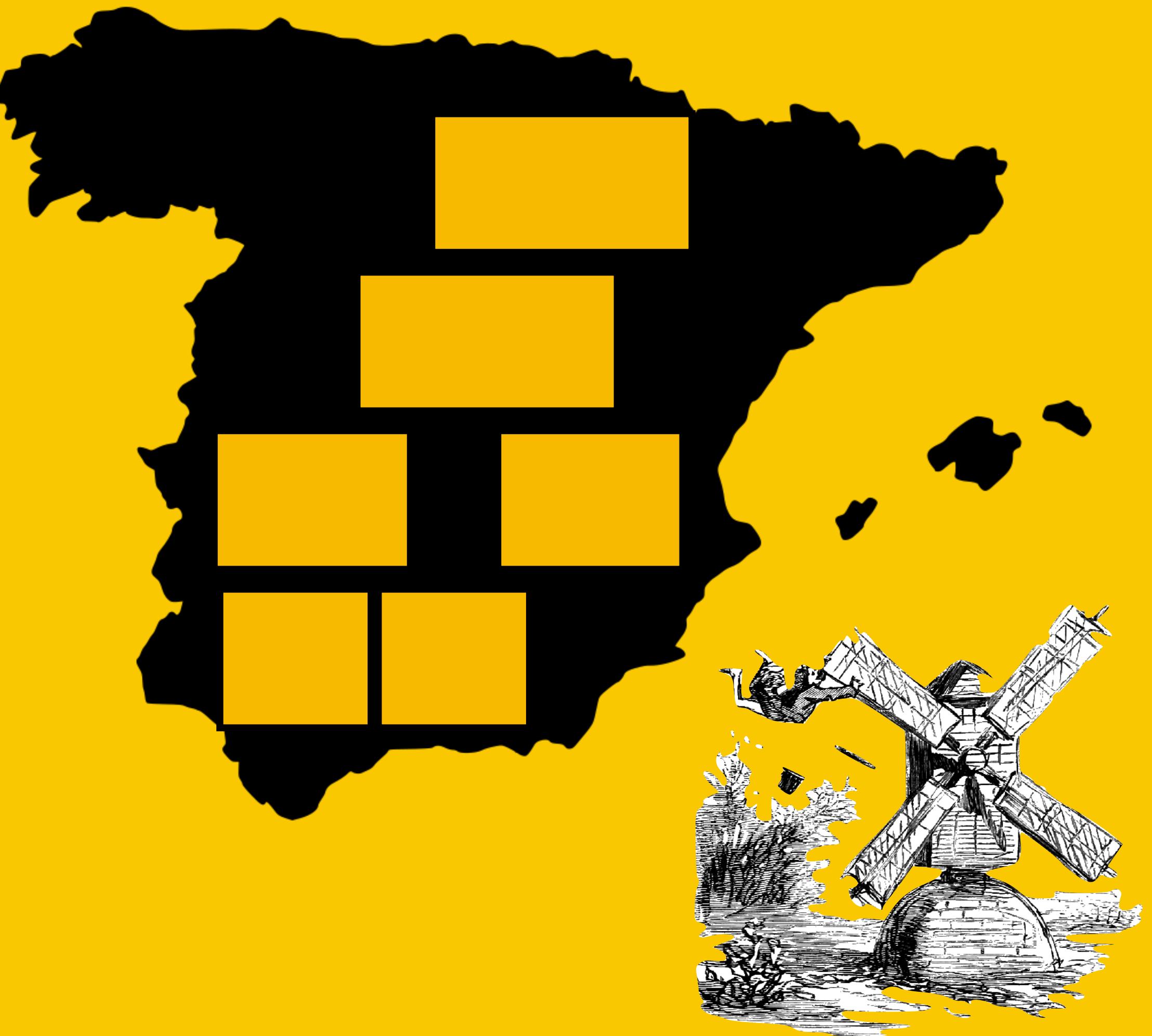


PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



MALLOC



!

PIANO • VOCAL • GUITAR



The Impossible Dream (The Quest)

Lyrics by JOE DARIEN
Music by MITCH LEIGH



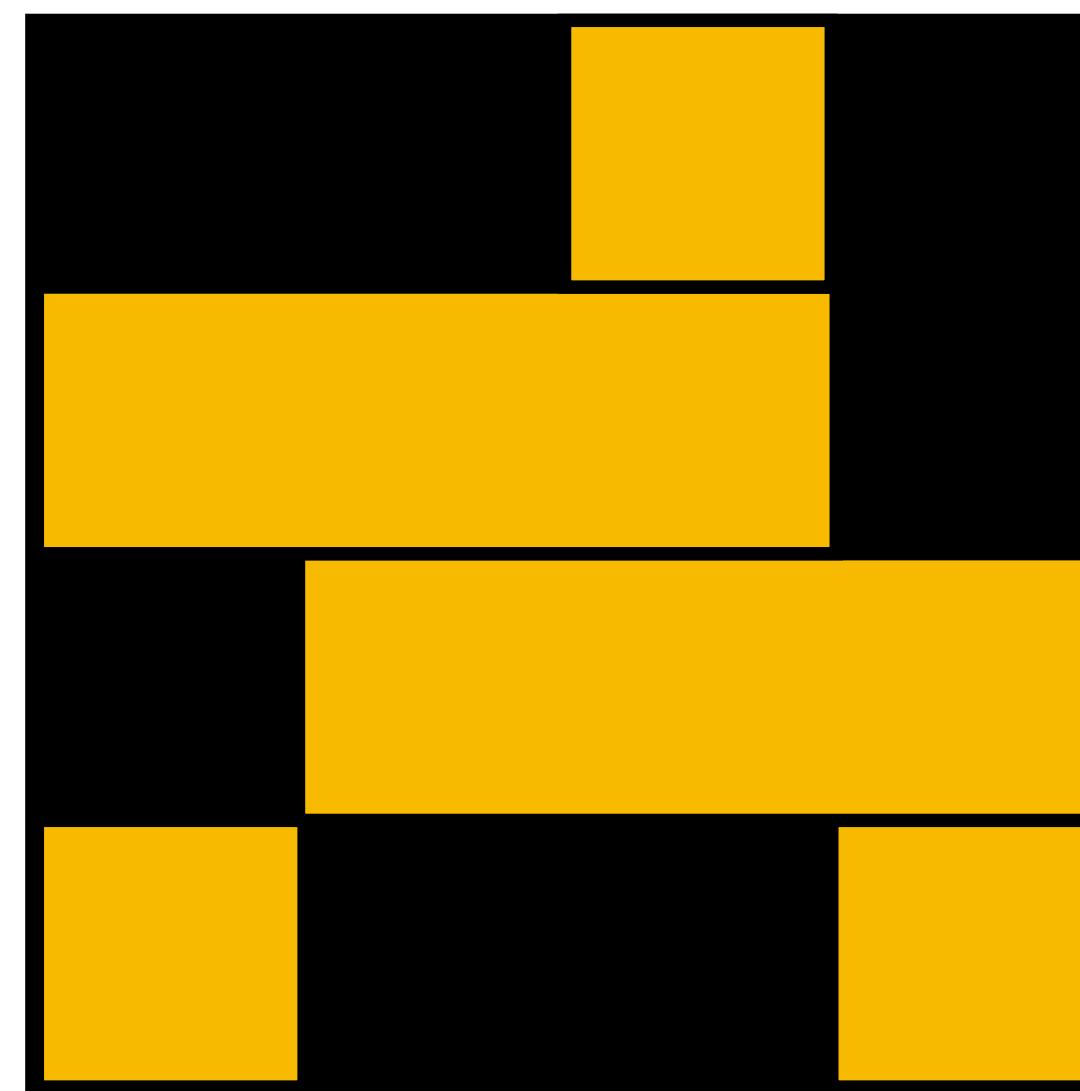
MALLOC



i !



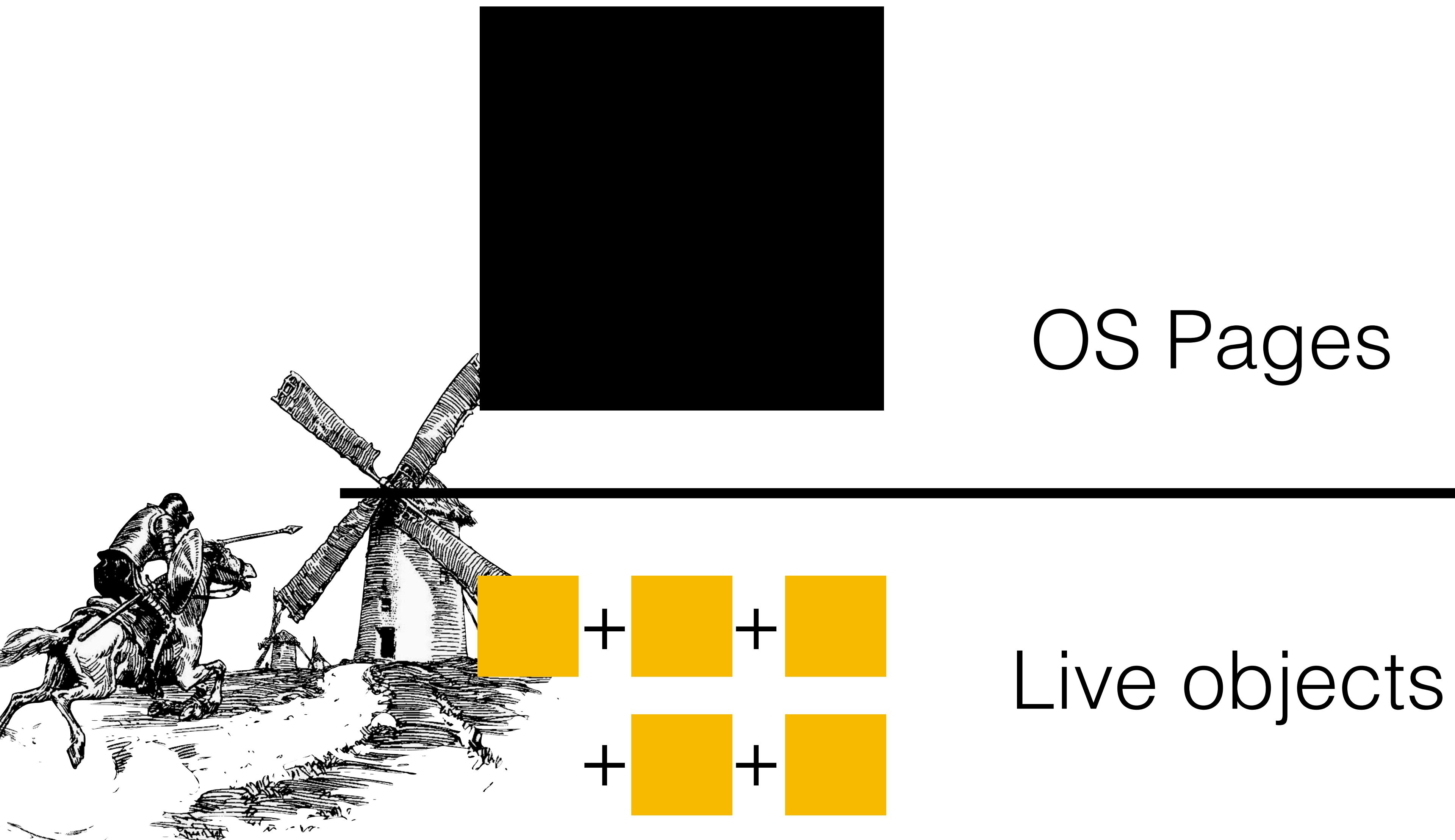
¡Fragmentación!



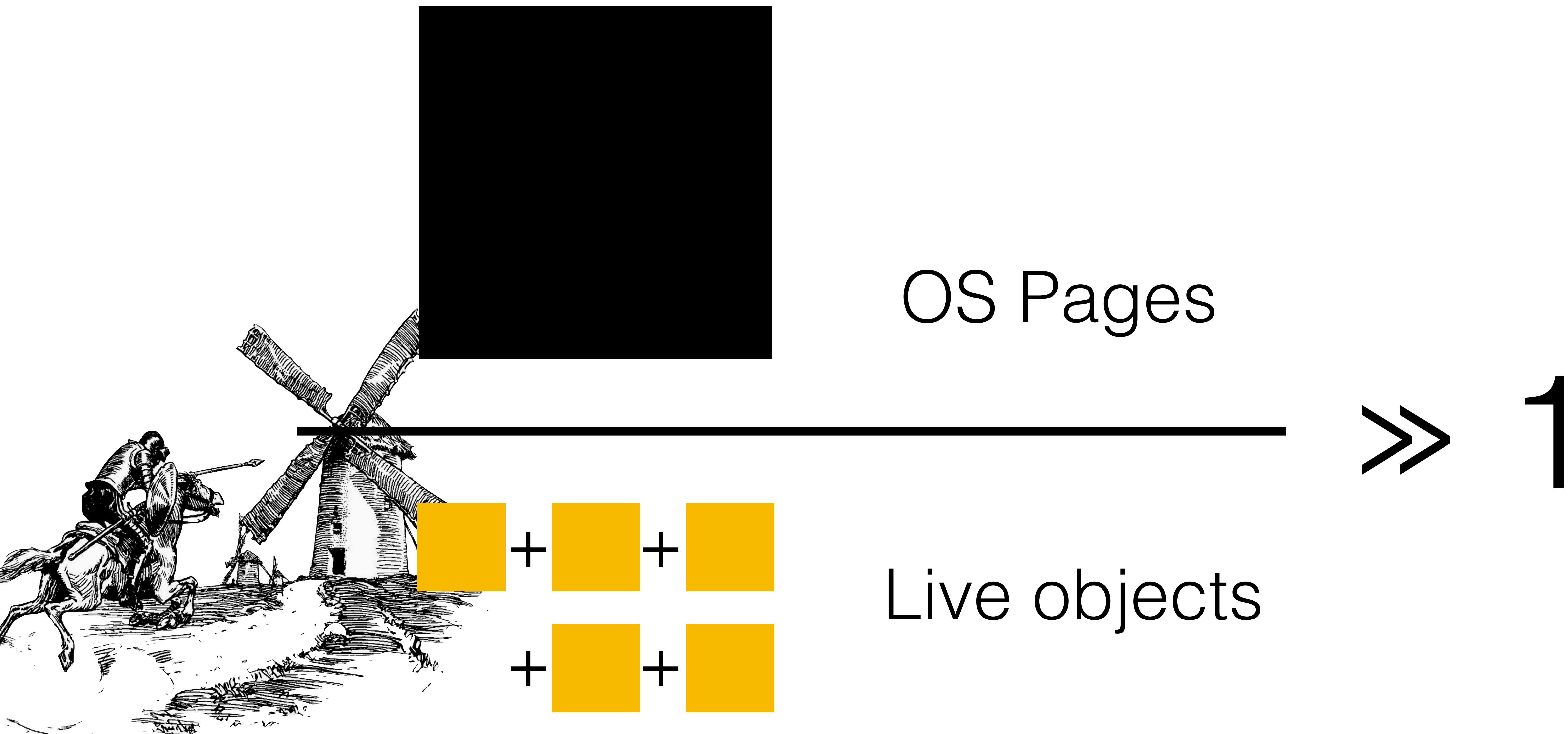
OS Page (“Spain”)

Live objects (malloc'd)

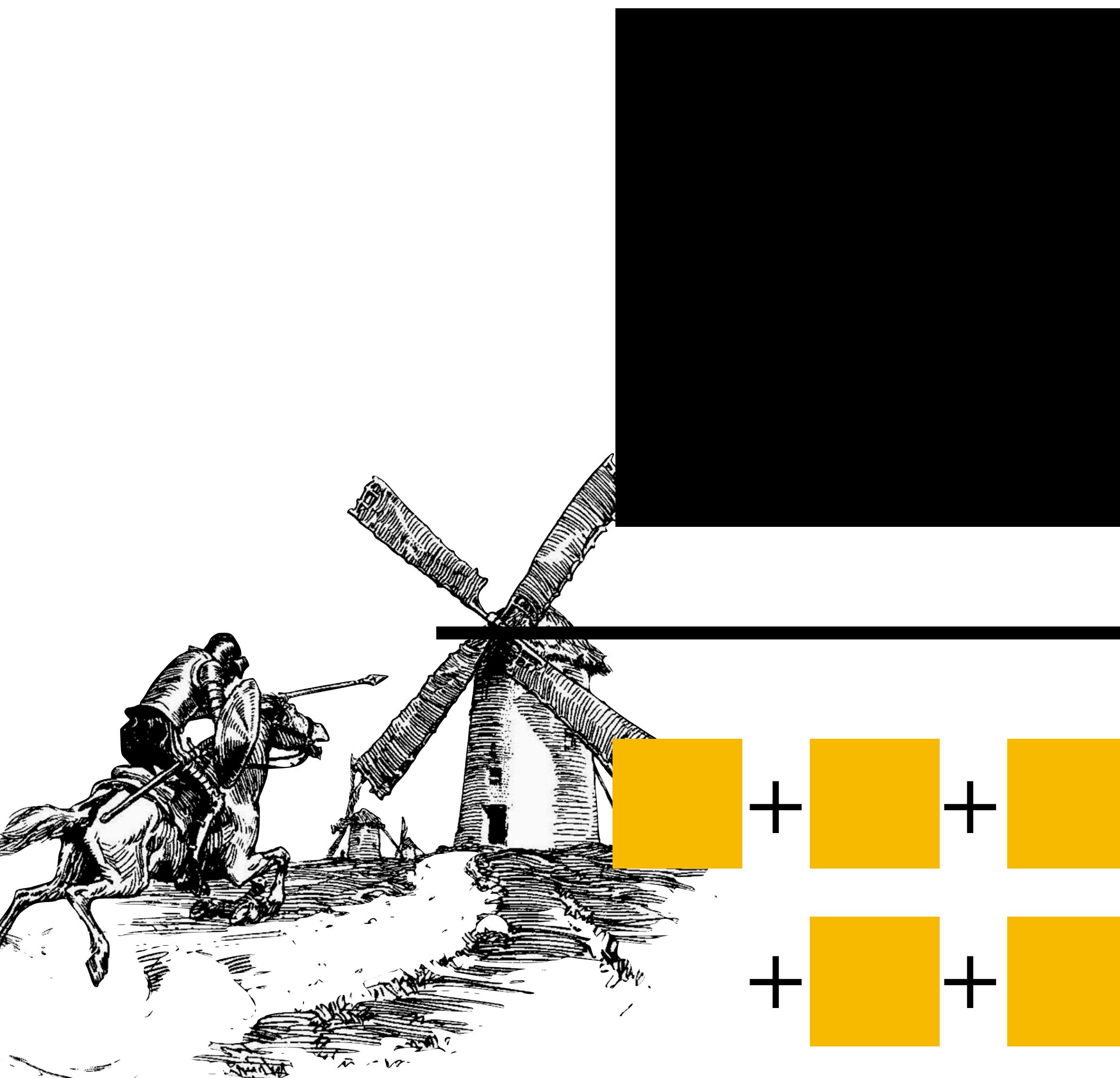
¡Fragmentación!



¡Fragmentación!



¡Fragmentación!

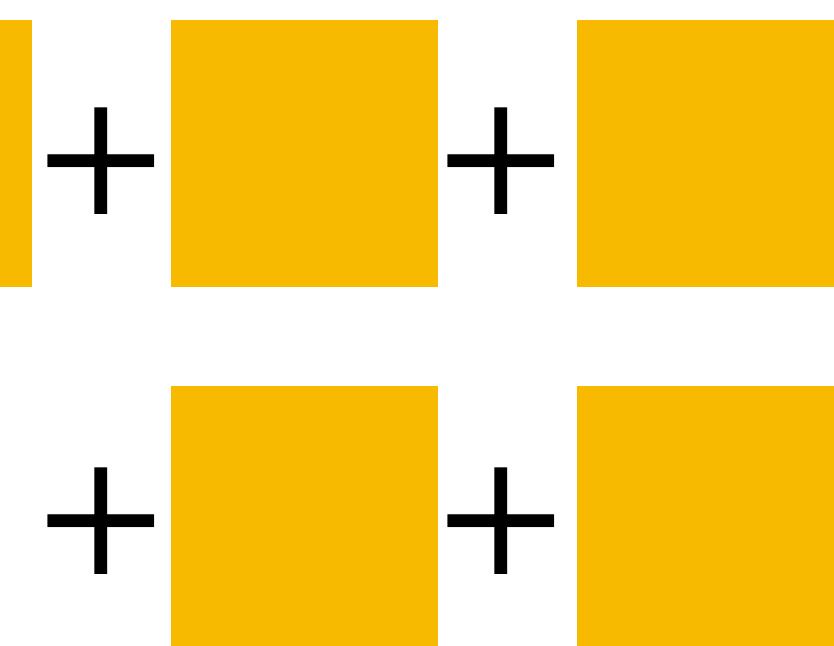


OS Pages

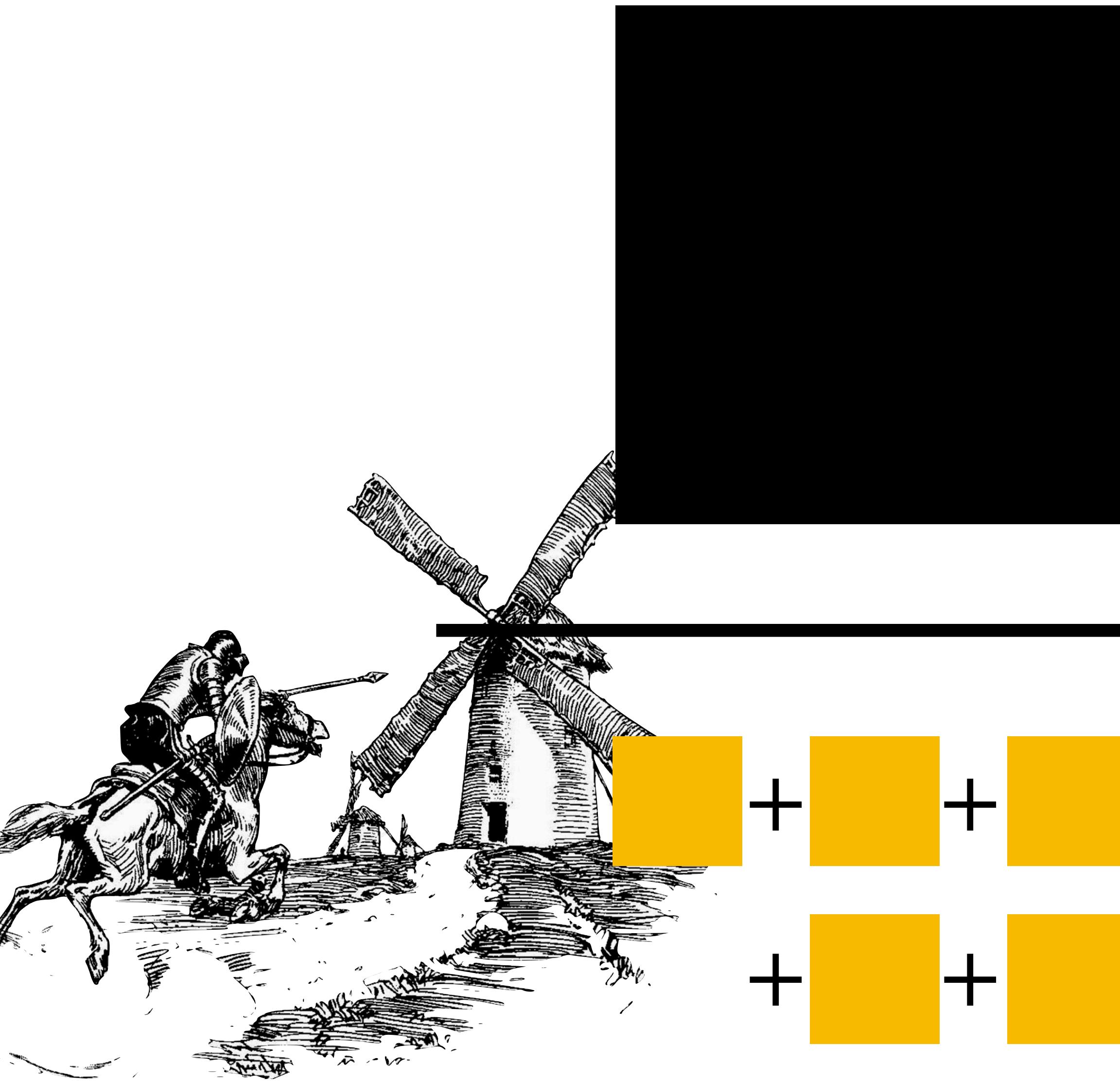
$O(\log \frac{\text{---}}{\text{---}})$ 13x!

[Robson '77]

Live objects



¿Compactación?



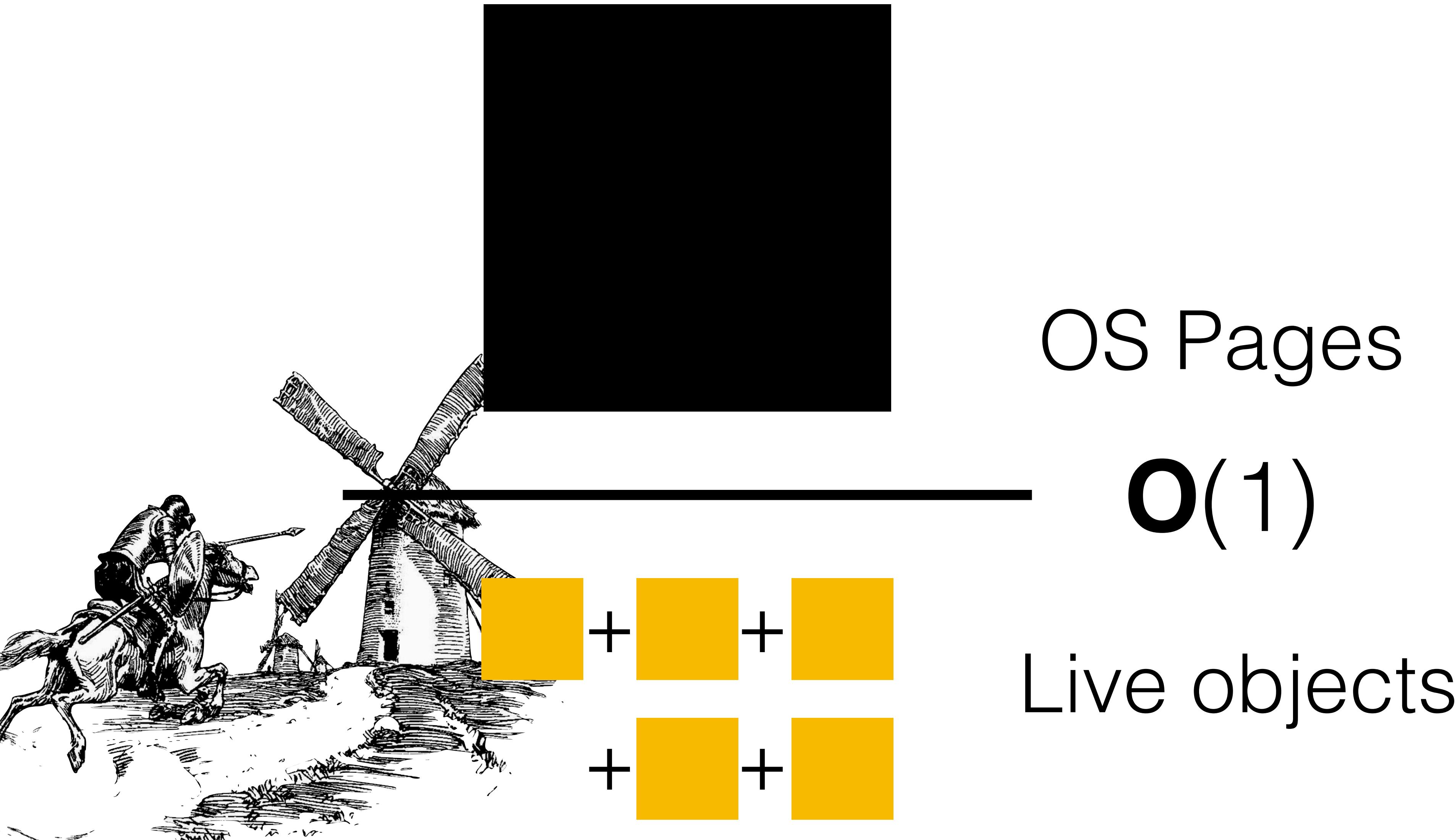
OS Pages

$O(\log \frac{\text{---}}{\text{---}})$ 13x!

[Robson '77]

Live objects

¿Compactación?









THE
C
PROGRAMMING
LANGUAGE

Brian W. Kernighan • Dennis M. Ritchie

PRENTICE HALL SOFTWARE SERIES



NG
E

M. Ritchie



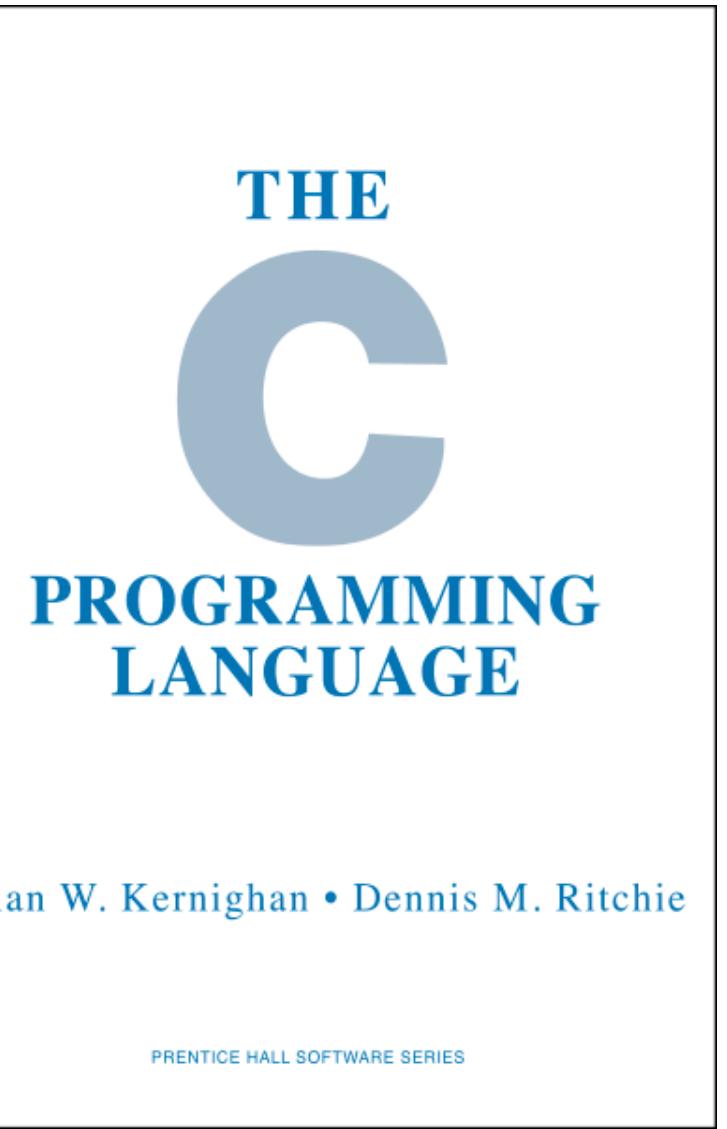
NG
E

M. Ritchie



NG
E

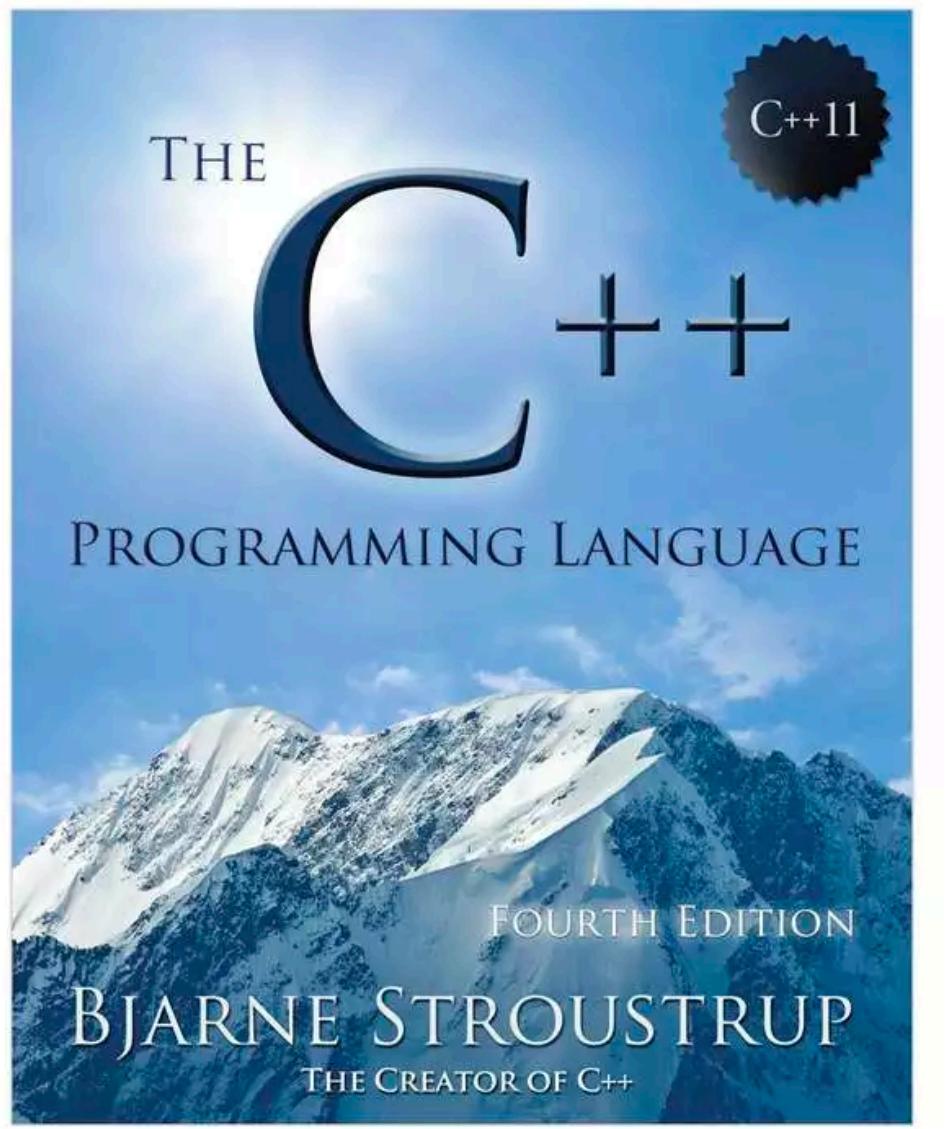
M. Ritchie



THE
C
PROGRAMMING
LANGUAGE

Brian W. Kernighan • Dennis M. Ritchie

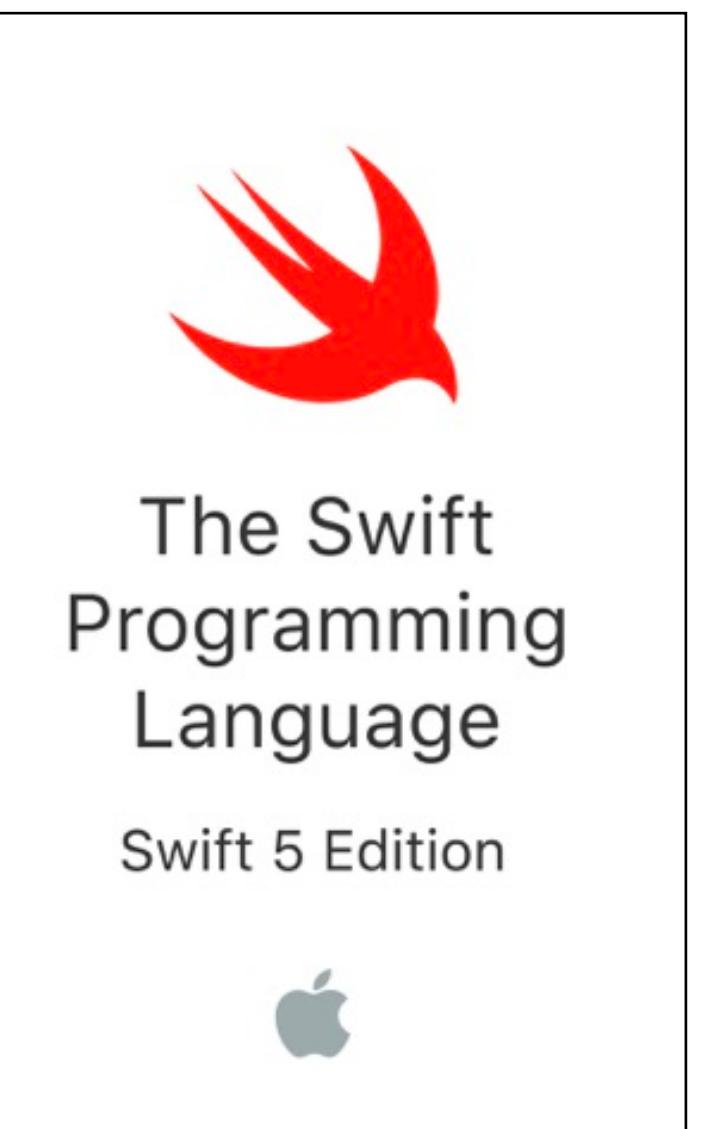
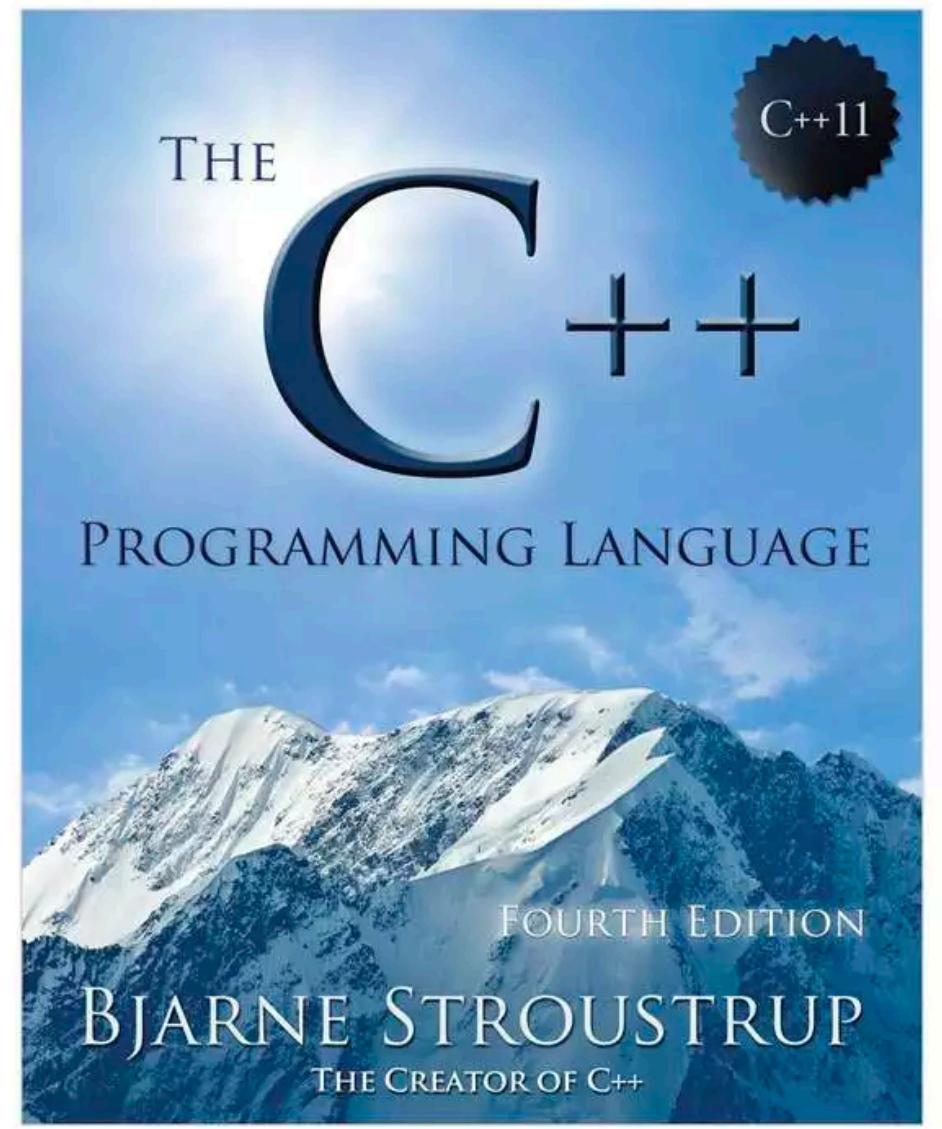
PRENTICE HALL SOFTWARE SERIES

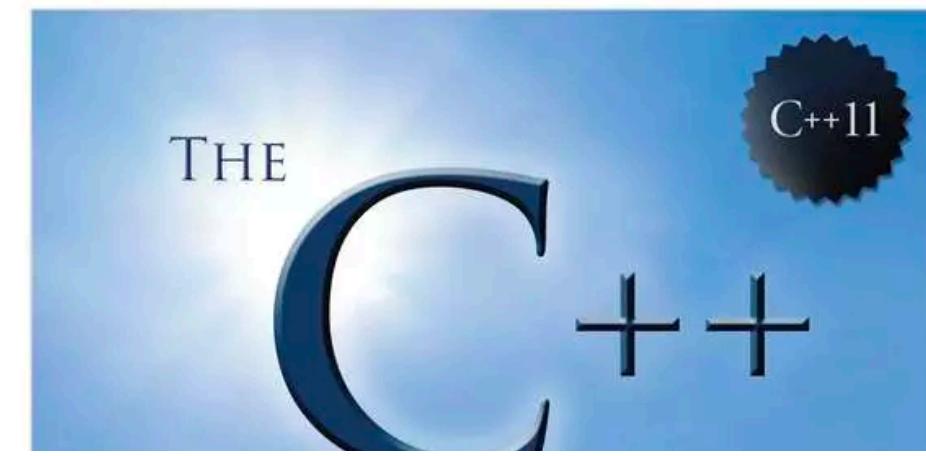


THE
C
PROGRAMMING
LANGUAGE

Brian W. Kernighan • Dennis M. Ritchie

PRENTICE HALL SOFTWARE SERIES





[Lattner, 2016]

Why not a tracing GC?

- Native interoperability with unmanaged code
- Deterministic destruction provides:
 - No “finalizer problems” like resurrection, threading, etc.
 - Deterministic performance: can test/debug performance stutters
- Performance:
 - GC use ~3-4x more memory than ARC to achieve good performance
 - Memory usage is very important for mobile and cloud apps
 - Incremental/concurrent GCs slow the mutator like ARC does

Quantifying the Performance of Garbage Collection vs. Explicit Memory Management
Matthew Hertz, Emery D. Berger. OOPSLA'05

Native interoperability with unmanaged code

Deterministic destruction provides:

- No “finalizer problems” like resurrection, threading, etc.
- Deterministic performance: can test/debug performance stutters

Performance:

- GC use ~3-4x more memory than ARC to achieve good performance
- Memory usage is very important for mobile and cloud apps
- Incremental/concurrent GCs slow the mutator like ARC does

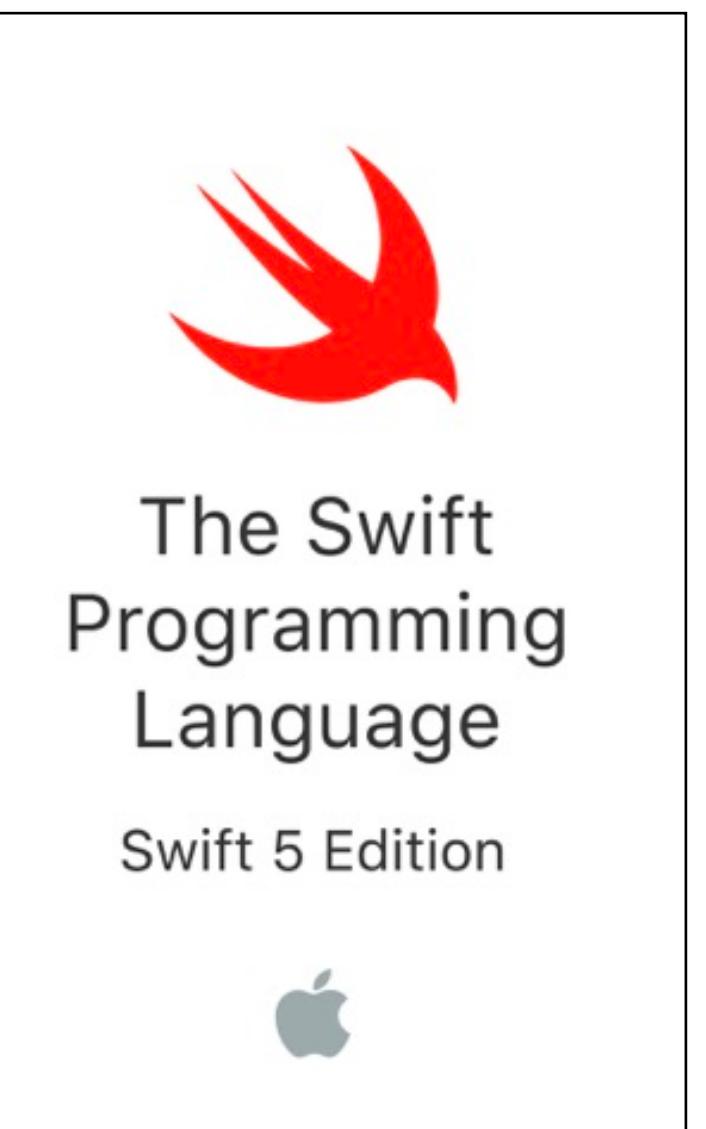
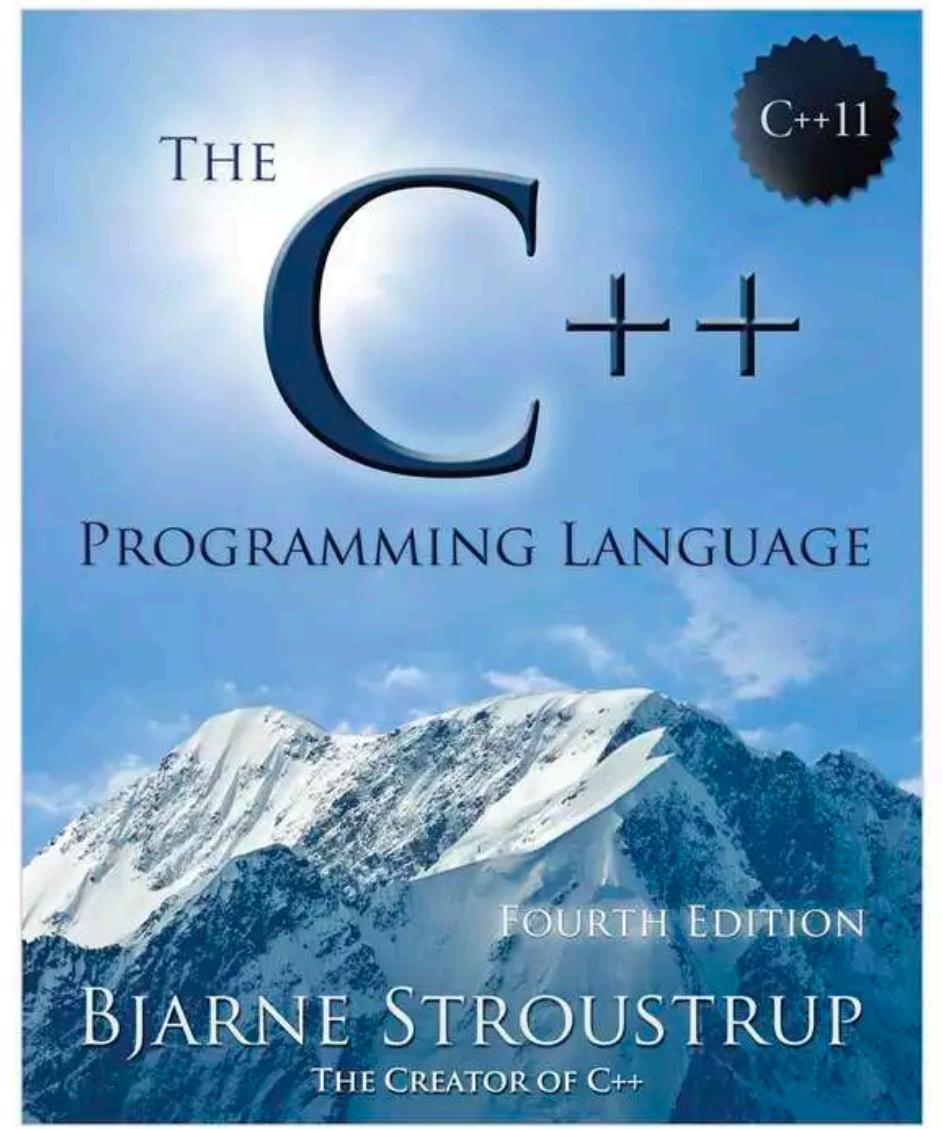
Quantifying the Performance of Garbage Collection vs. Explicit Memory Management

Matthew Hertz, Emery D. Berger. OOPSLA'05

THE
C
PROGRAMMING
LANGUAGE

Brian W. Kernighan • Dennis M. Ritchie

PRENTICE HALL SOFTWARE SERIES



The Swift
Programming
Language

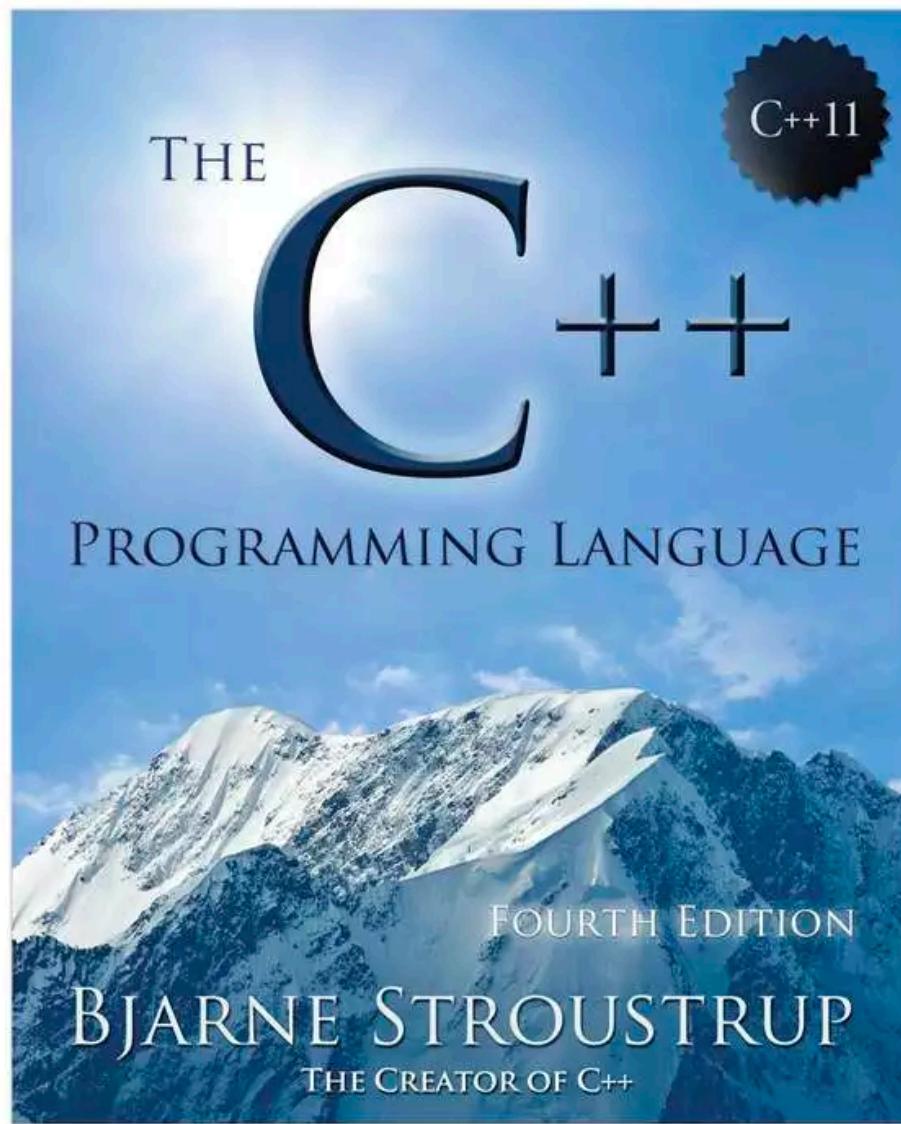
Swift 5 Edition



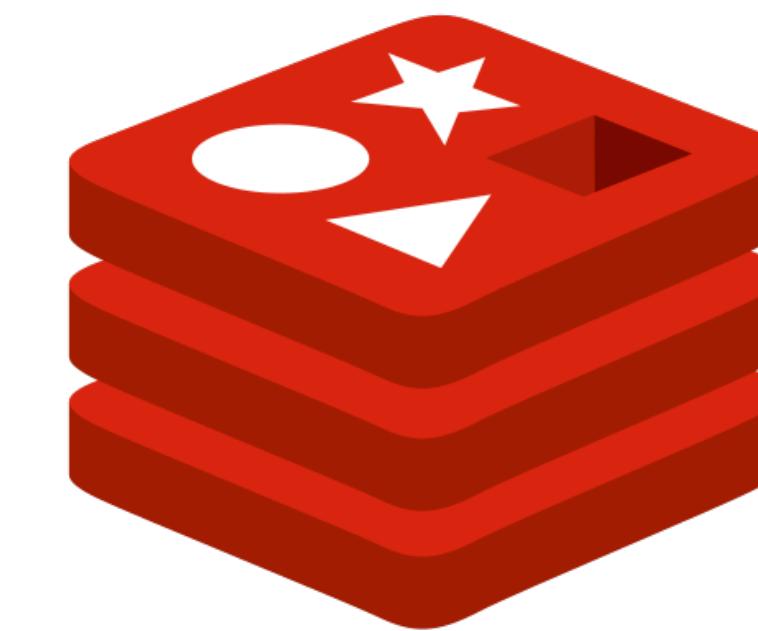
THE
C
PROGRAMMING
LANGUAGE

Brian W. Kernighan • Dennis M. Ritchie

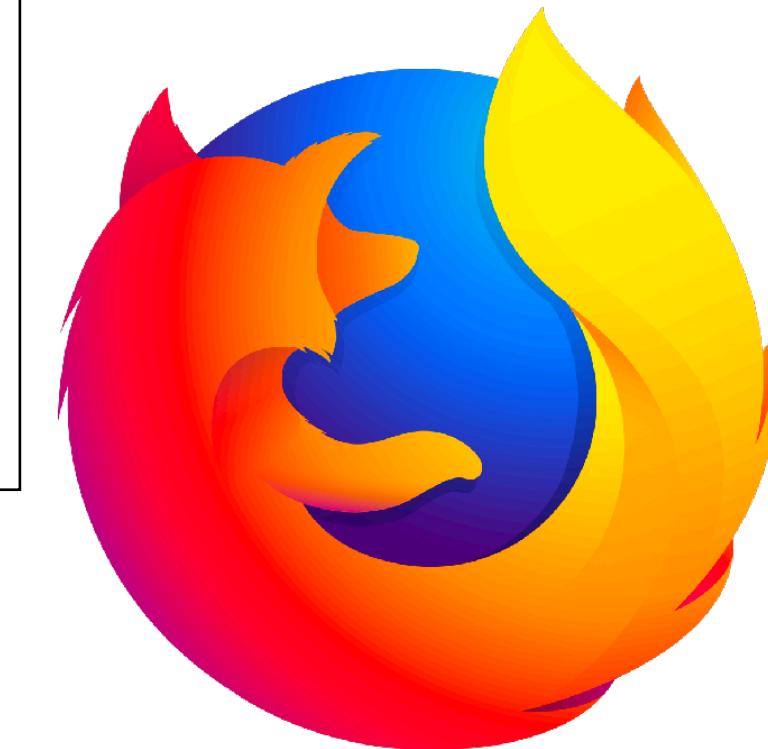
PRENTICE HALL SOFTWARE SERIES



The Swift
Programming
Language
Swift 5 Edition



redis

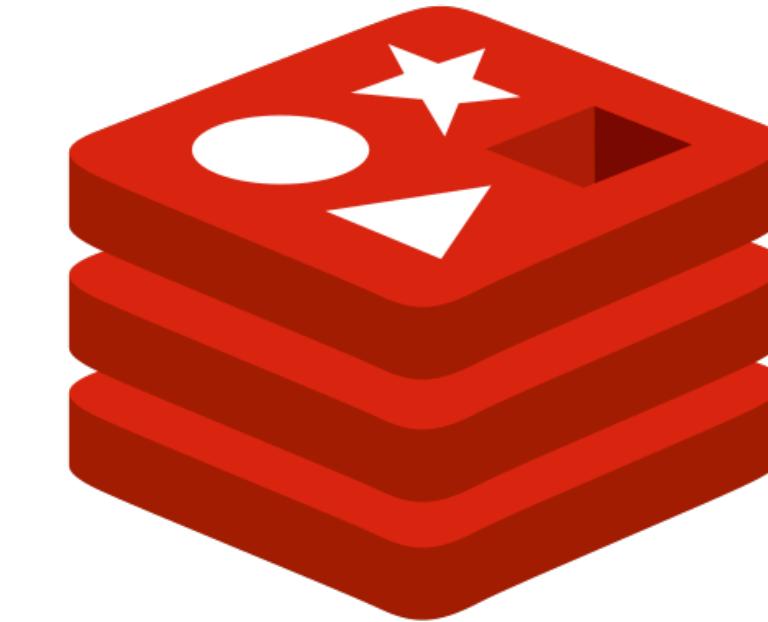
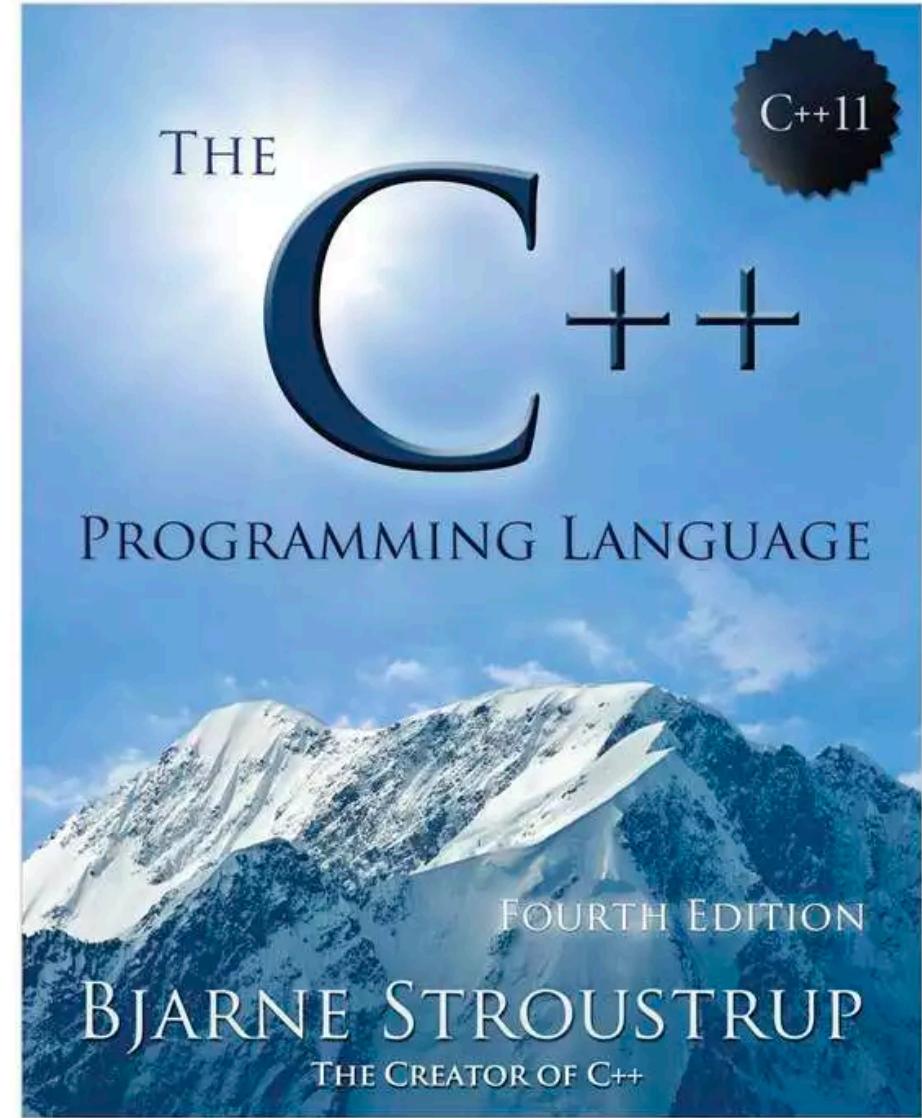


MySQL™

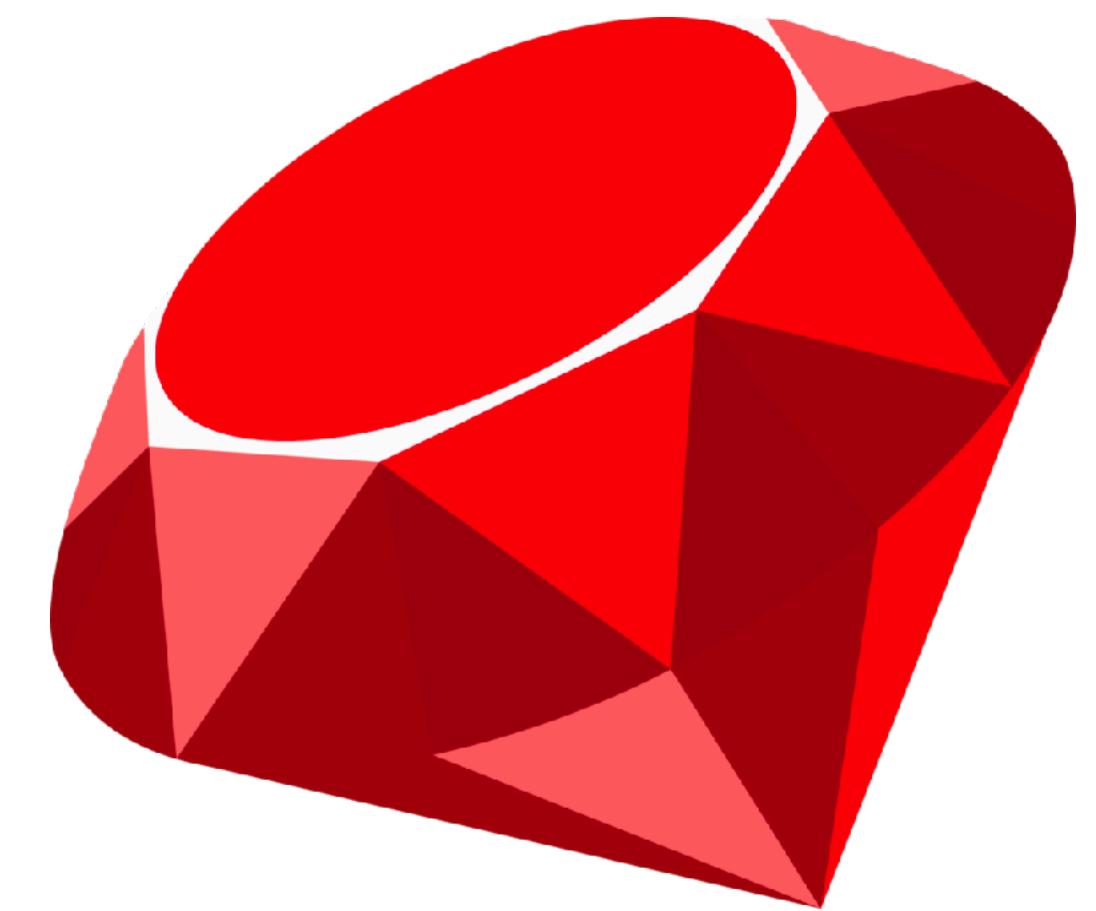


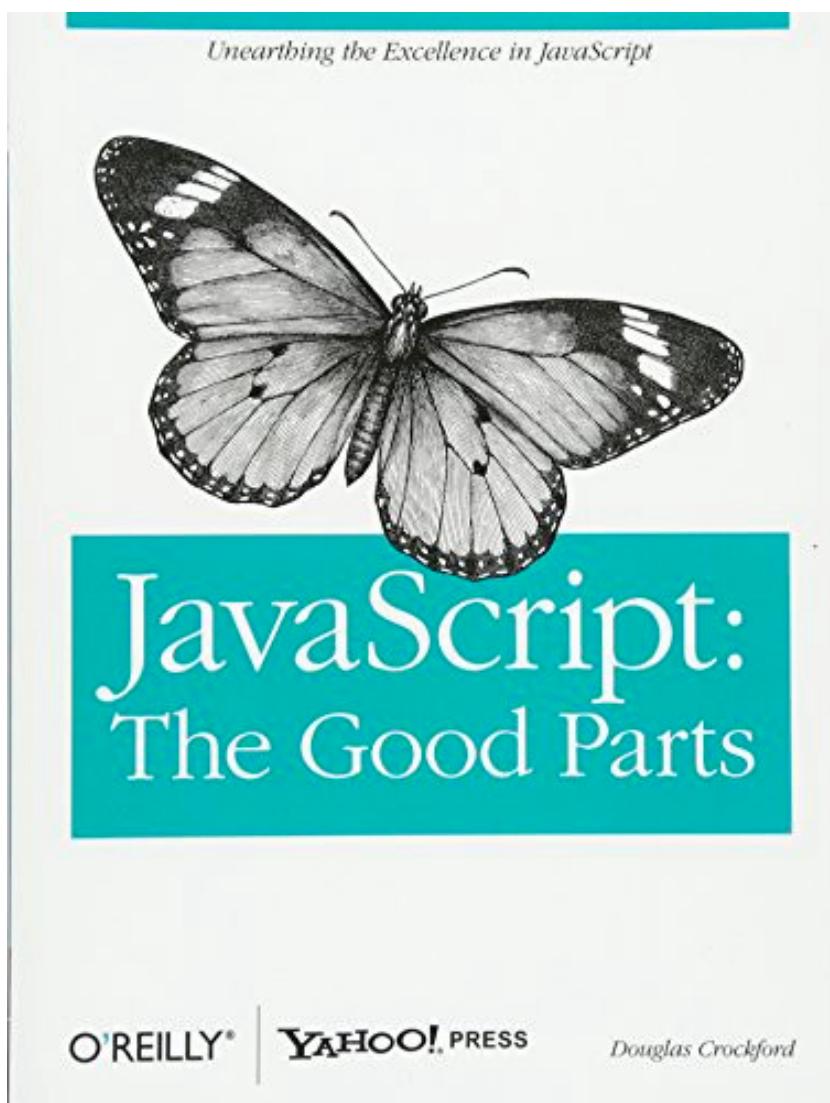
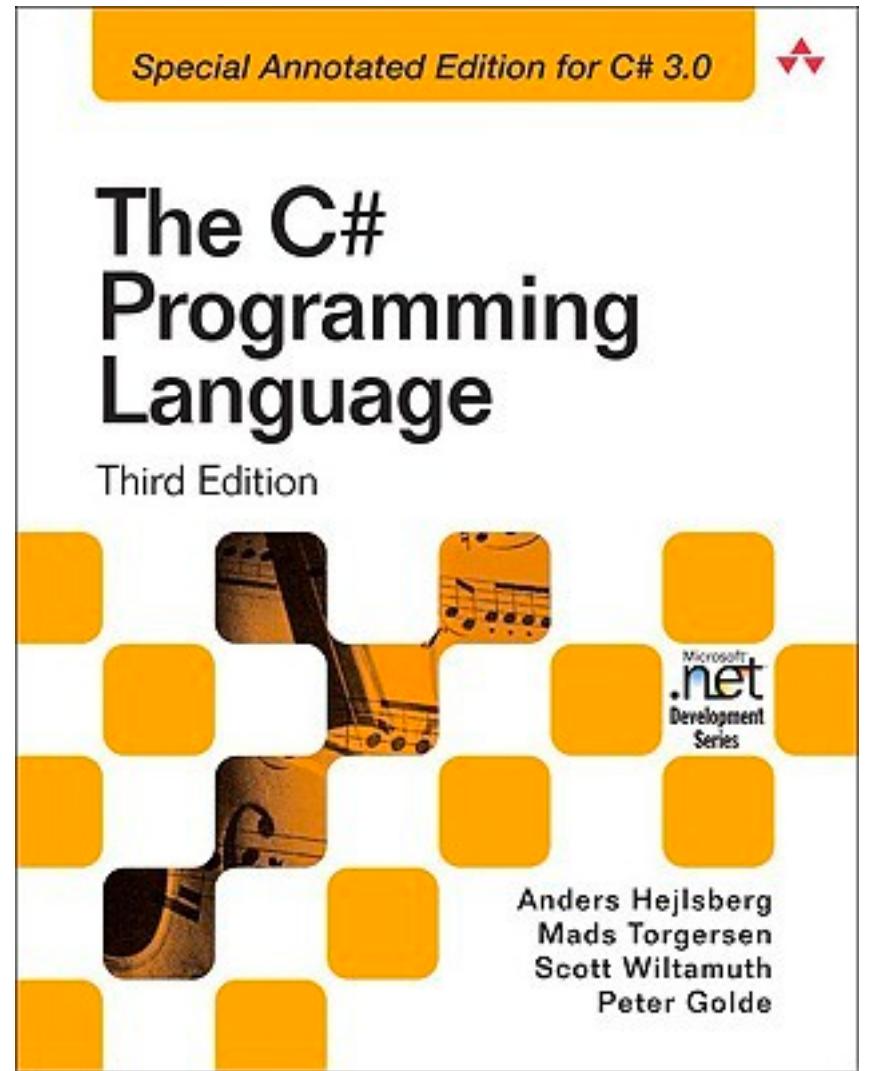
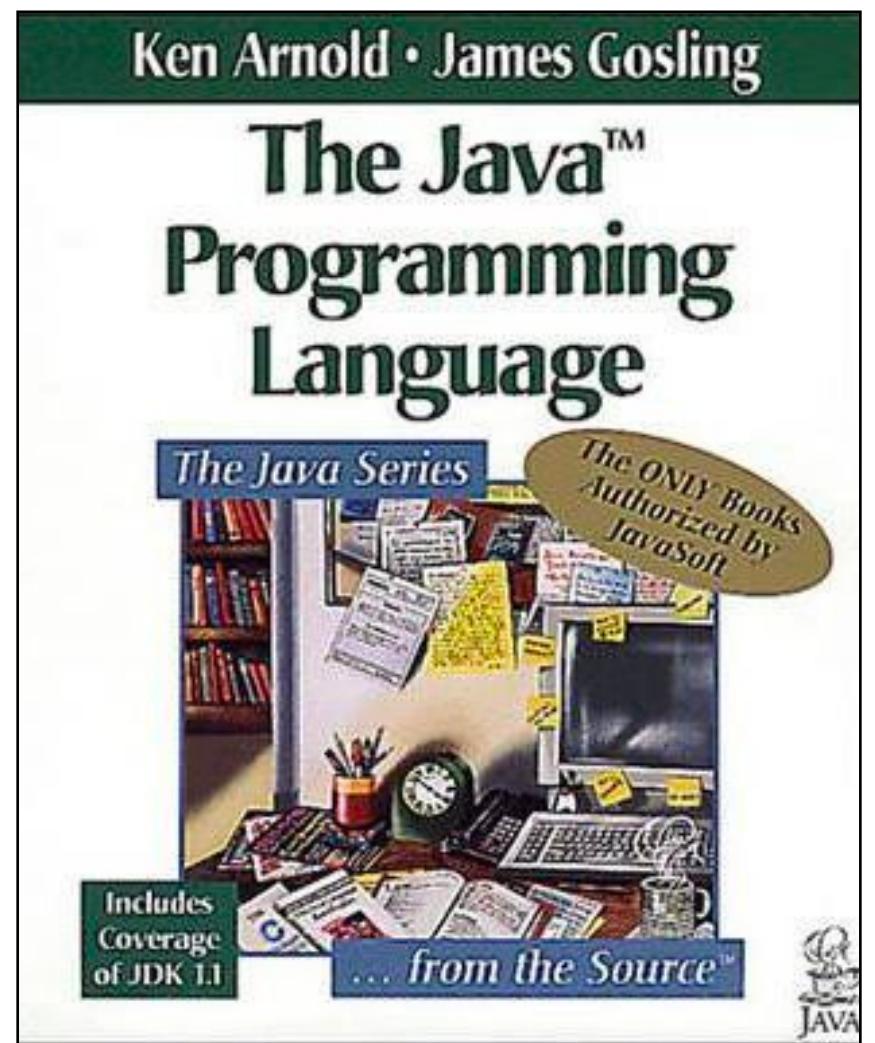
Office

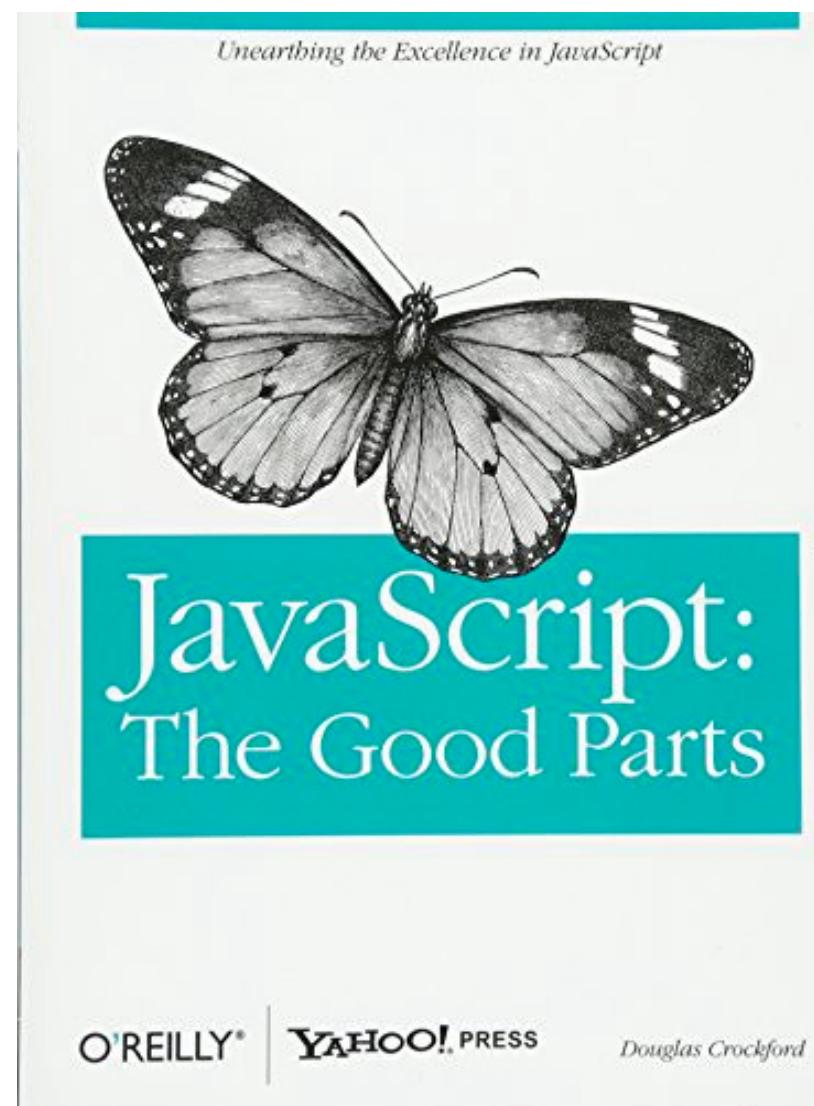
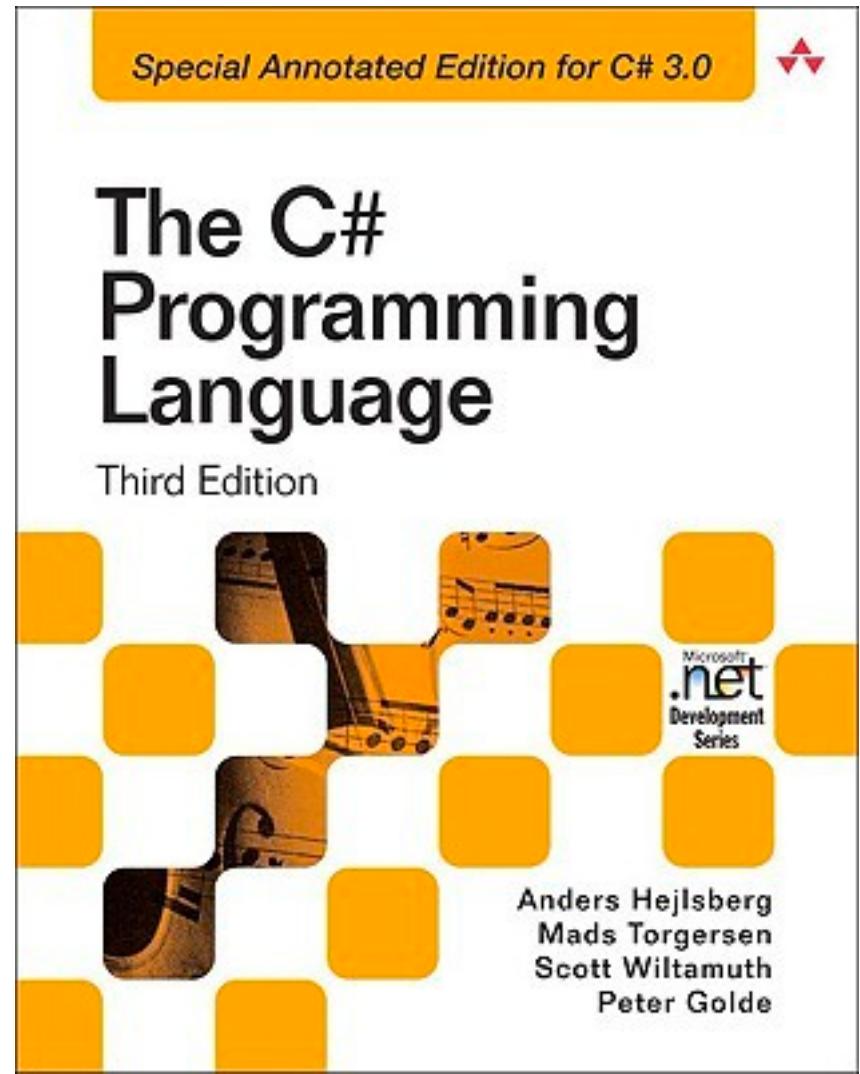
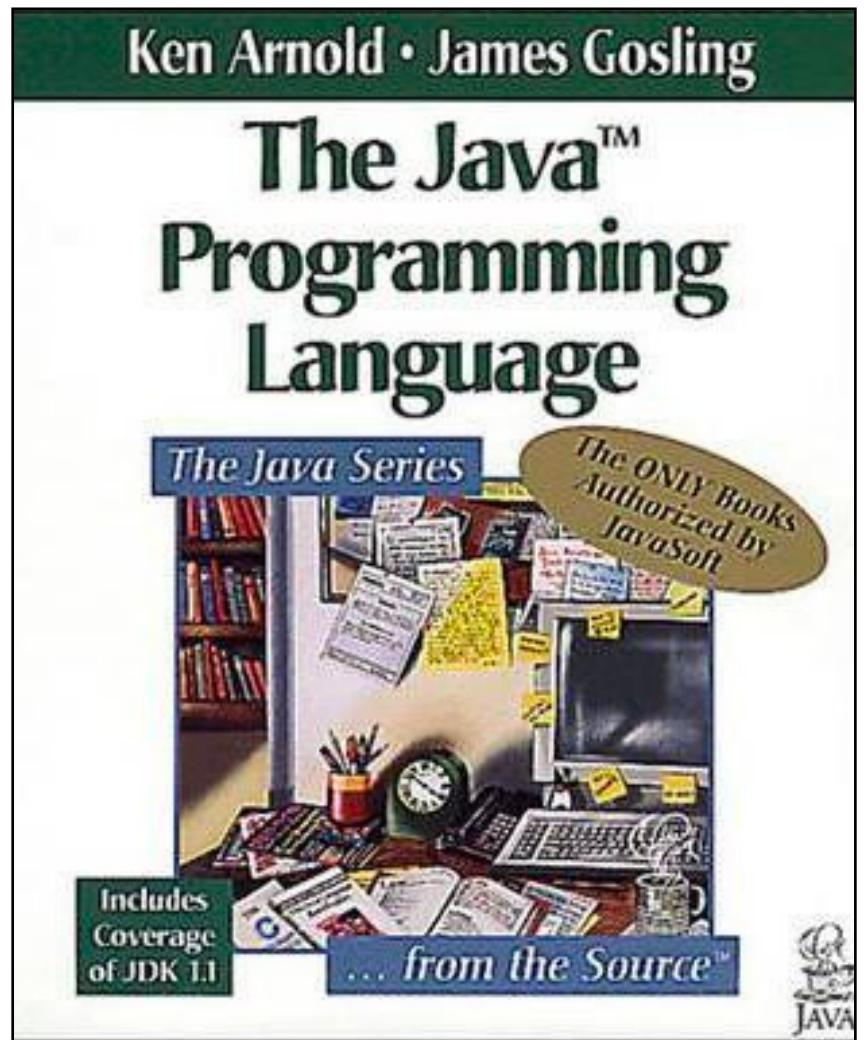




redis

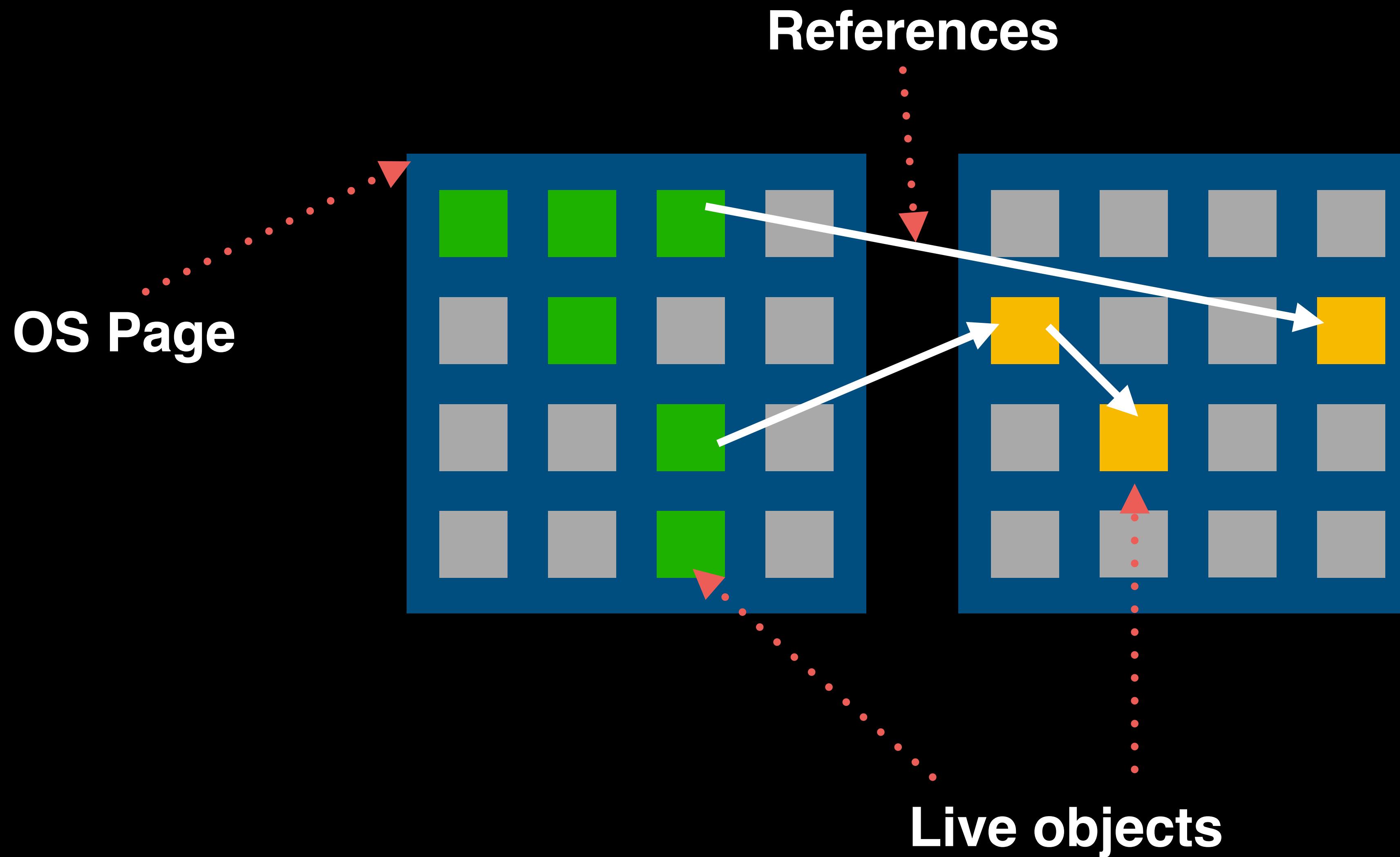


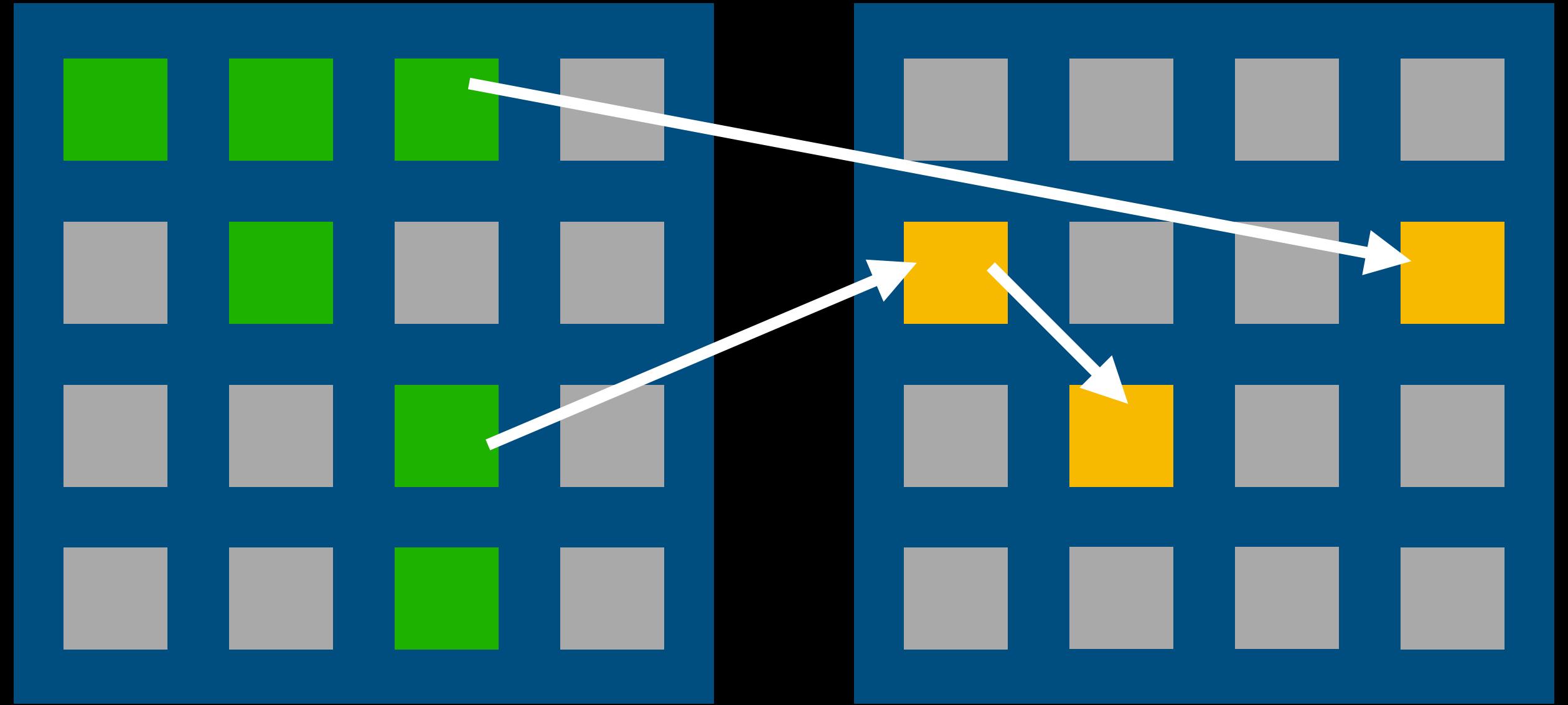


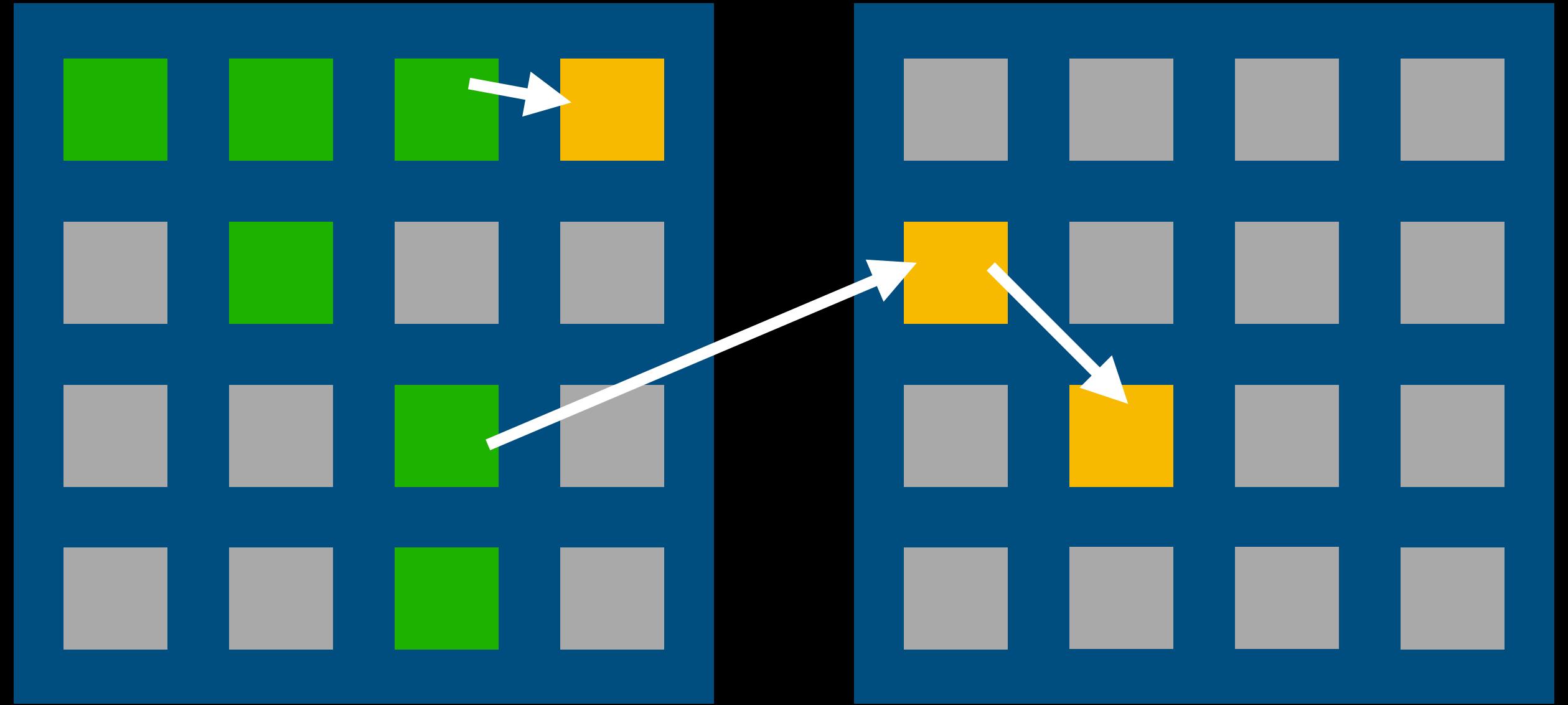


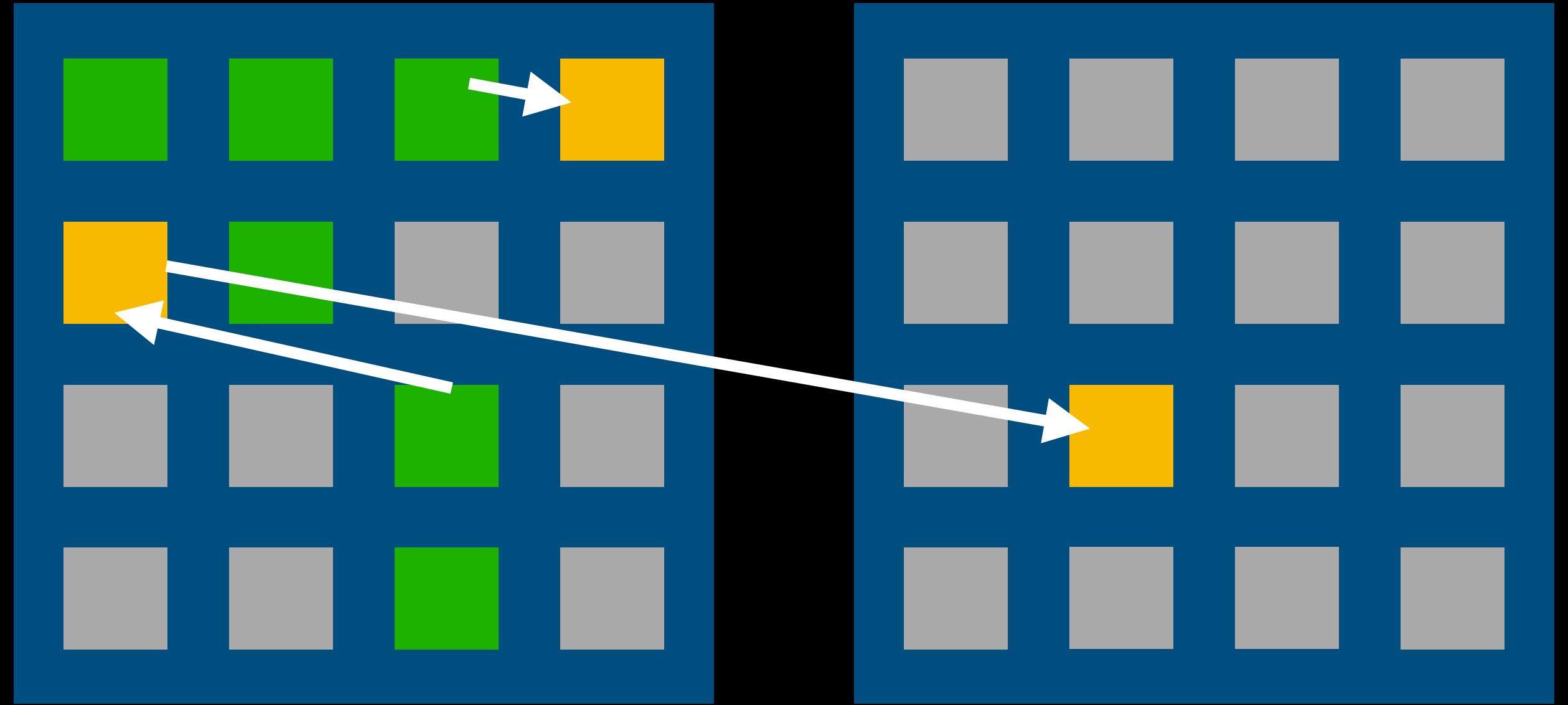
¡Compactación!

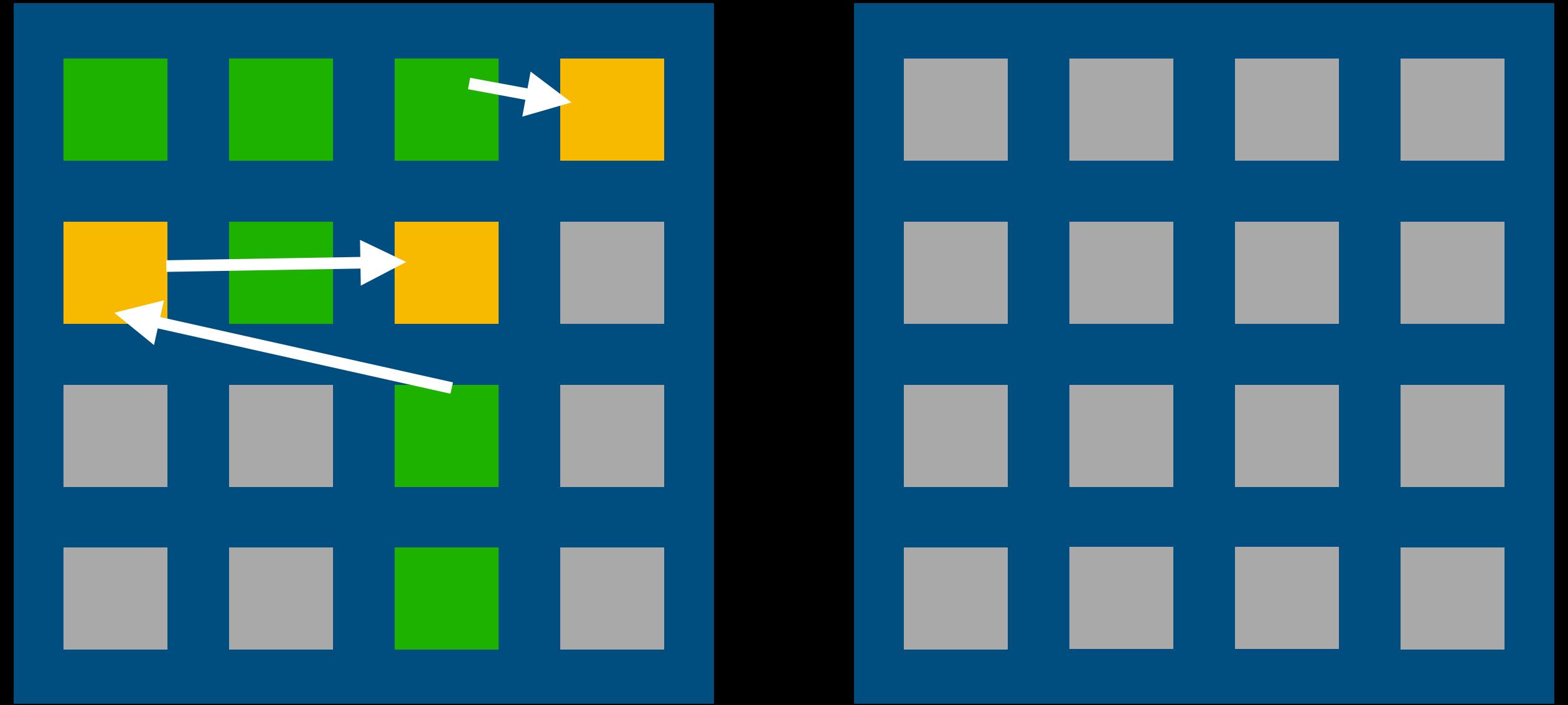
Compaction In Action

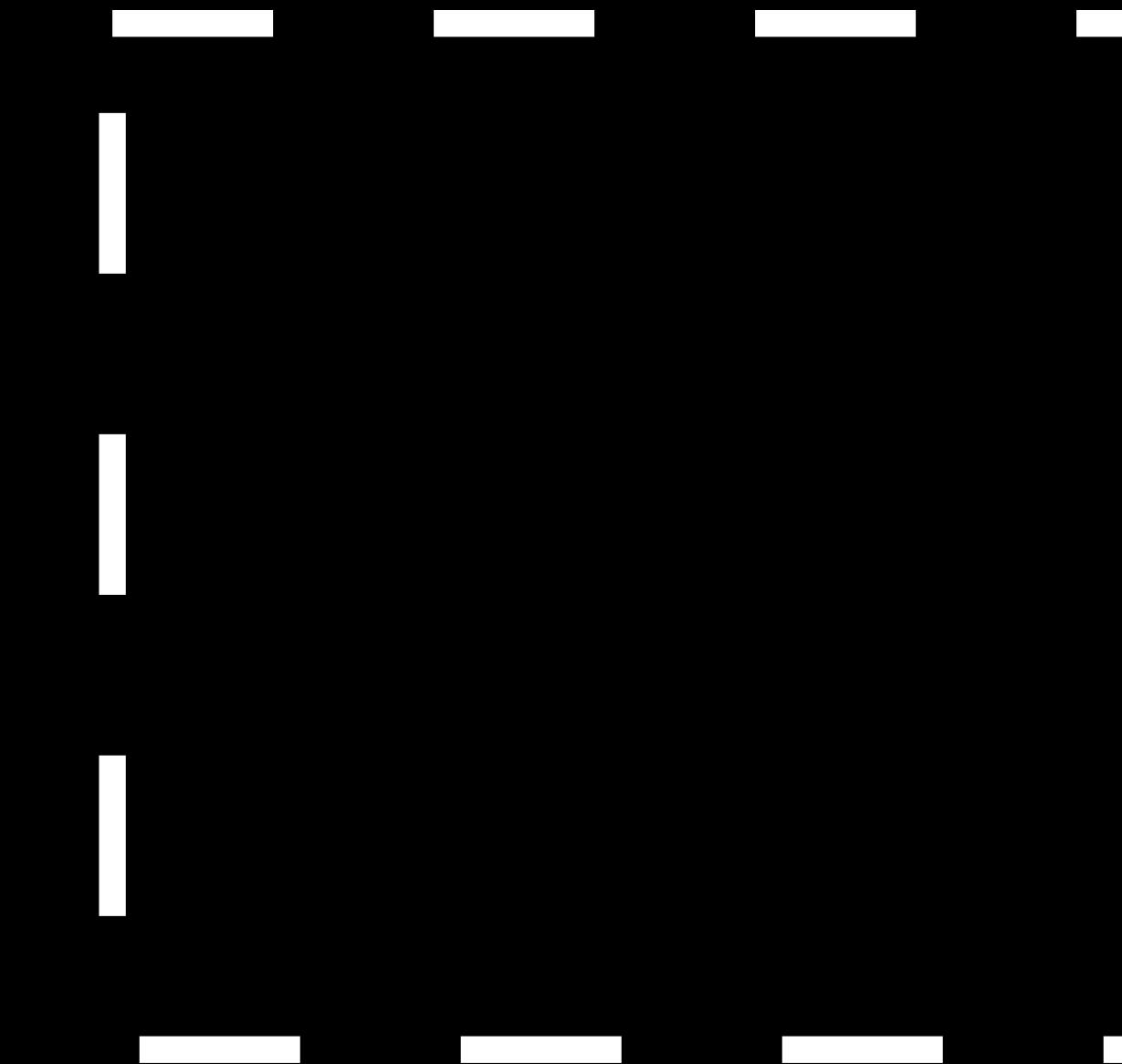
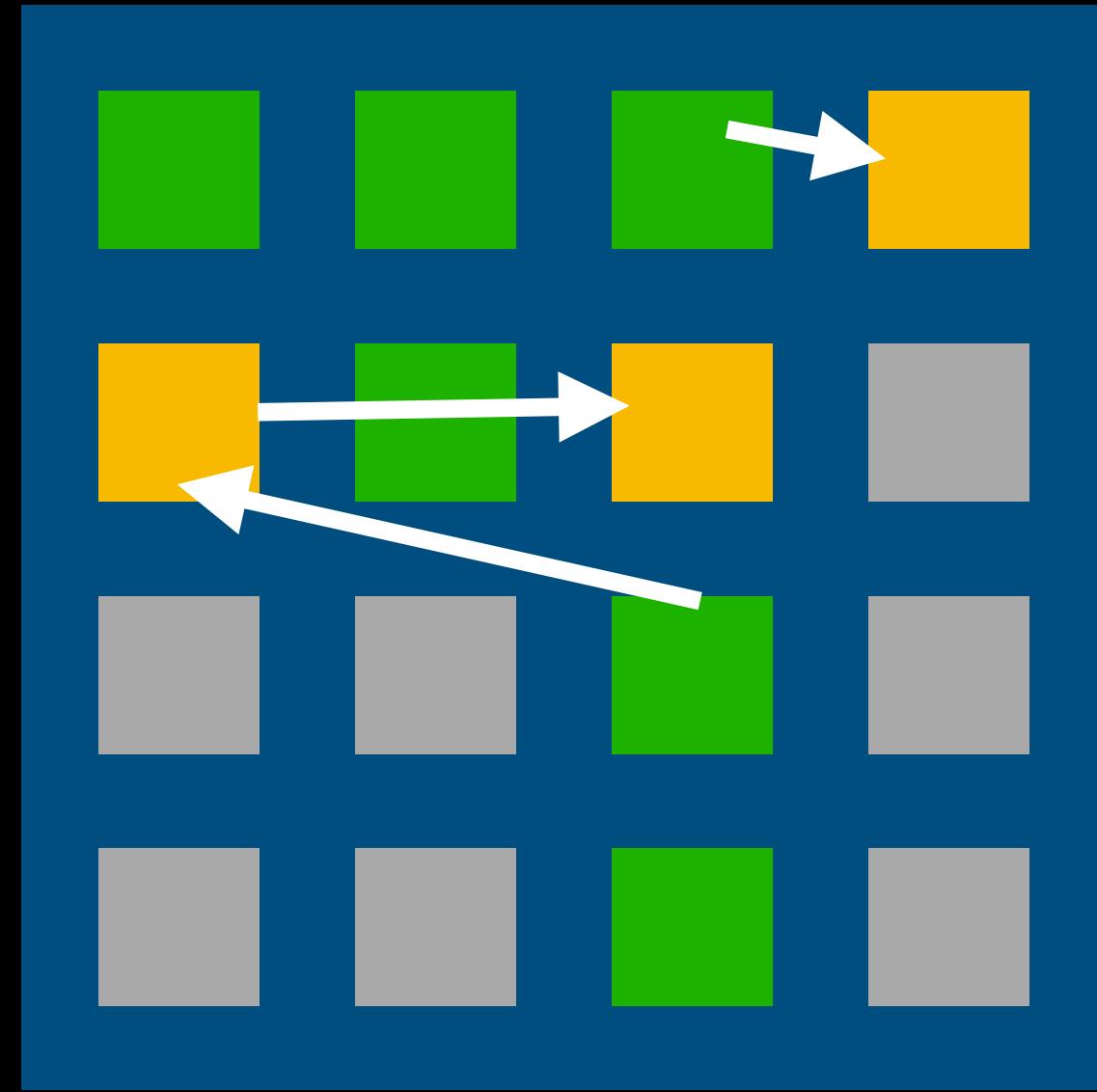












Page returned
to the OS

```
$ clang++ -o yolo main.cc
```



```
$ clang++ -o yolo main.cc  
$ strip yolo
```



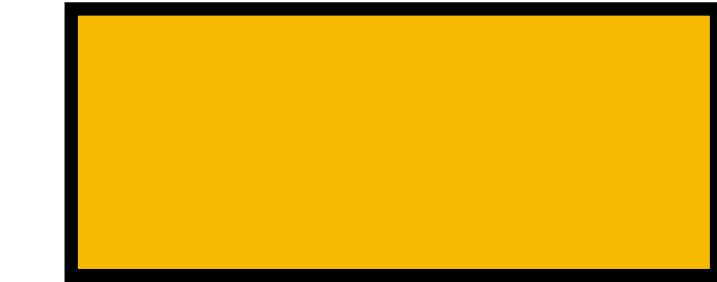
```
$ clang++ -o yolo main.cc  
$ strip yolo  
$ ./yolo
```



**No way to precisely
distinguish pointers
from integers**

```
$ clang++ -o yolo main.cc  
$ strip yolo  
$ ./yolo
```

0xDEADC000



**No way to precisely
distinguish pointers
from integers**



```
$ clang++ -o yolo main.cc
```

```
$ strip yolo
```

```
$ ./yolo
```

0xDEADC000



**No way to precisely
distinguish pointers
from integers**

0xBEEFC000



0xDEADC000

```
$ clang++ -o yolo main.cc
```

```
$ strip yolo
```

```
$ ./yolo
```

0xDEADC000



0xBEEFC000



No way to precisely
distinguish pointers
from integers

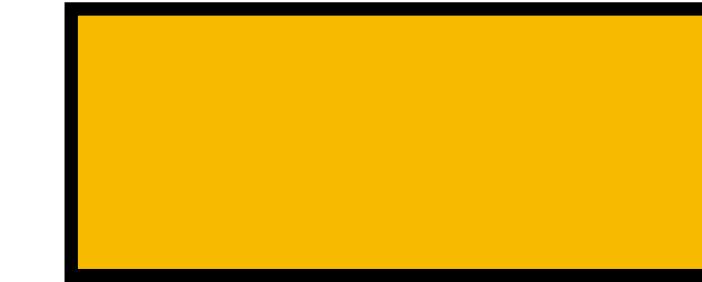


```
$ clang++ -o yolo main.cc  
$ strip yolo  
$ ./yolo
```

0xDEADC000



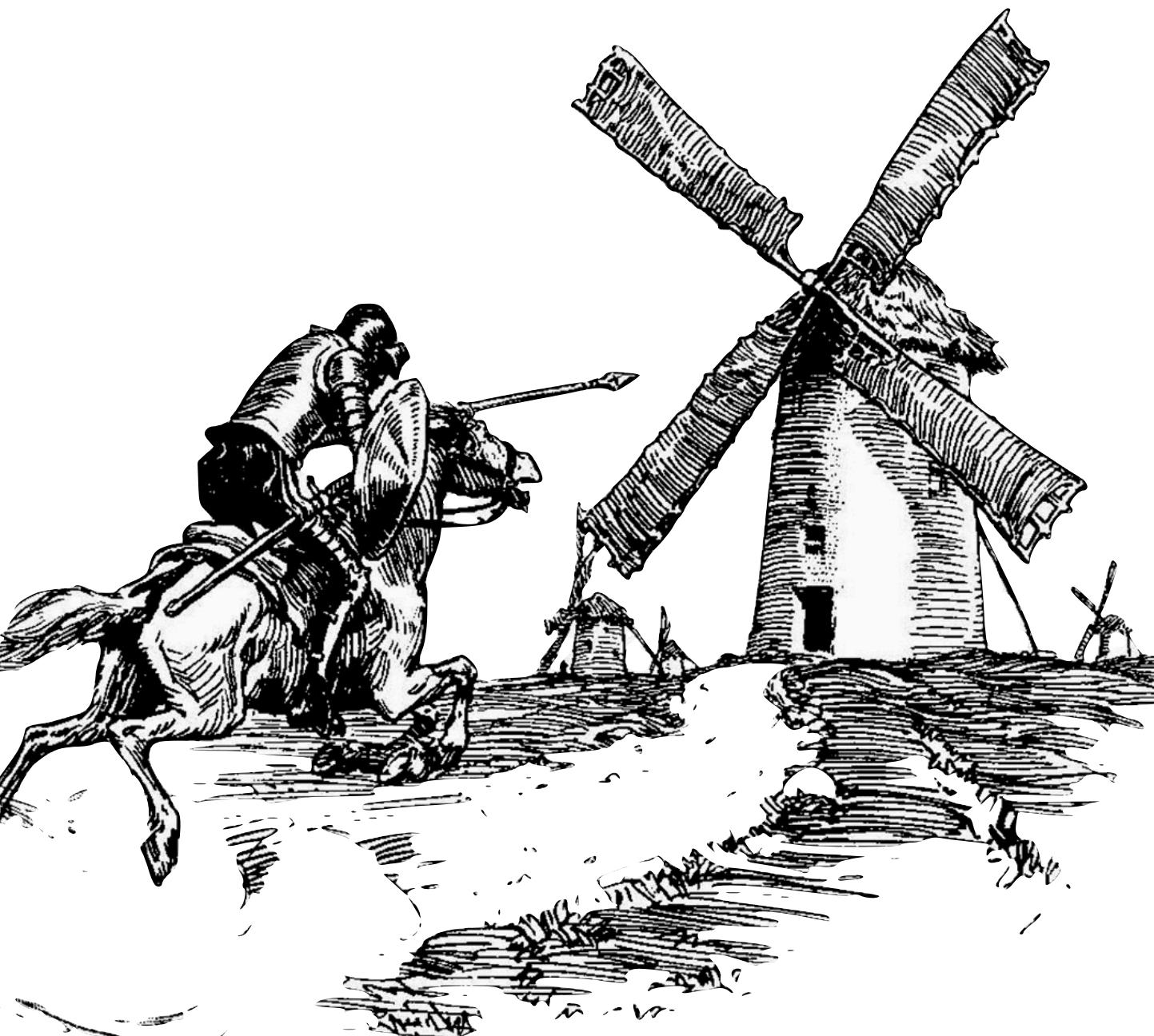
0xDEADC000



0xBEEFC000



**No way to precisely
distinguish pointers
from integers**



```
union tiny
{
    int * ptr;
    uintptr_t flag;
};

tiny x;

// initialize
x.ptr = new int;

// set flag true
x.flag |= 1;
```

```
union tiny
{
    int * ptr;
    uintptr_t flag;
};

tiny x;
```

0xDEADC001



0xDEADC000



```
// initialize
x.ptr = new int;
```

0xBEEFC000



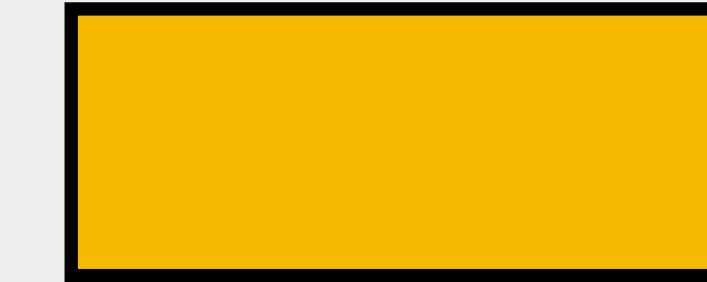
```
// set flag true
x.flag |= 1;
```

```
union tiny
{
    int * ptr;
    uintptr_t flag;
};
tiny x;
```

0xDEADC001



0xDEADC000



```
// initialize
x.ptr = new int;
```

0xBEEFC000



```
// set flag true
x.flag |= 1;
```



MESH

Compaction without
Relocation for C/C++

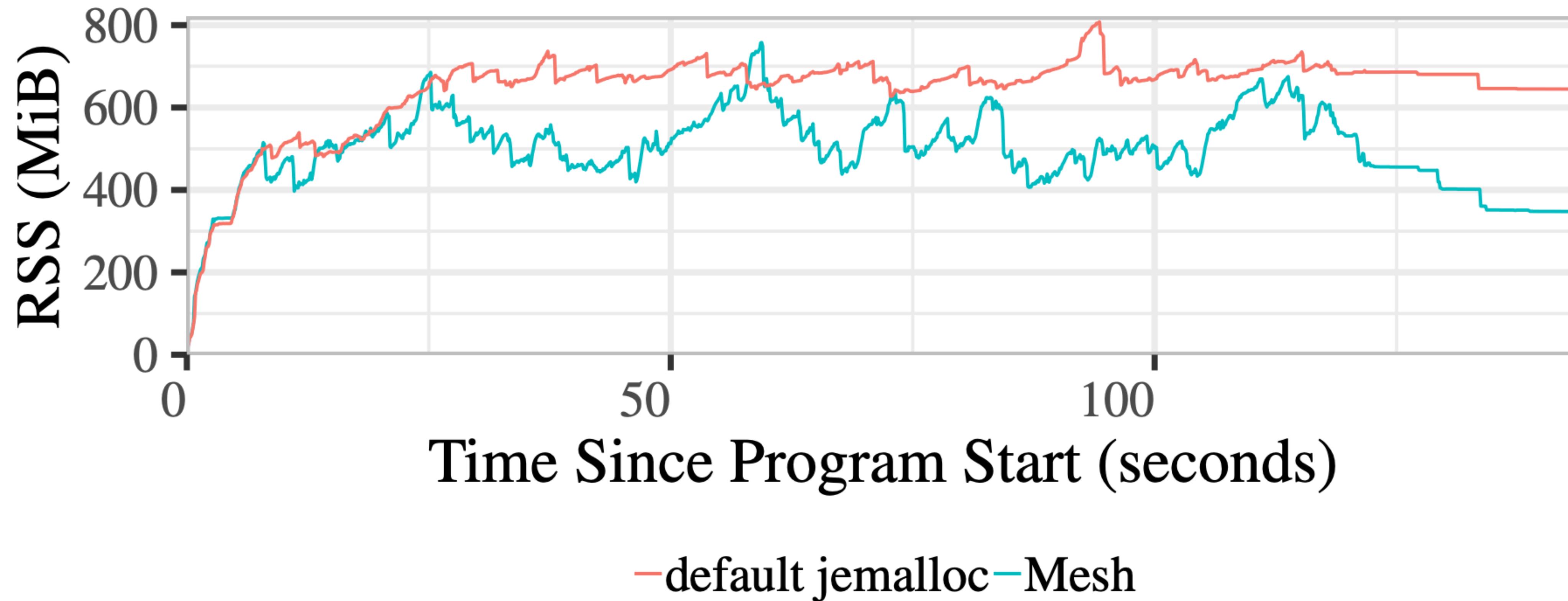
No code changes

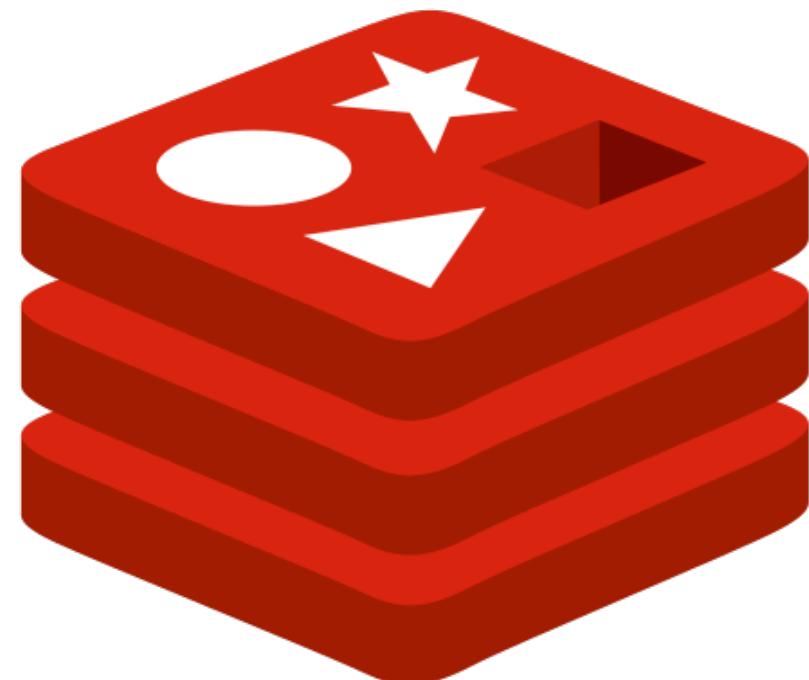
No recompilation

LD_PRELOAD and go

17% heap size reduction

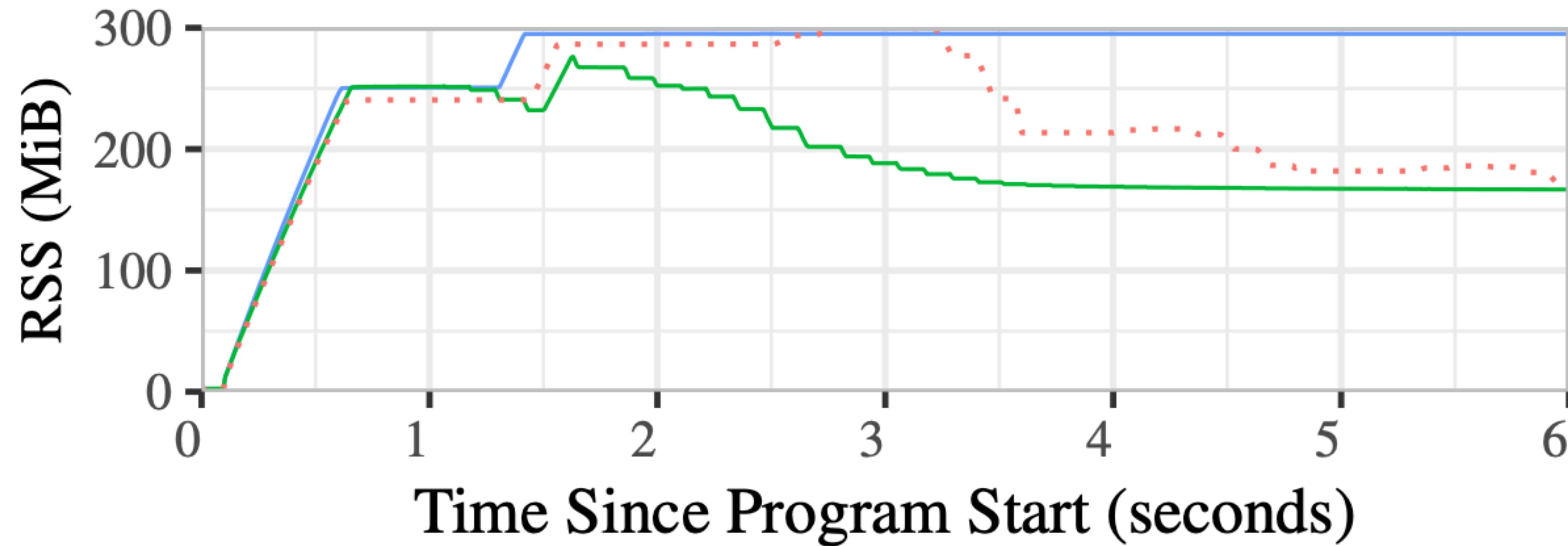
< 1% performance overhead

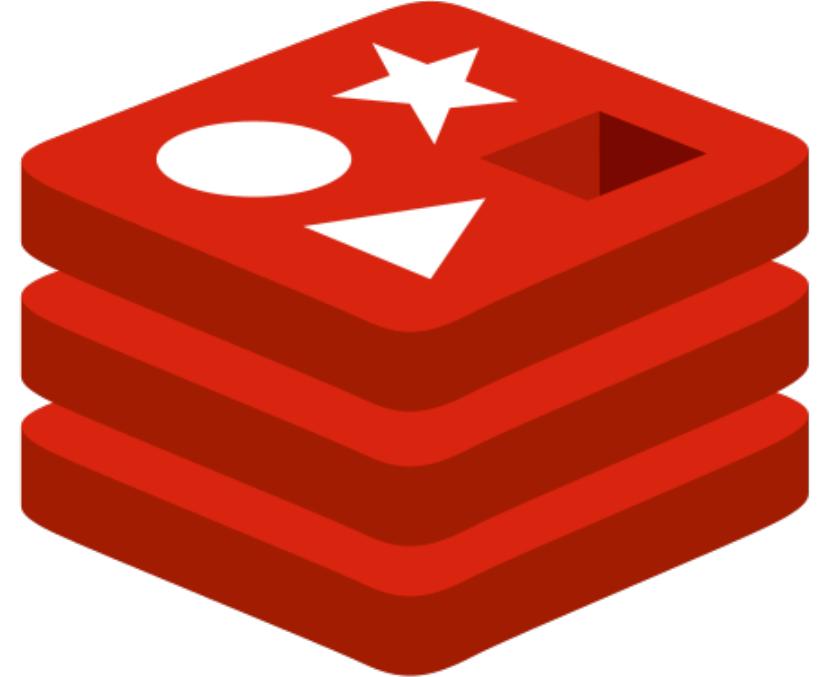




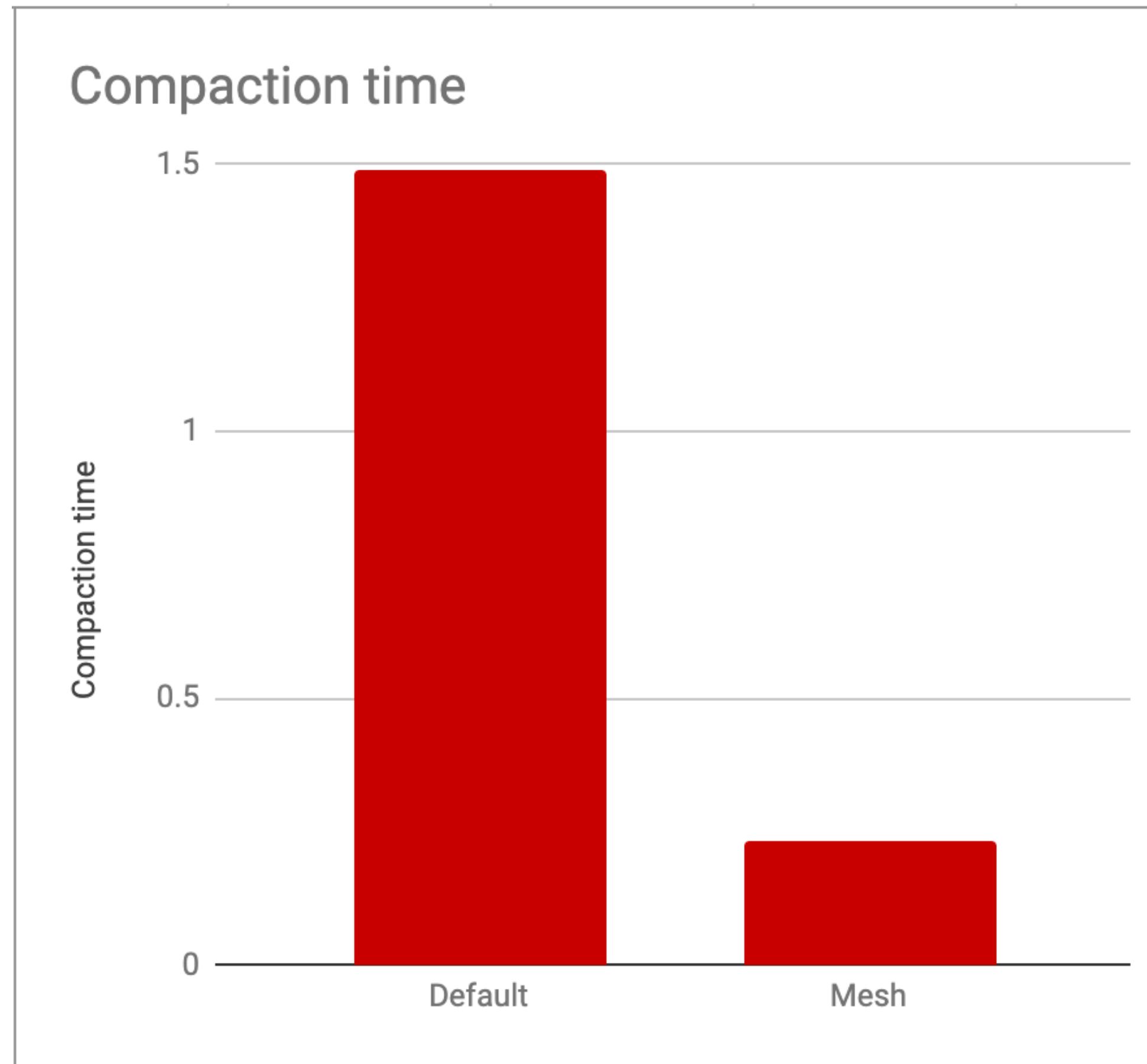
redis

- jemalloc + activatedefrag
- Mesh
- no compaction



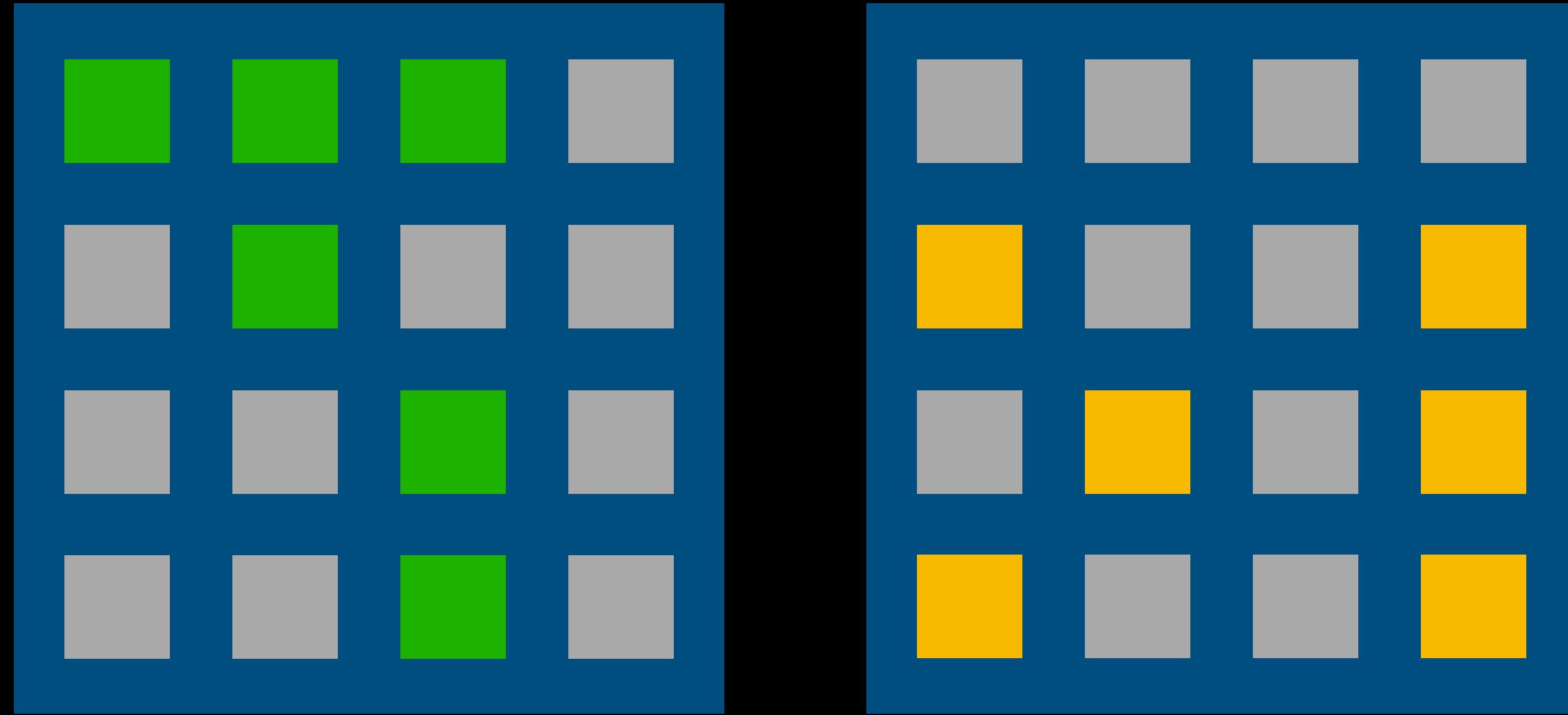


redis



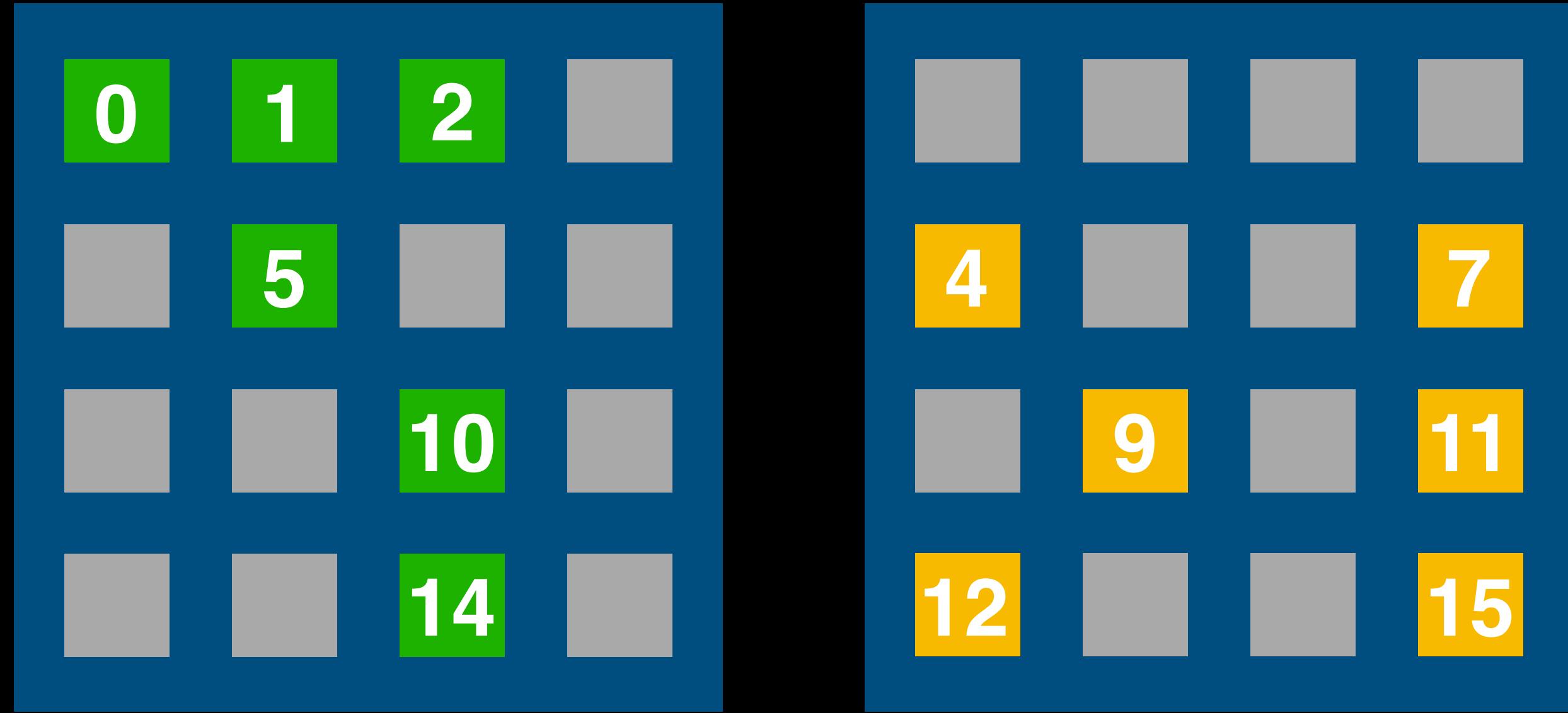
5x reduction in
time spent
compacting

Meshing: compaction without
(virtual!) relocation



Pages **meshable** when they:

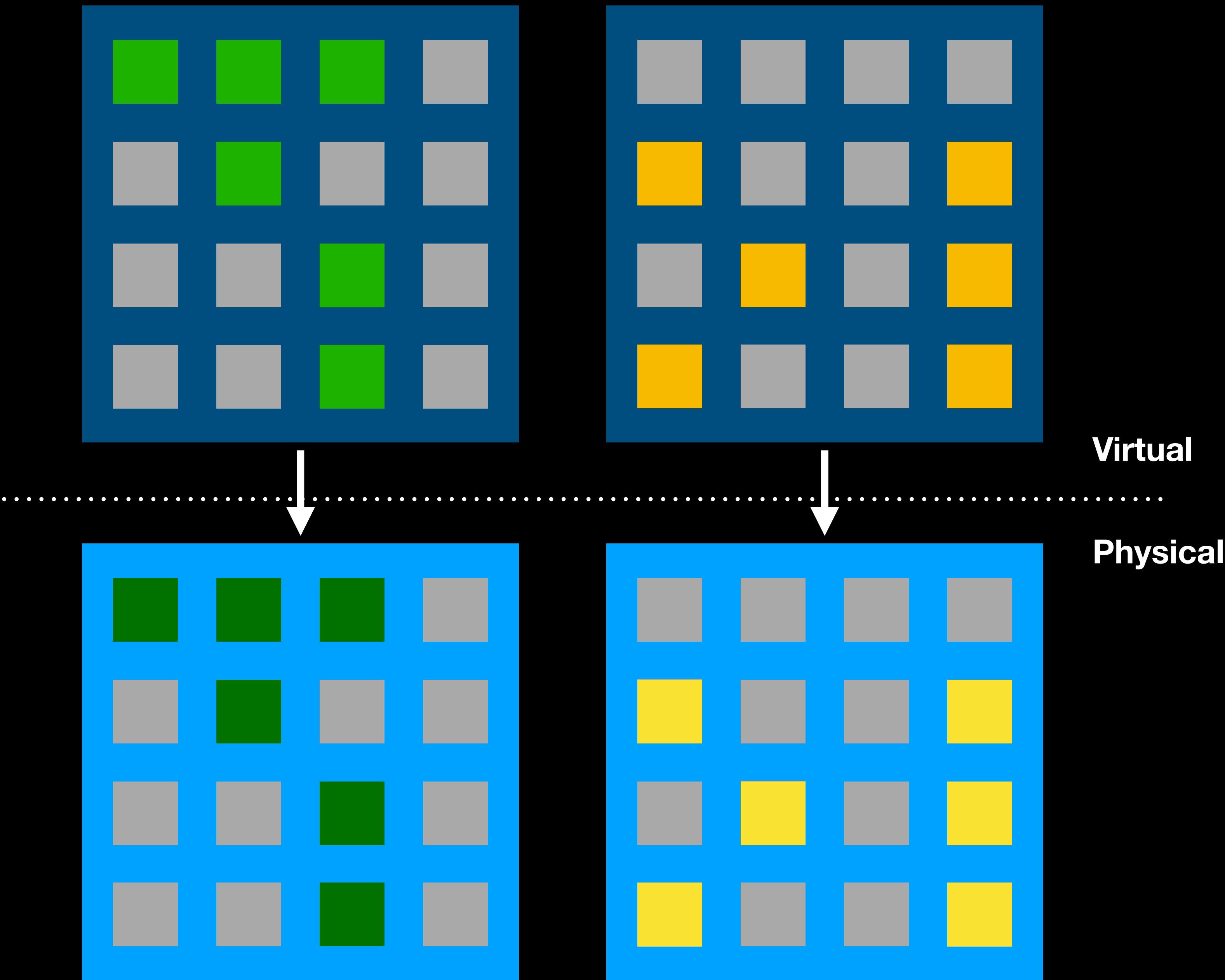
- Hold objects of same size class



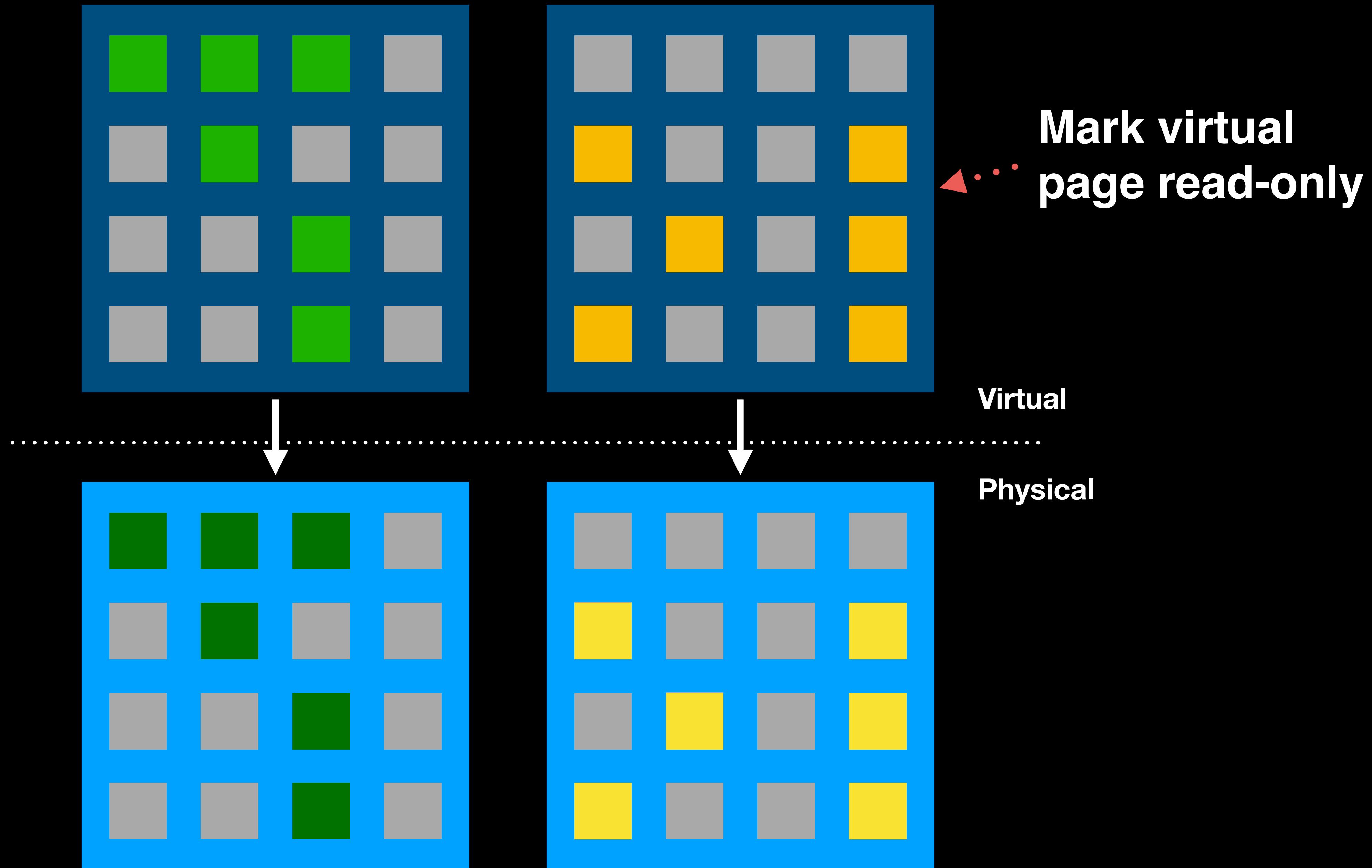
Pages **meshable** when they:

- Hold objects of same size class
- Have non-overlapping object offsets

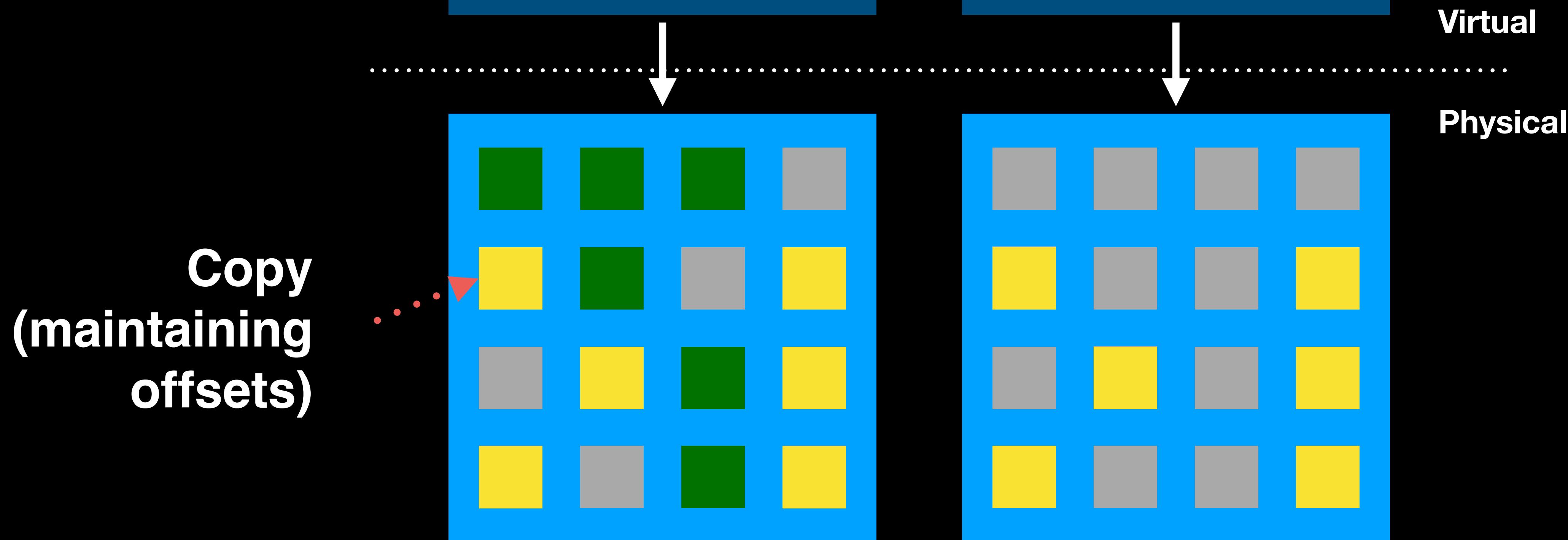
Meshing



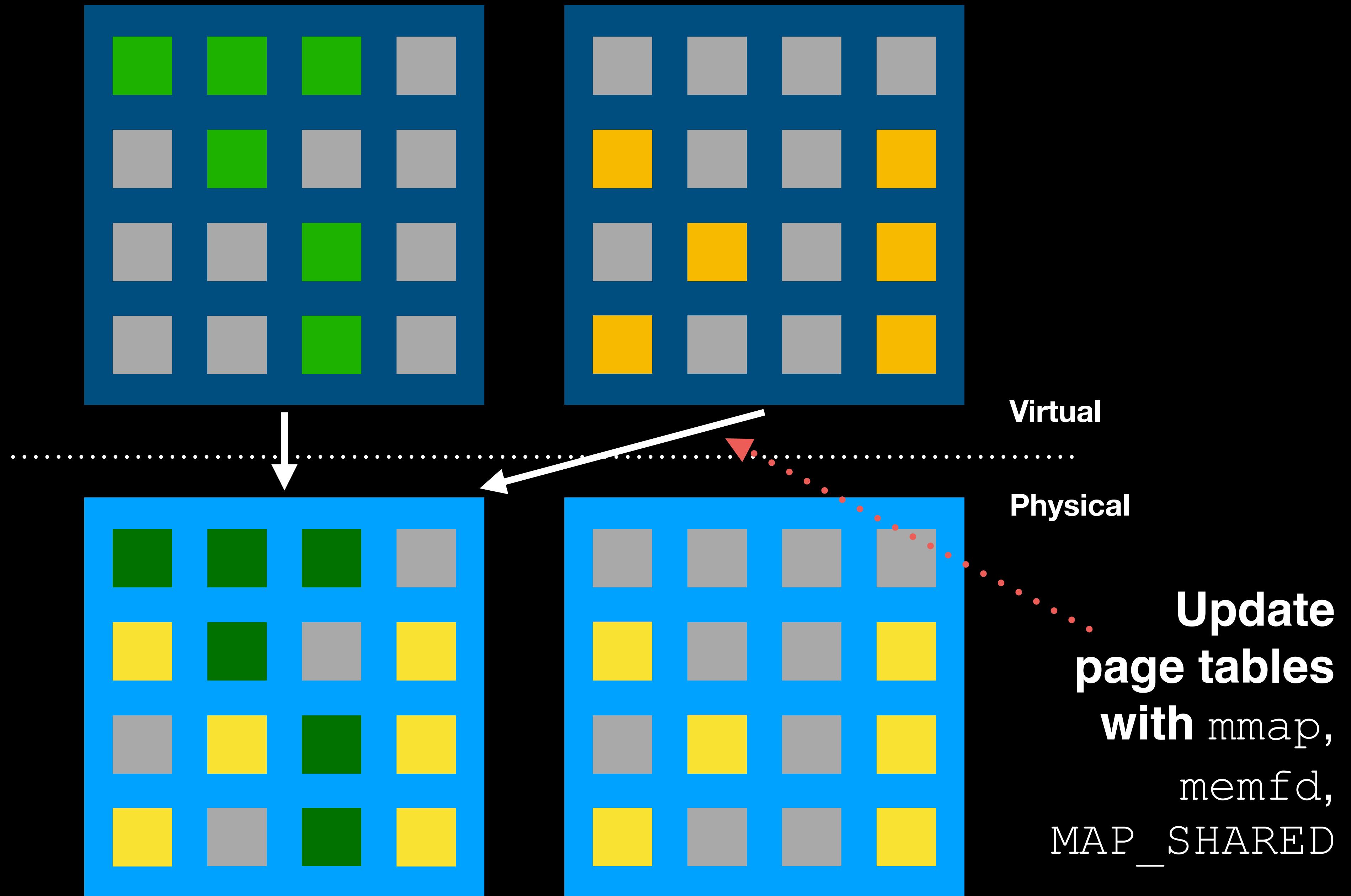
Meshing



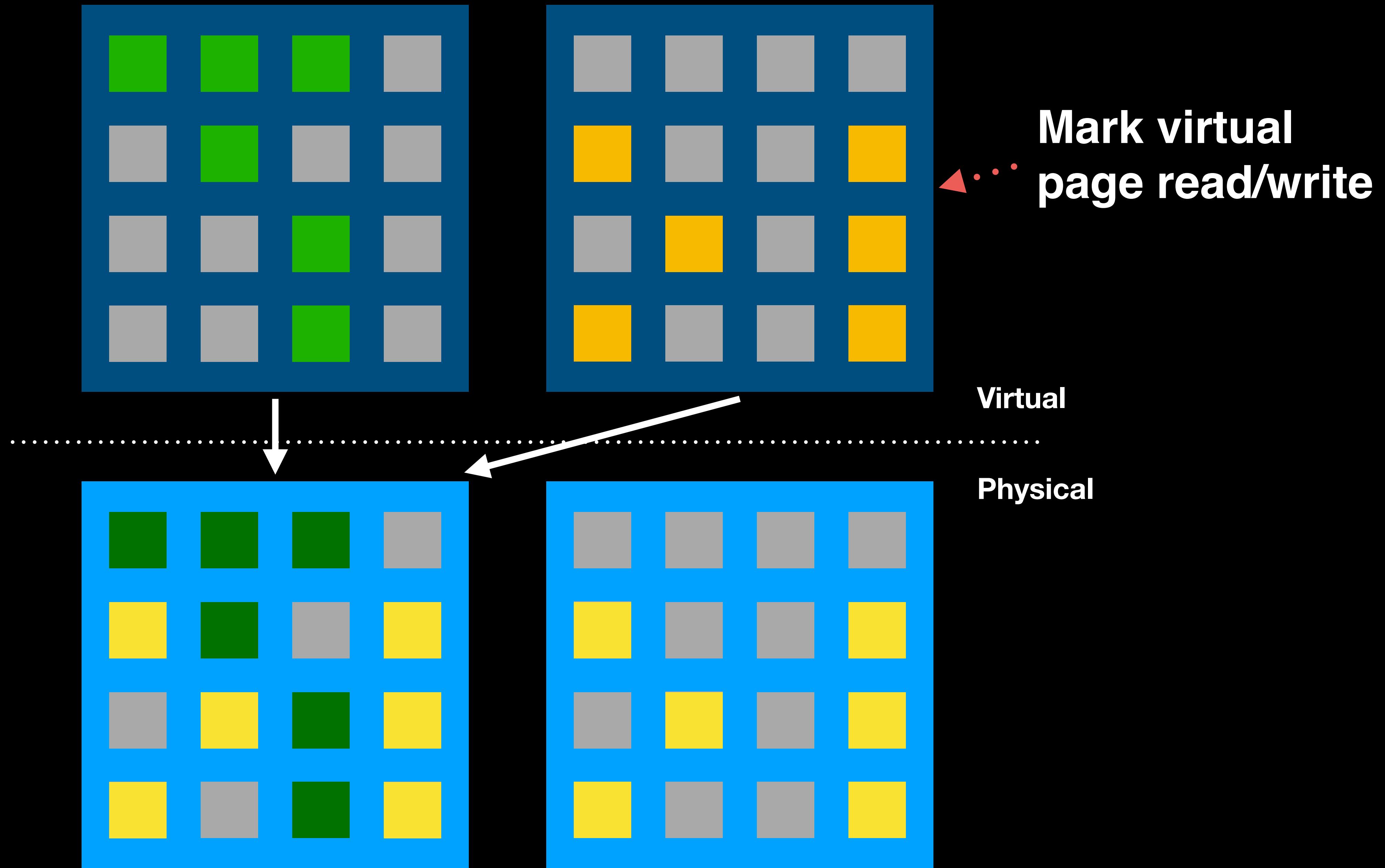
Meshing



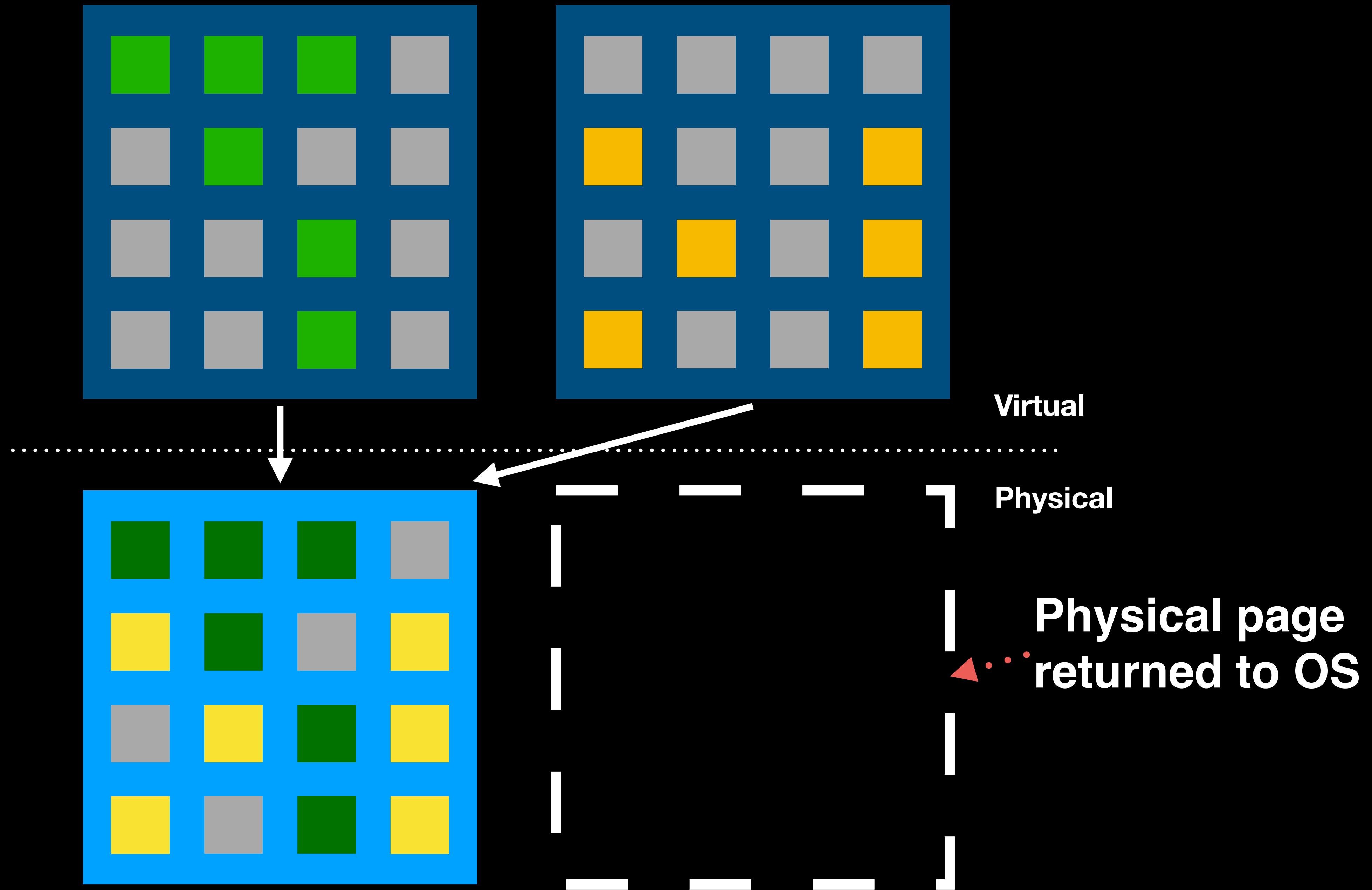
Meshing



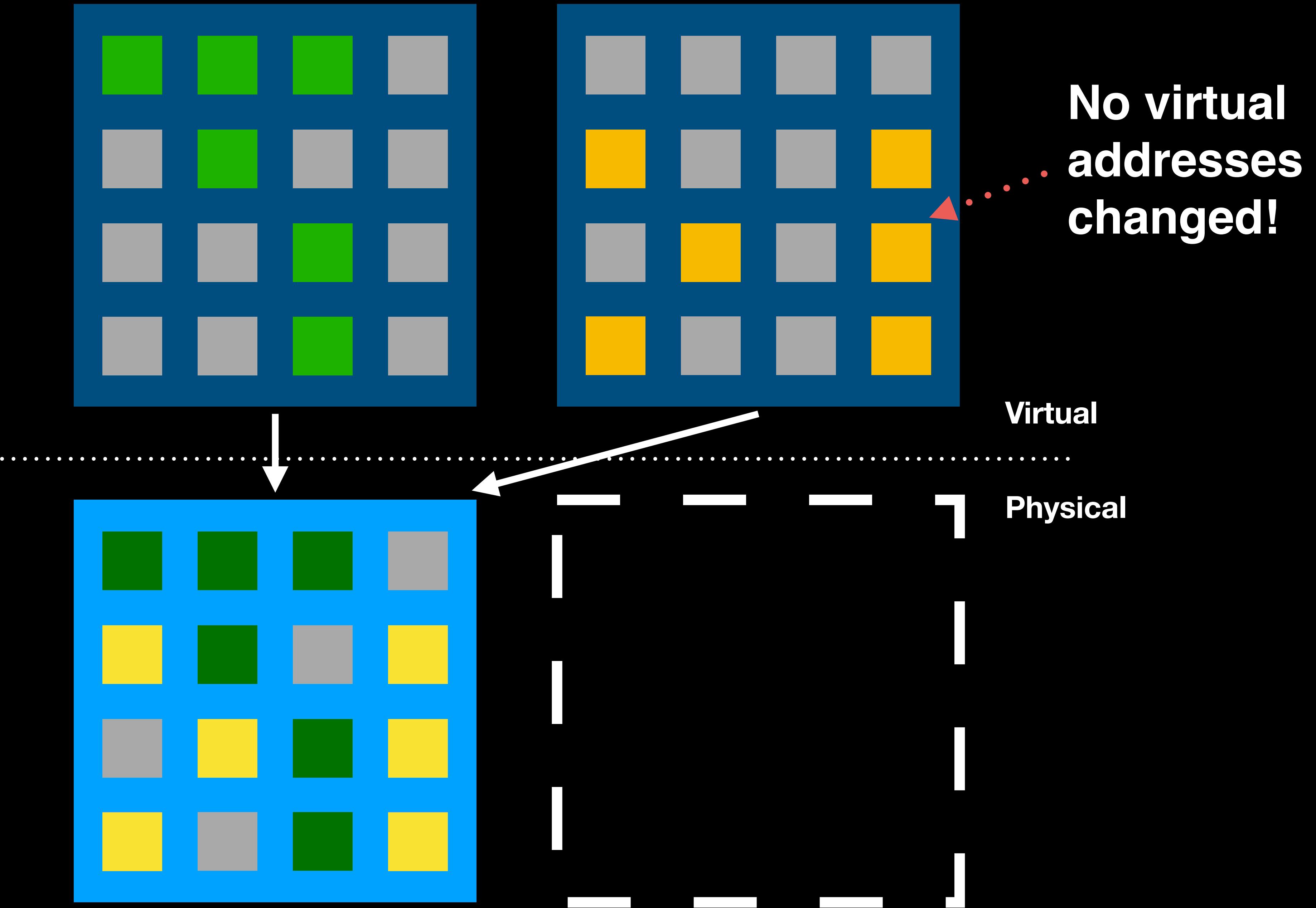
Meshing



Meshing



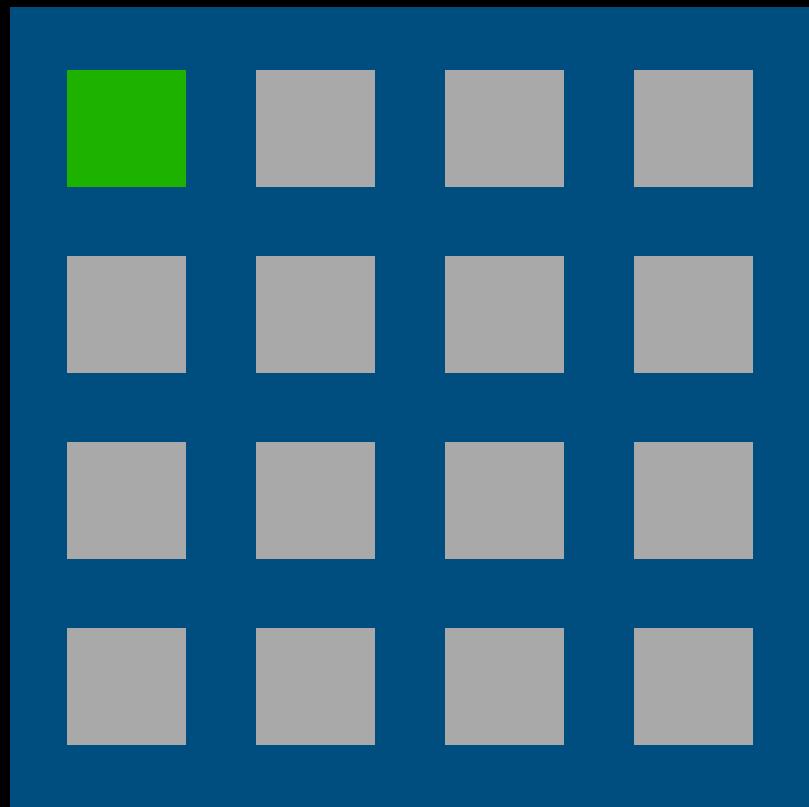
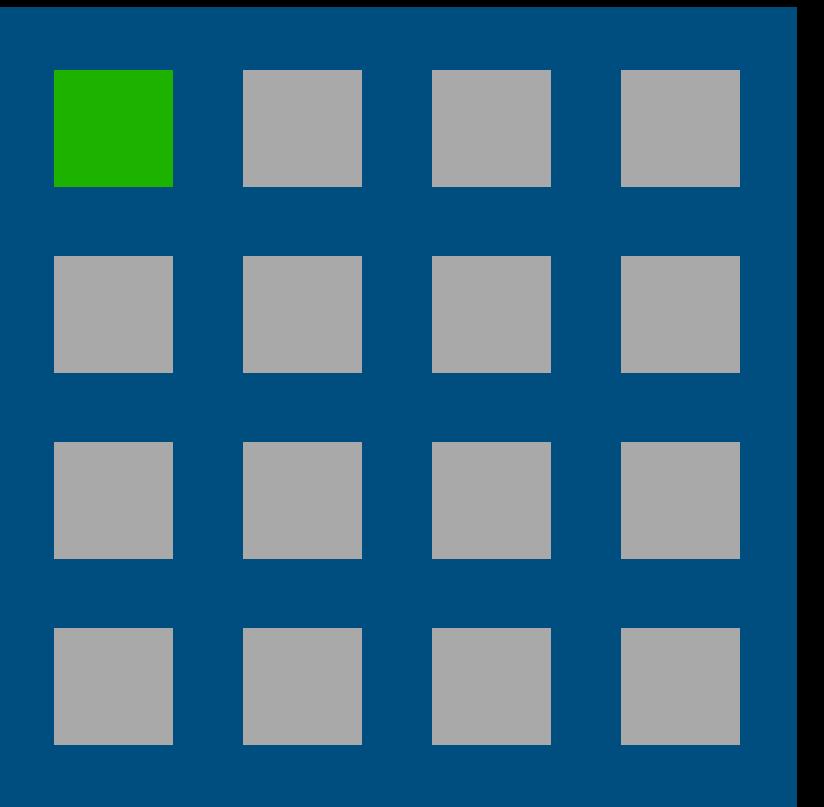
Meshing



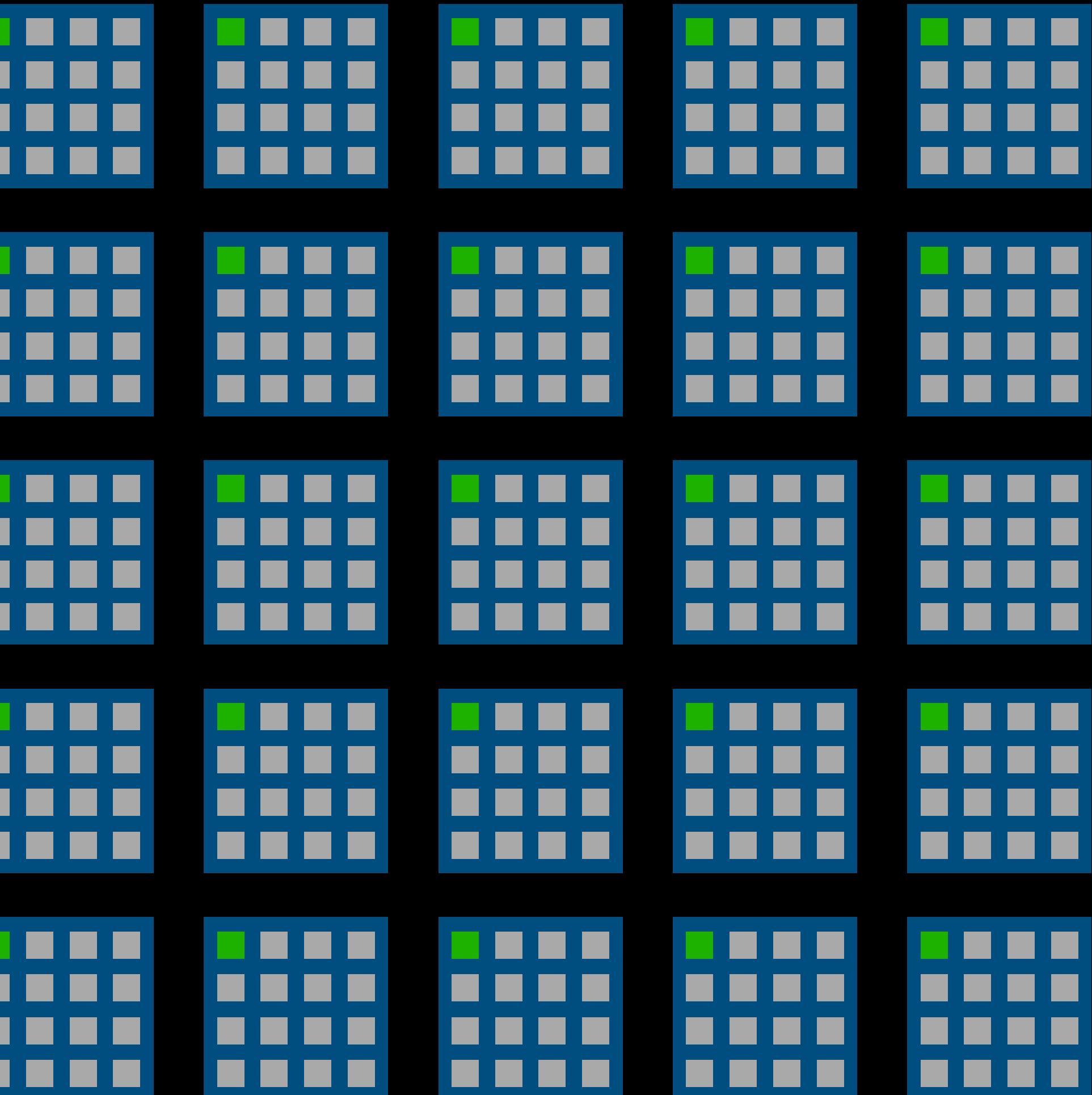
Worst Case:

Worst Case:

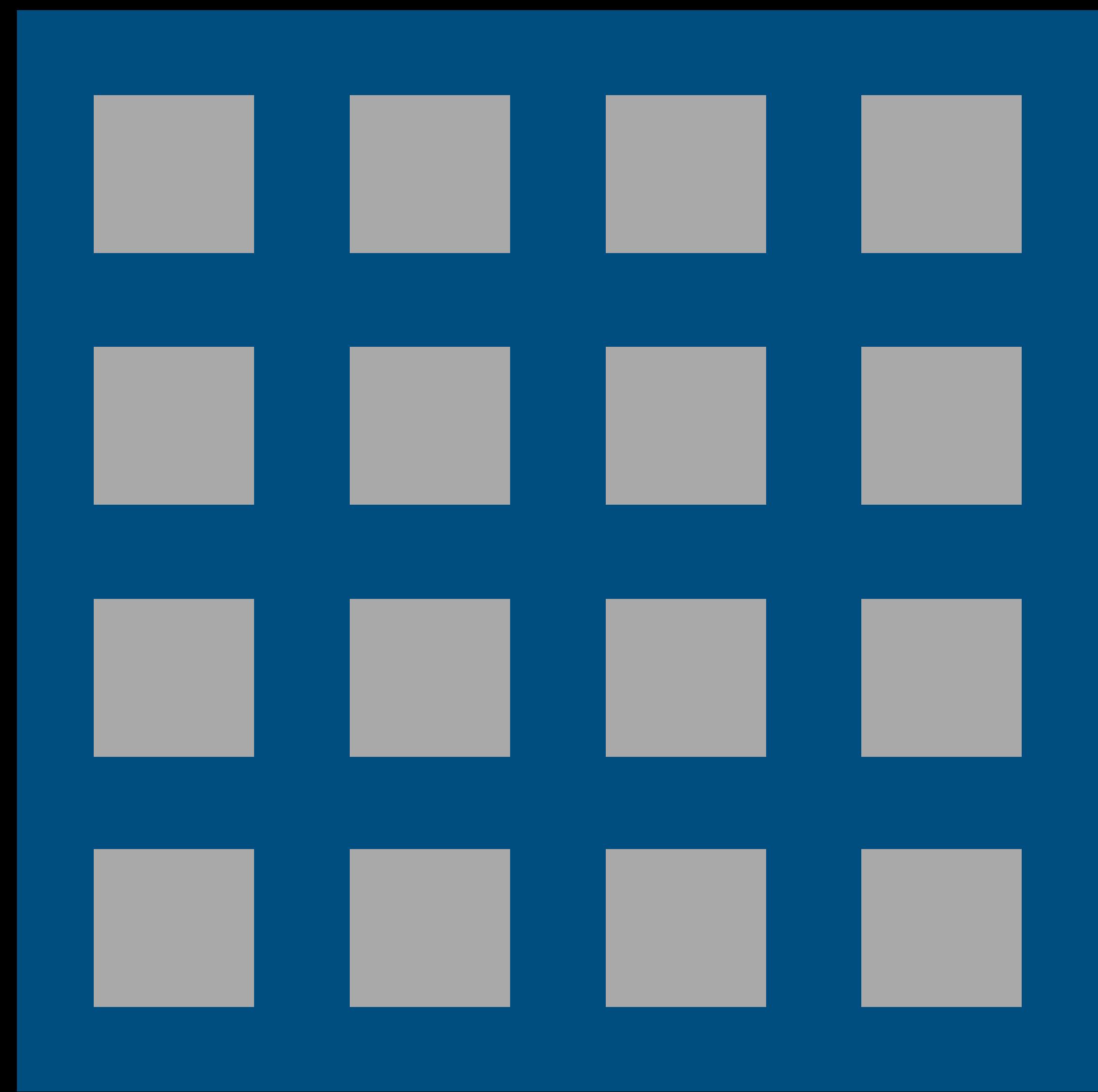
low occupancy,
non-meshable
pages



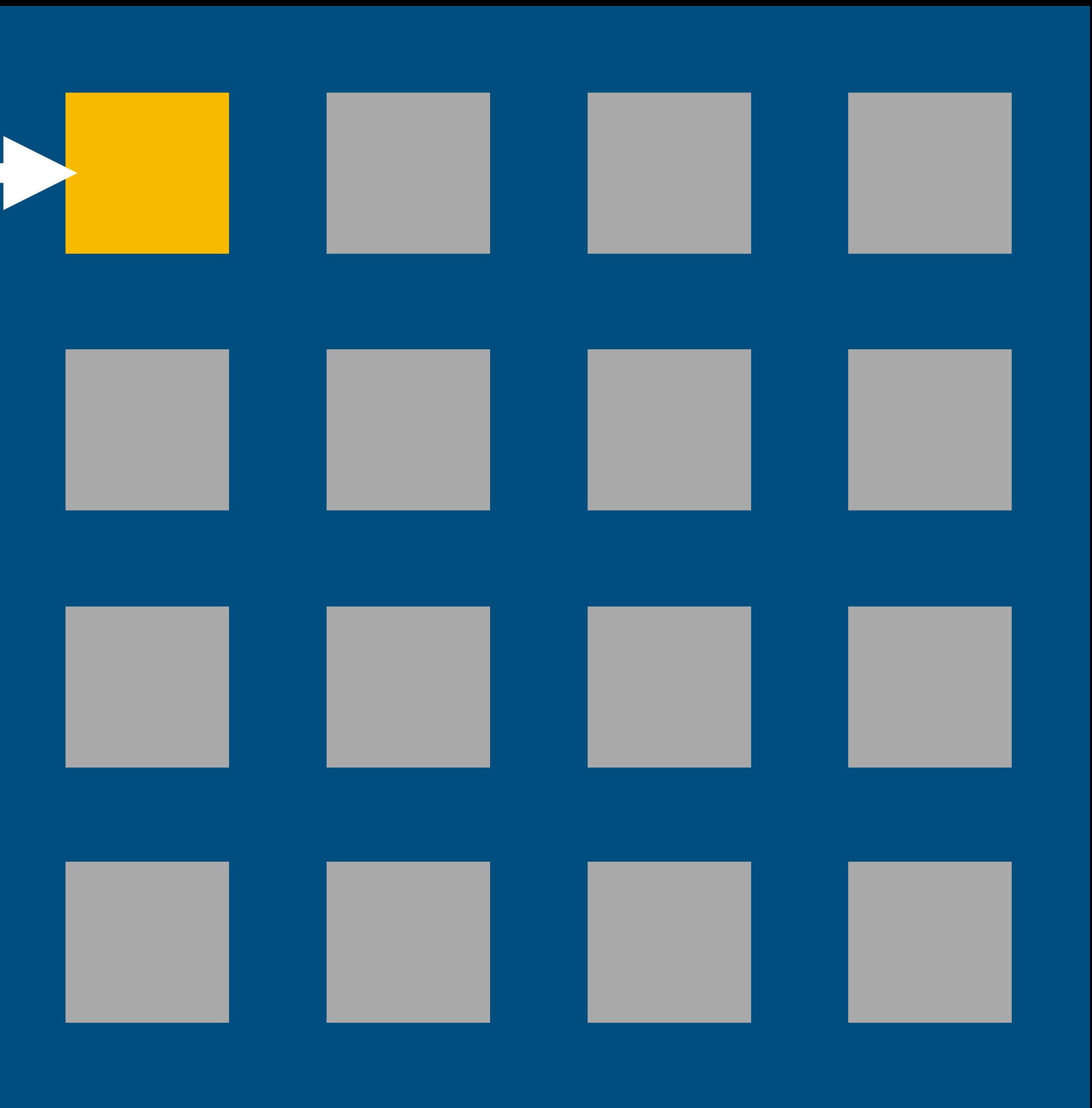
Worst Case:
many
low occupancy,
non-meshable
pages



Standard allocators



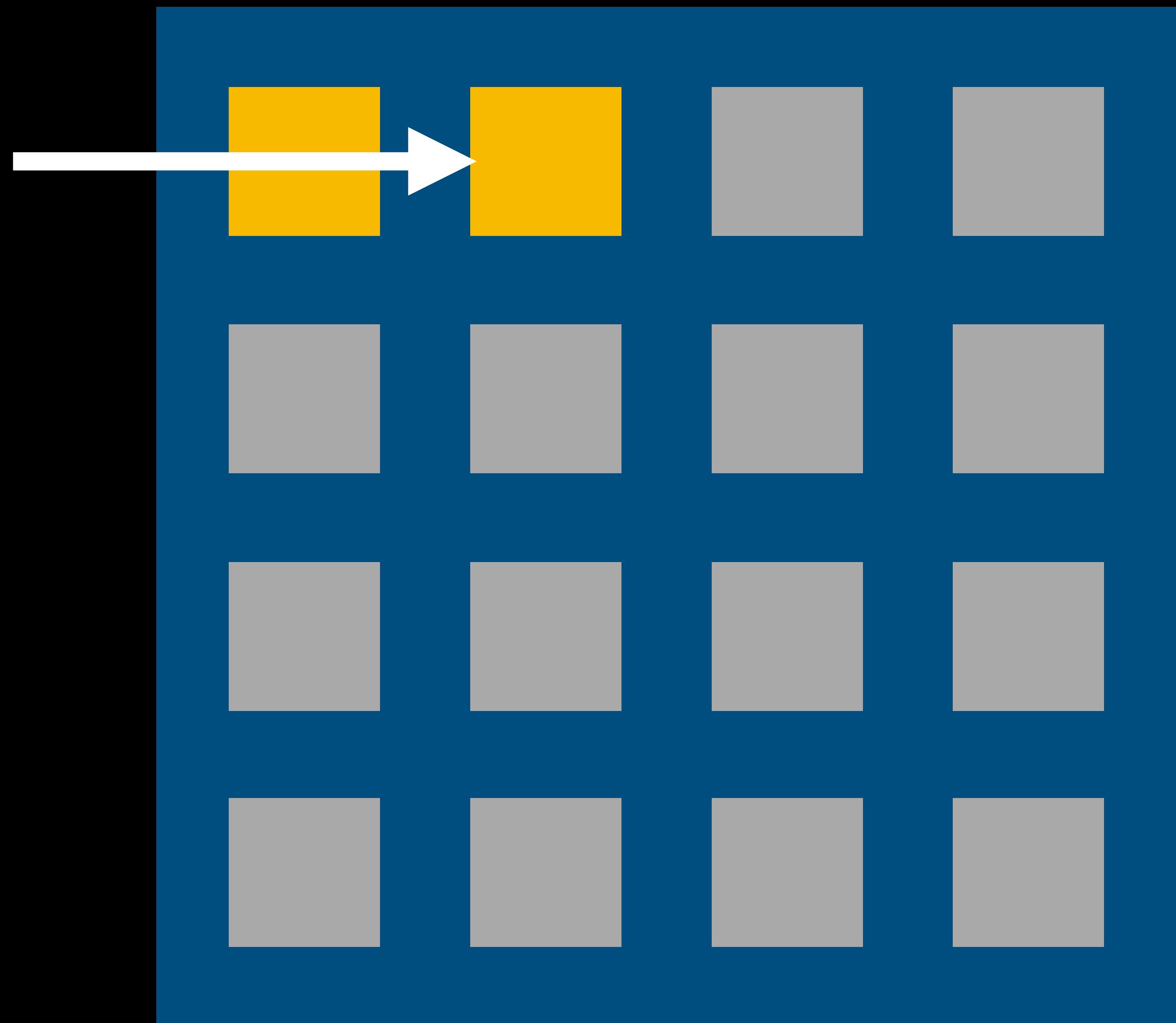
`malloc(256)`



**Standard
allocators**

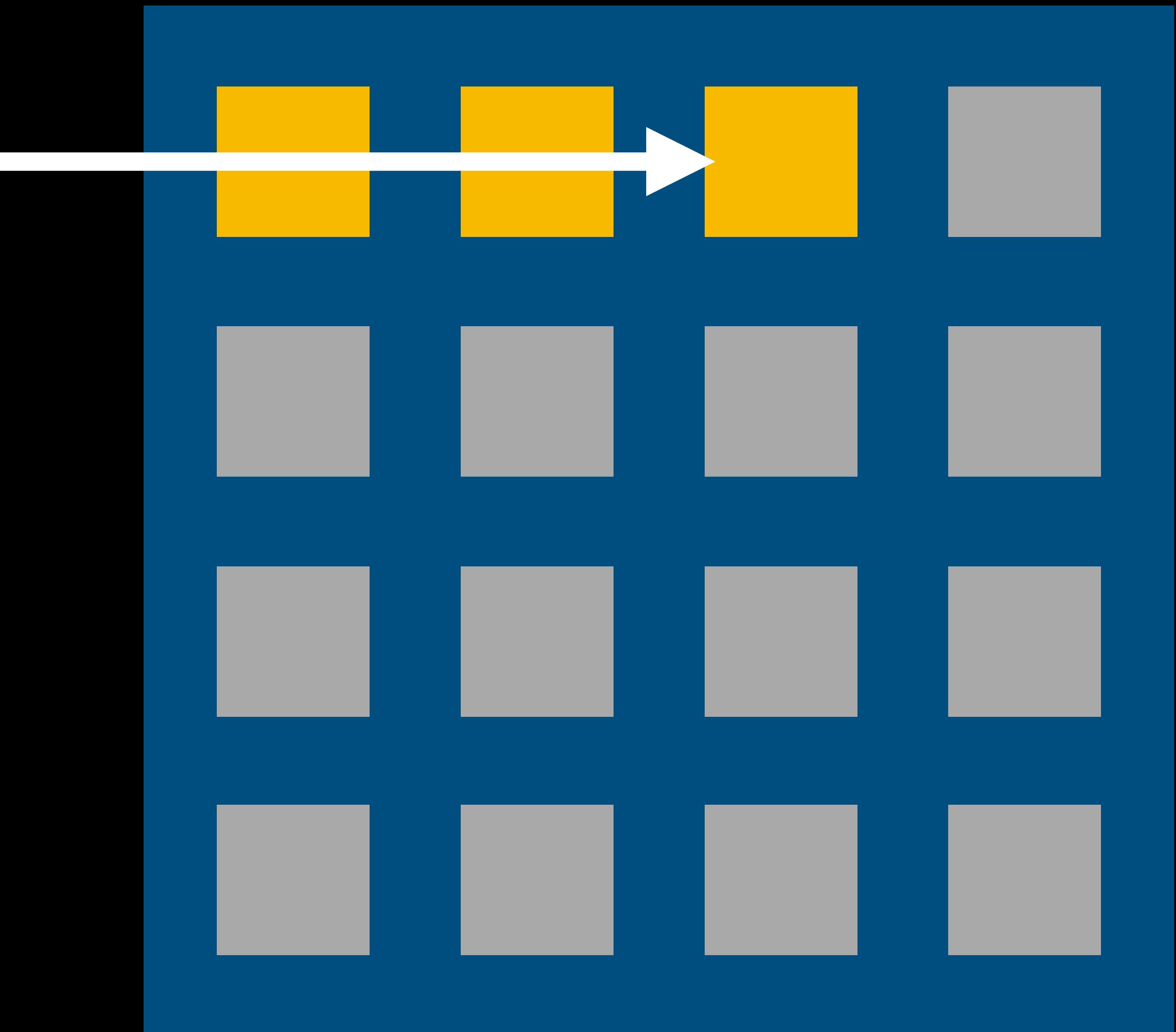
malloc(256)

**Standard
allocators**



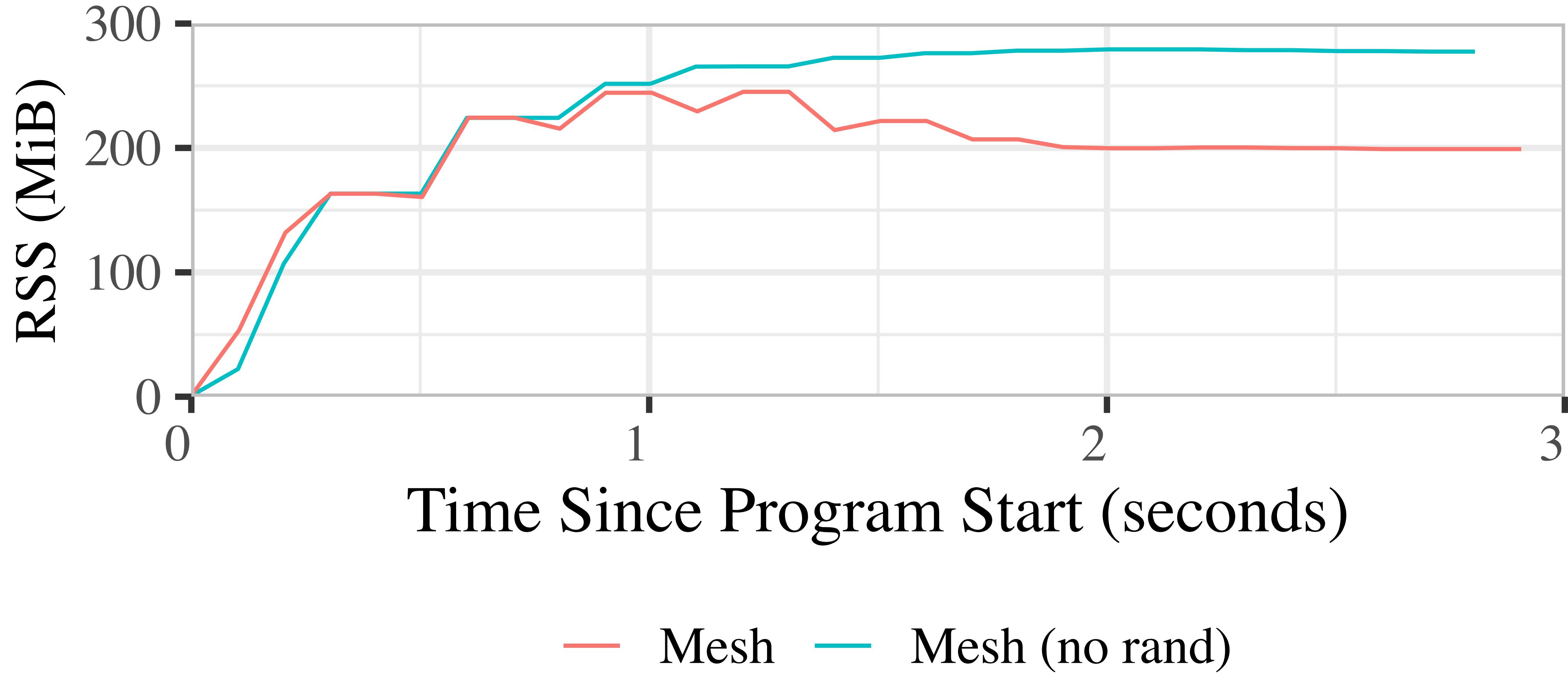
malloc(256)

**Standard
allocators**



Mesh uses **randomization** to ensure live objects are uniformly distributed

Regular allocation patterns are real



How to **randomize** allocation?

Random probing:

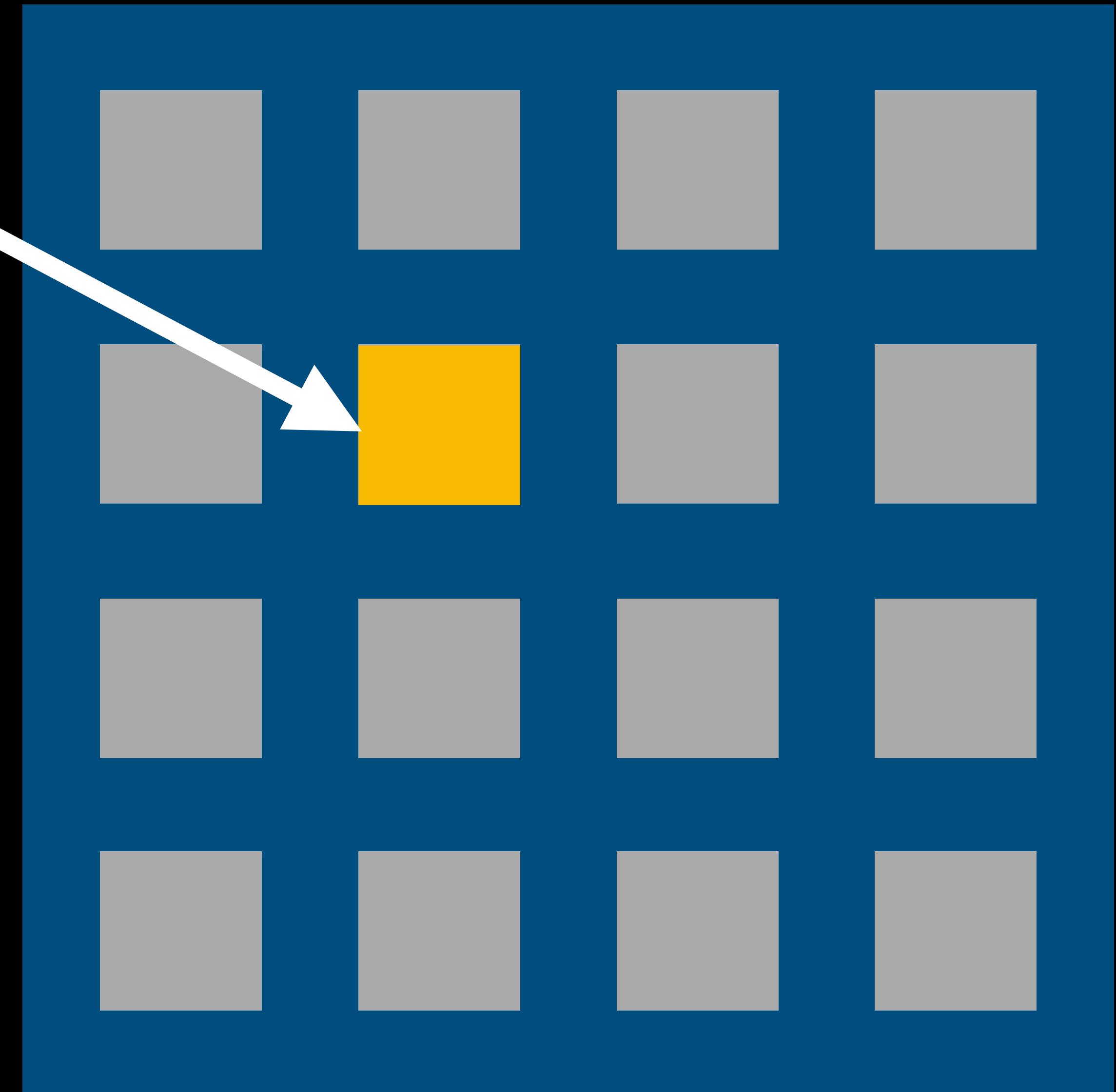
```
while true:  
    if rand_off().is_free:  
        return rand_off
```



malloc(256)

Random probing:

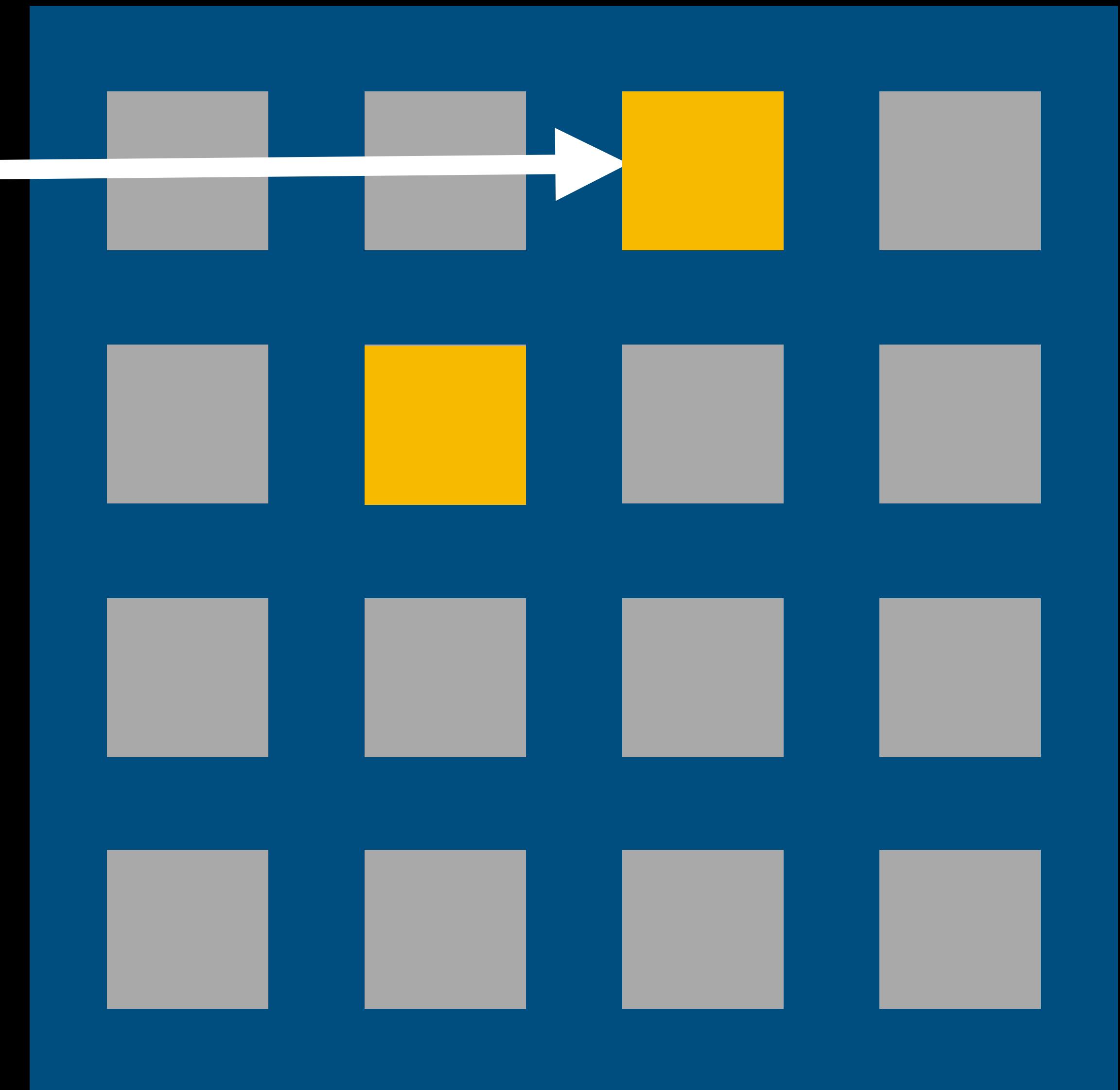
```
while true:  
    if rand_off().is_free:  
        return rand_off
```



malloc(256)

Random probing:

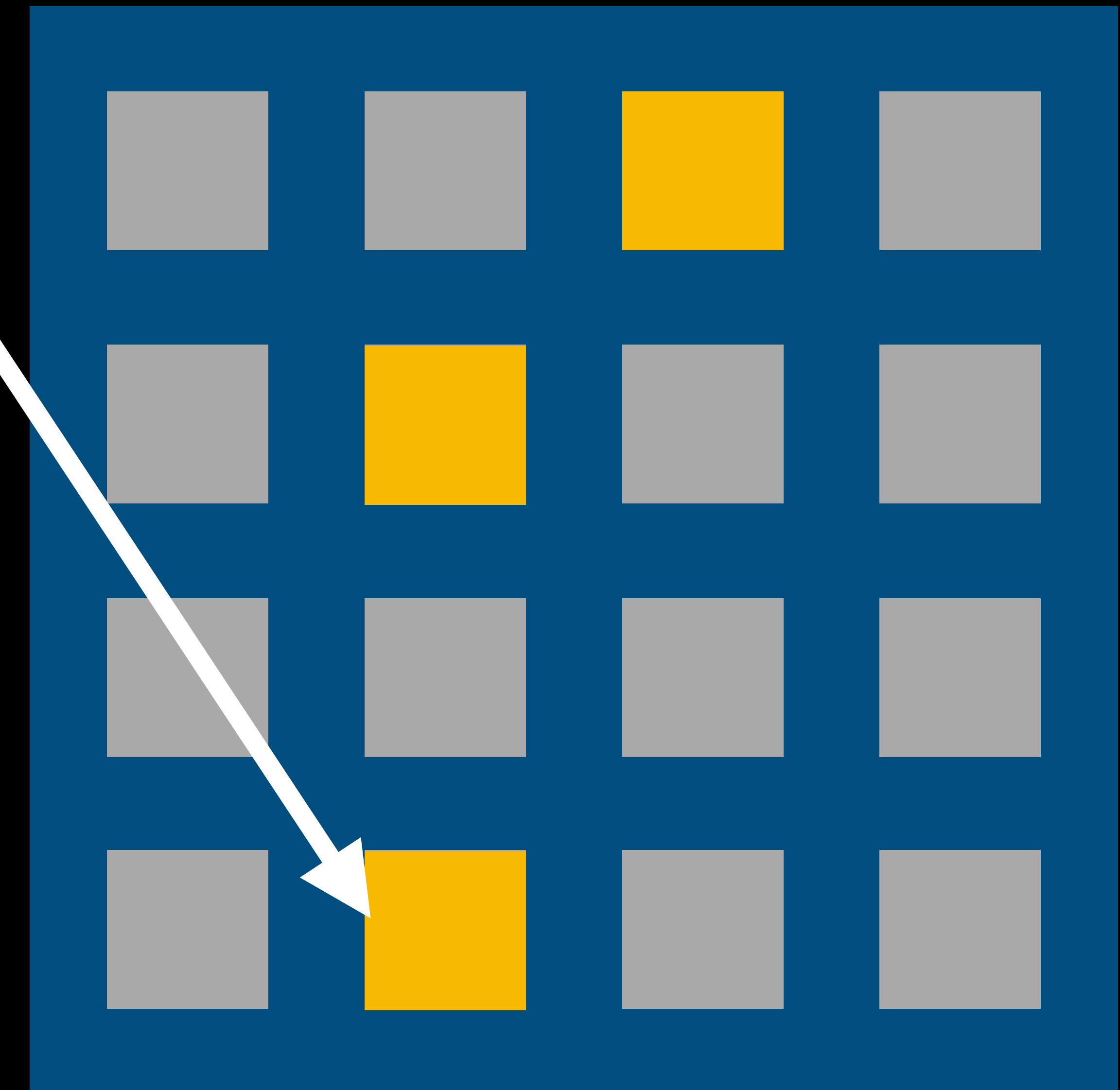
```
while true:  
    if rand_off().is_free:  
        return rand_off
```



```
malloc(256)
```

Random probing:

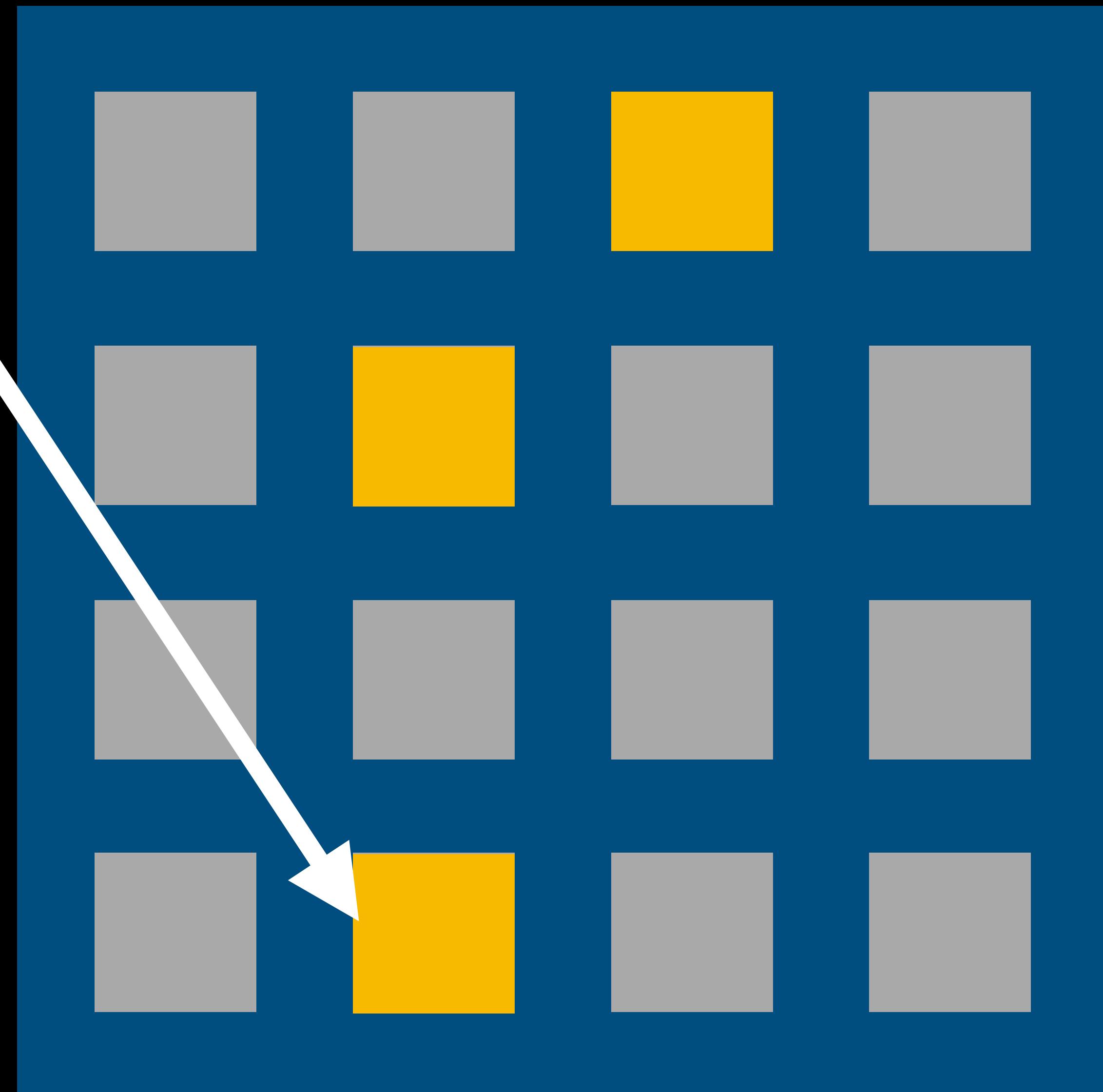
```
while true:  
    if rand_off().is_free:  
        return rand_off
```



malloc(256)

Random probing:

```
while true:  
    if rand_off().is_free:  
        return rand_off
```



(DieHard [Berger & Zorn 2006])

Random probing fast in
expectation *iff page*
occupancy is low

Random probing fast in
expectation *iff page*
occupancy is low

**but this is at odds with
minimizing heap size!**

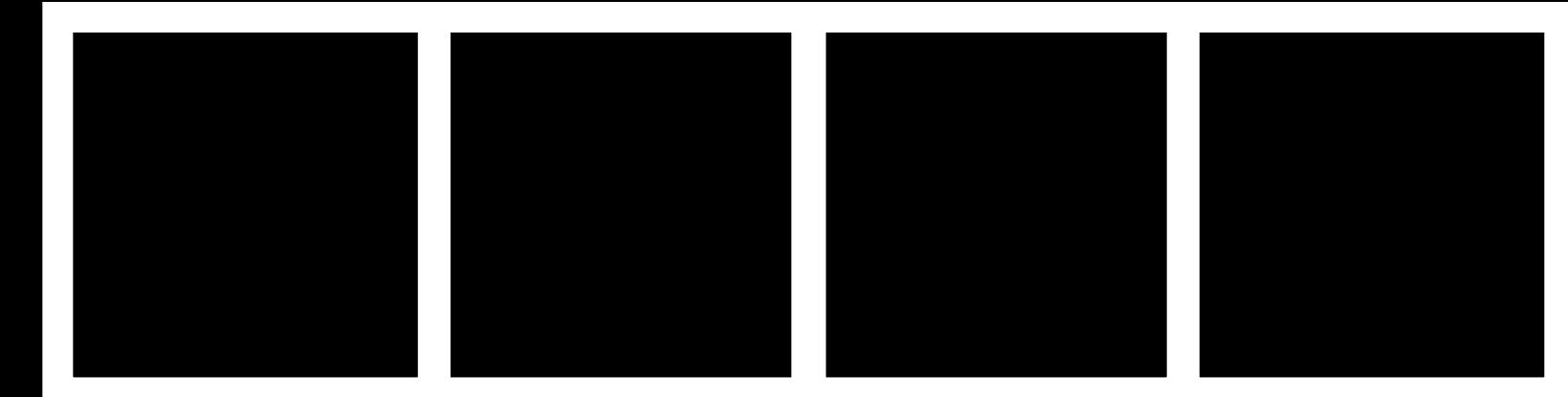
Shuffle Vector:
Fast randomized
allocation + full page utilization

Shuffle Vector: Fast randomized allocation

Page

0 1 2 3

load



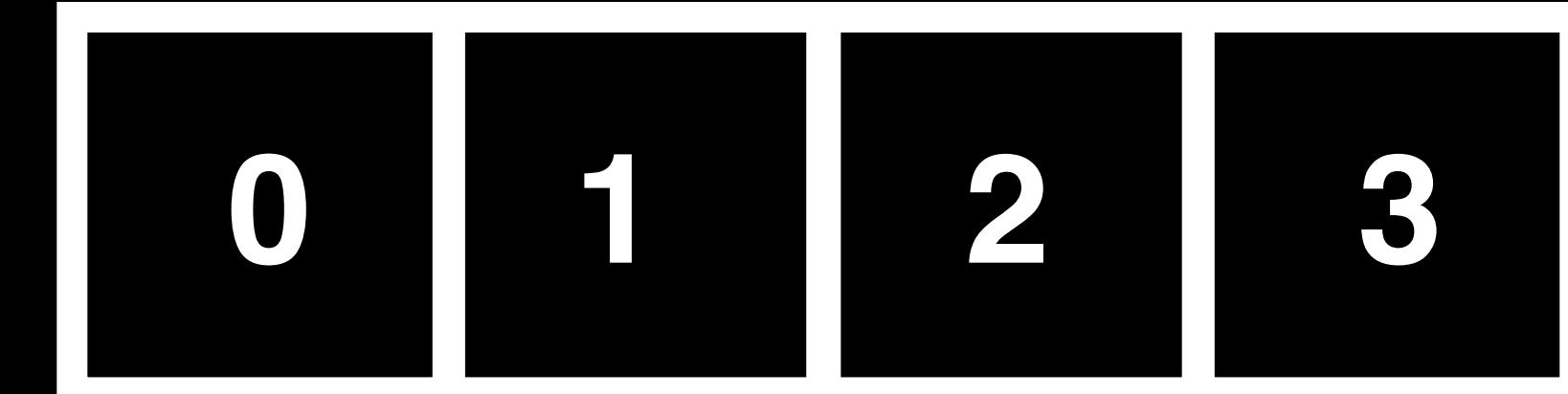
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



load



Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



shuffle (

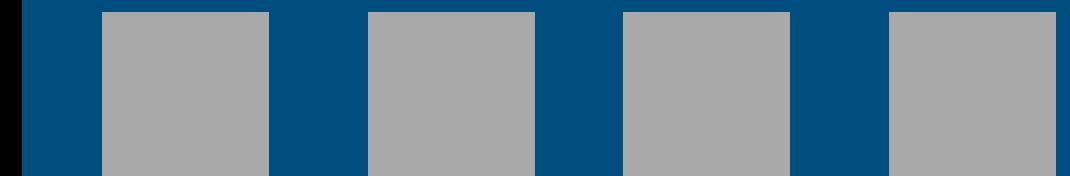
0	1	2	3
---	---	---	---

)

Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



shuffle (

2	3	1	0
---	---	---	---

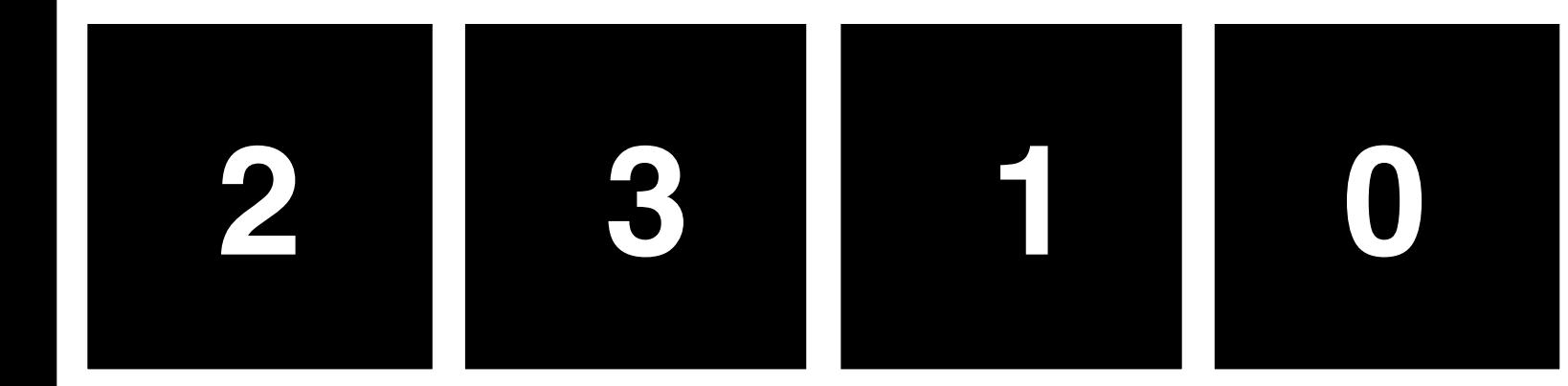
)

Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

`malloc()`

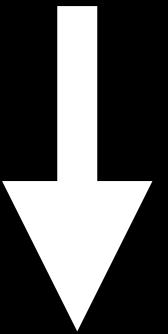
Page



Thread-local shuffle vector

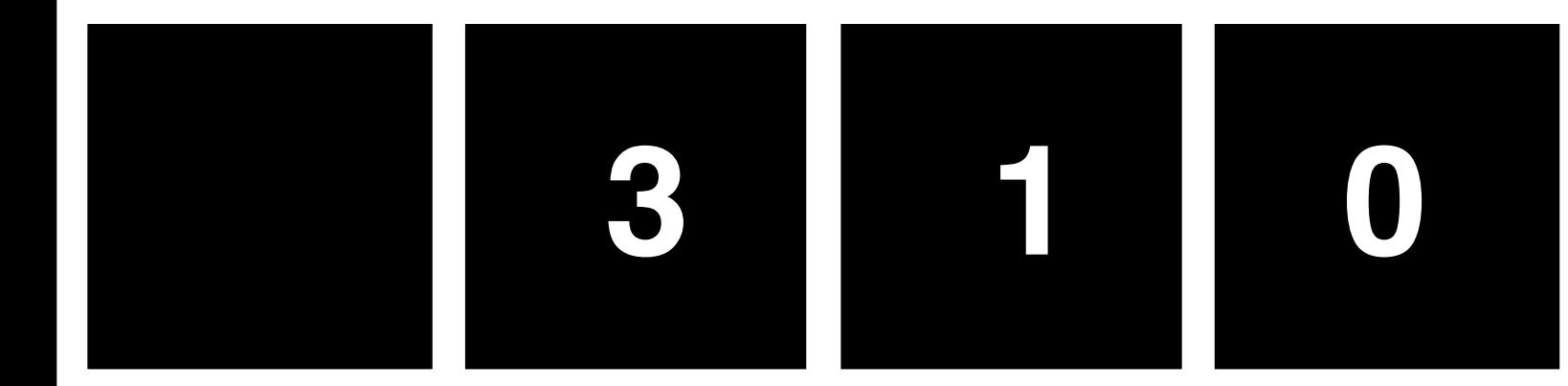
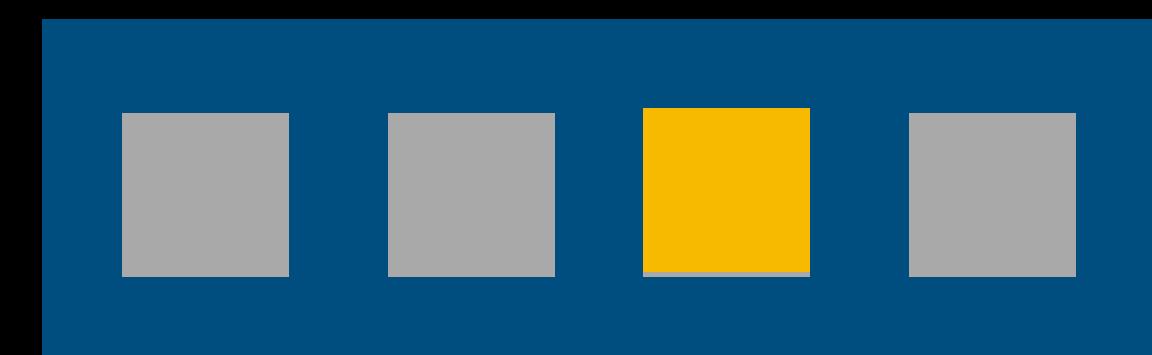
Shuffle Vector: Fast randomized allocation

malloc()



page_start +
2 * object_size

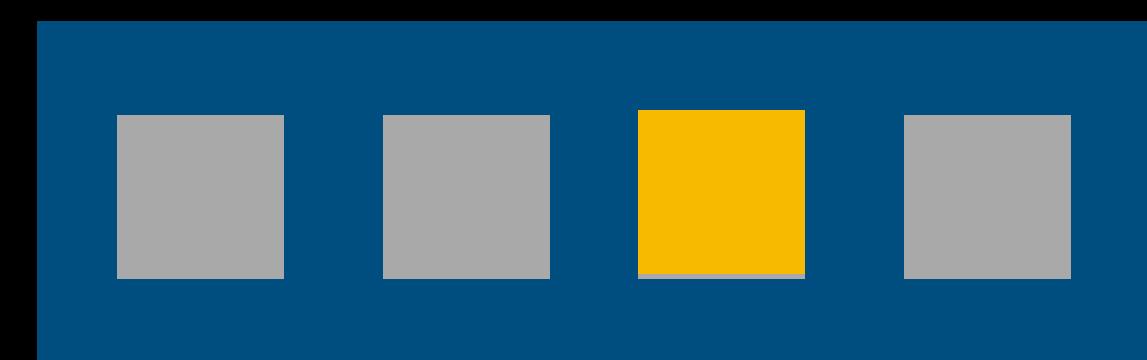
Page



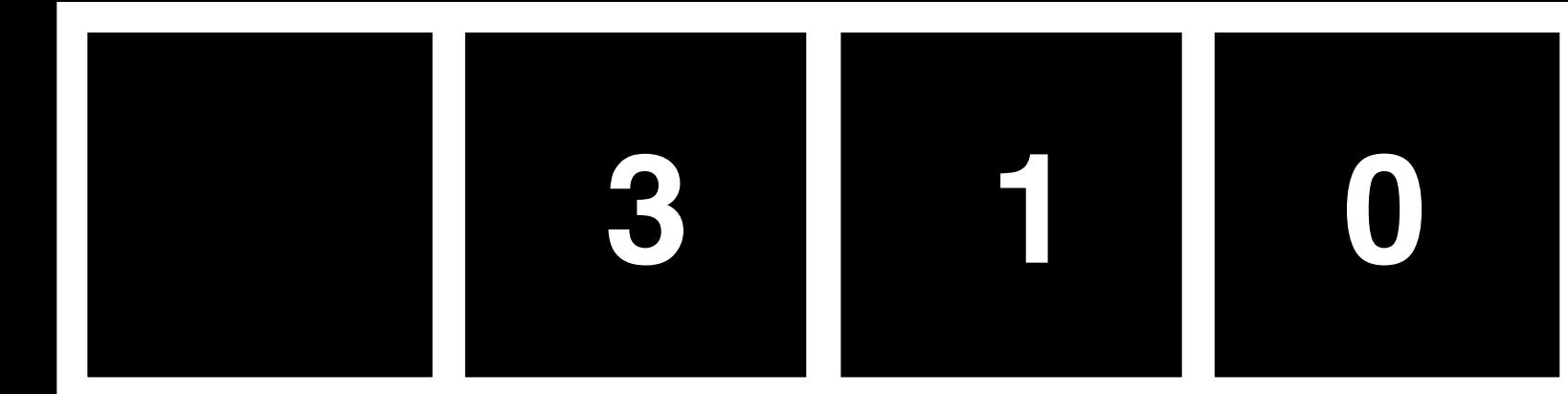
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



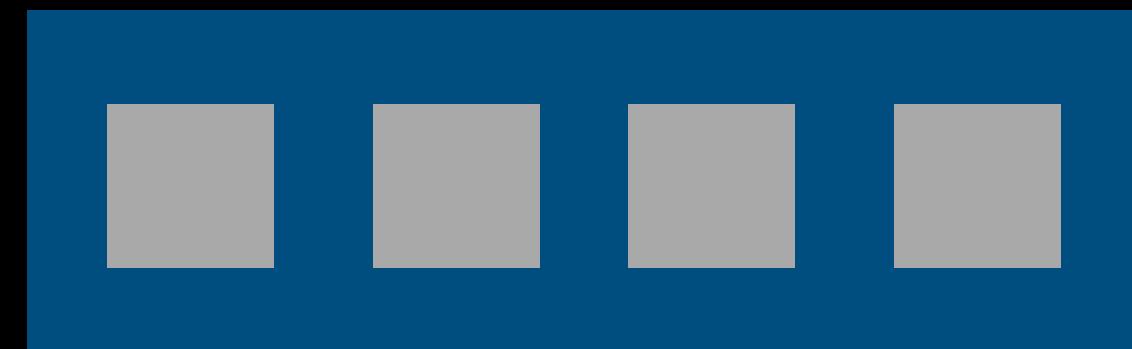
free (2)



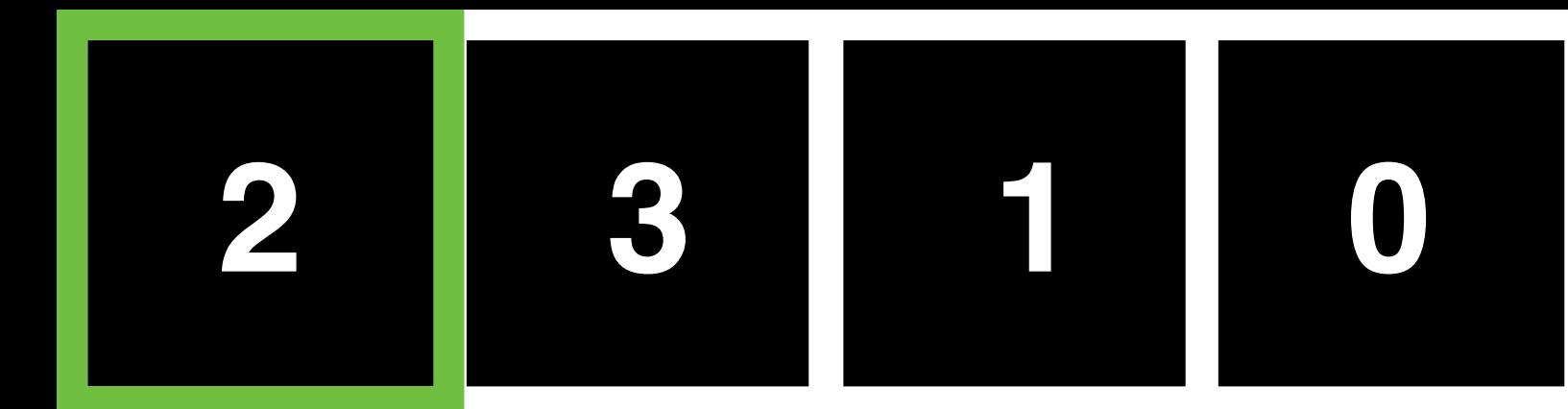
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



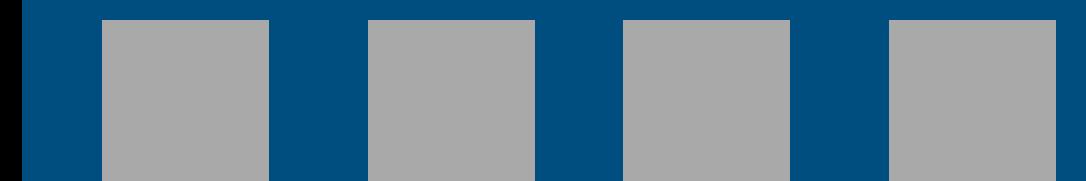
free()



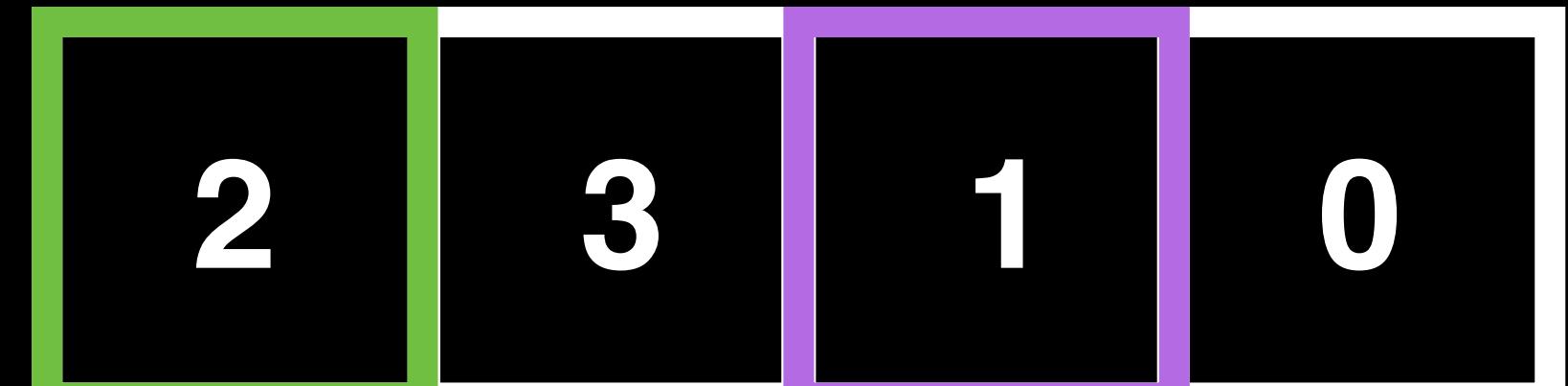
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



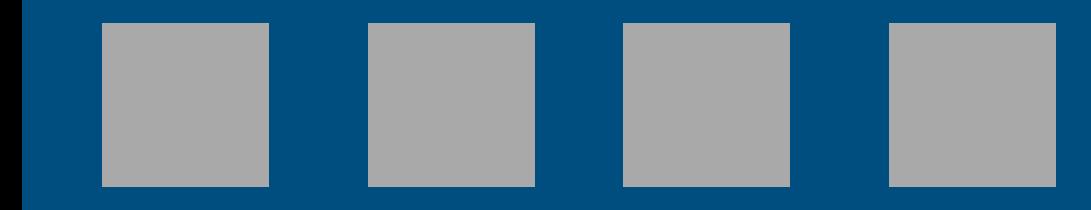
free()

shuffle_one( 0)

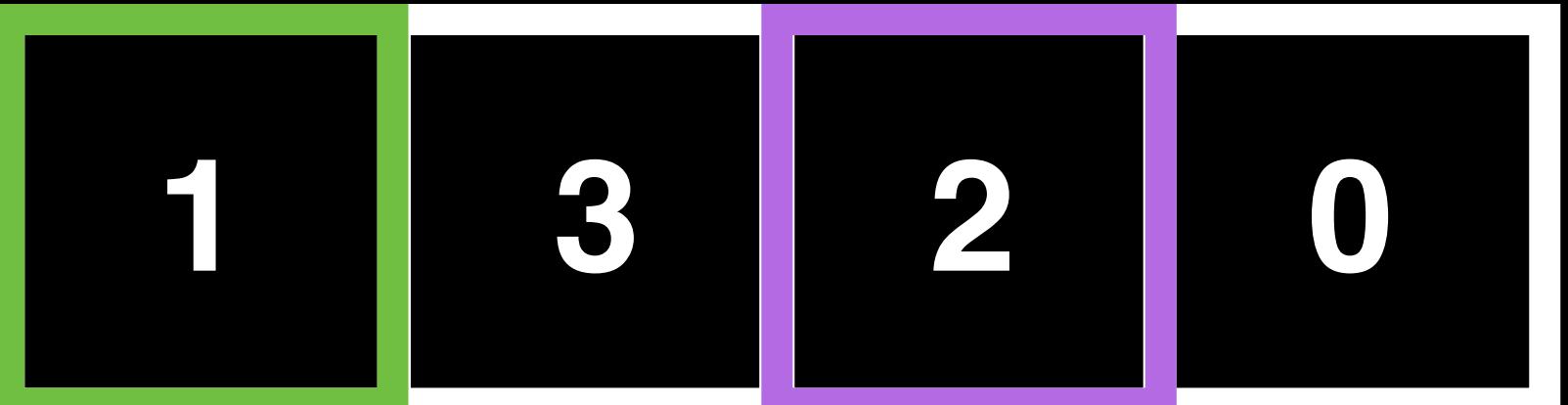
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page



free()

shuffle_one( 0)

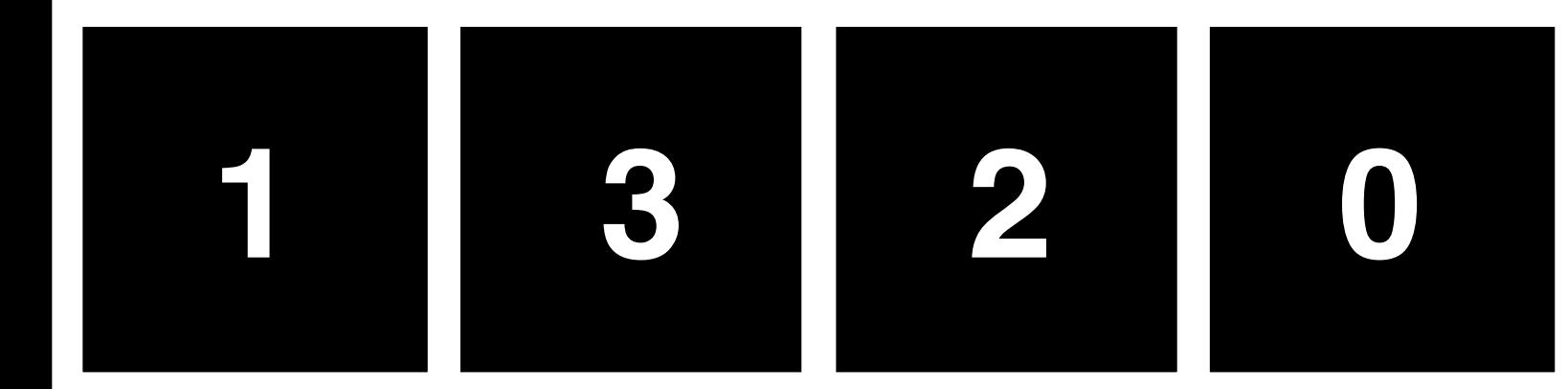
Thread-local shuffle vector

Shuffle Vector: Fast randomized allocation

Page

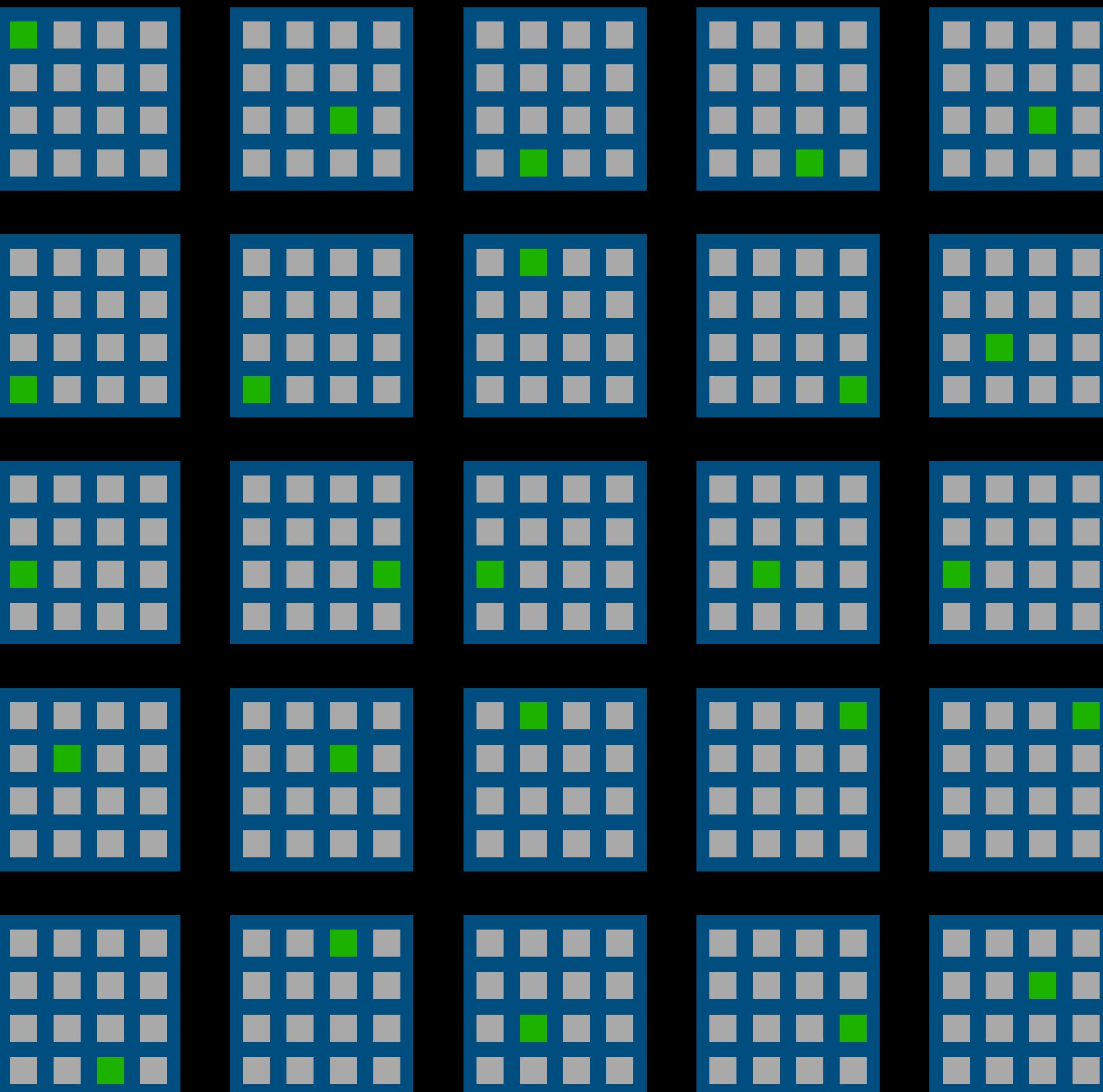


free()



Thread-local shuffle vector

**All
Pages
Meshable**



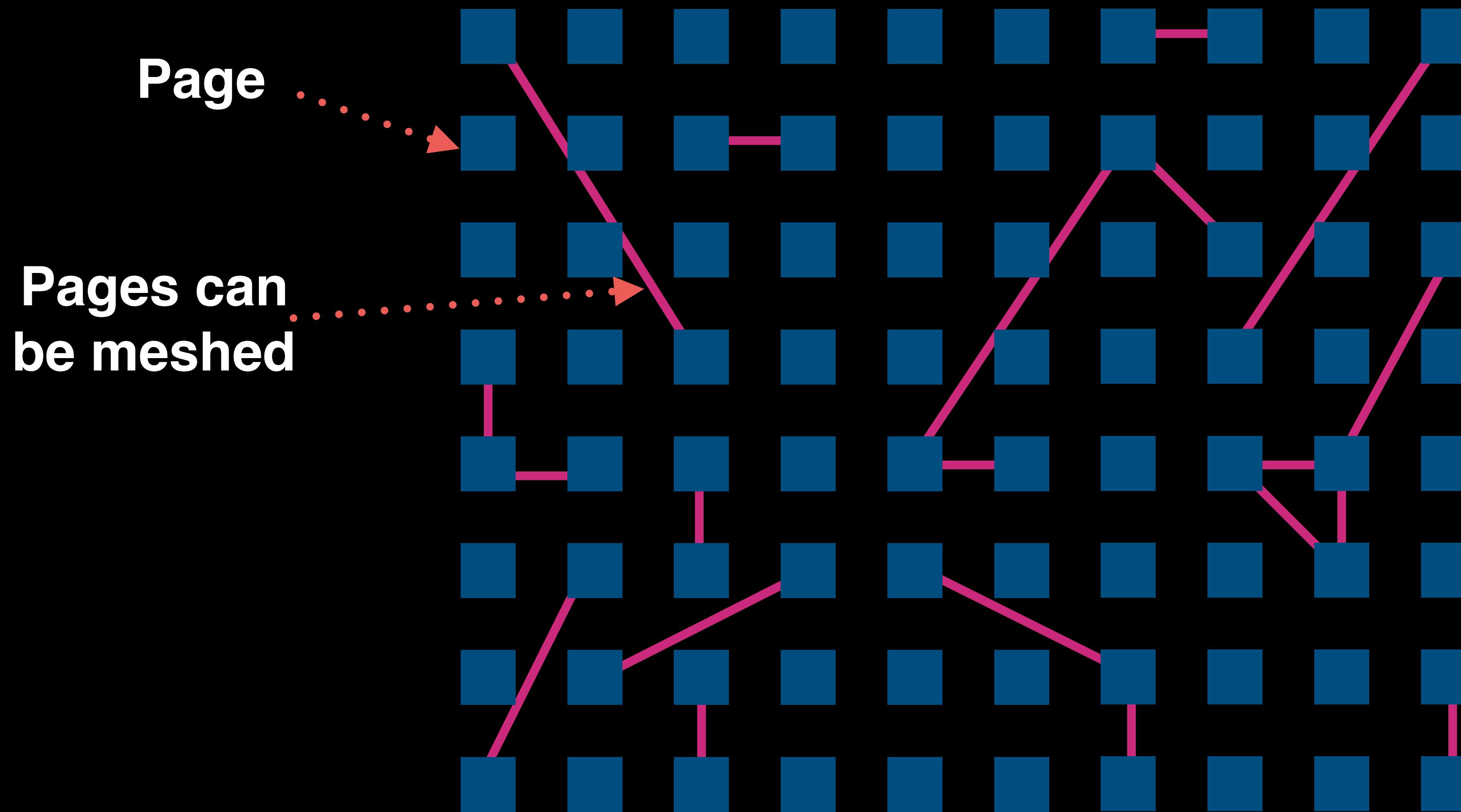
Finding pages to Mesh

Problem: Find meshing that releases maximum number of pages

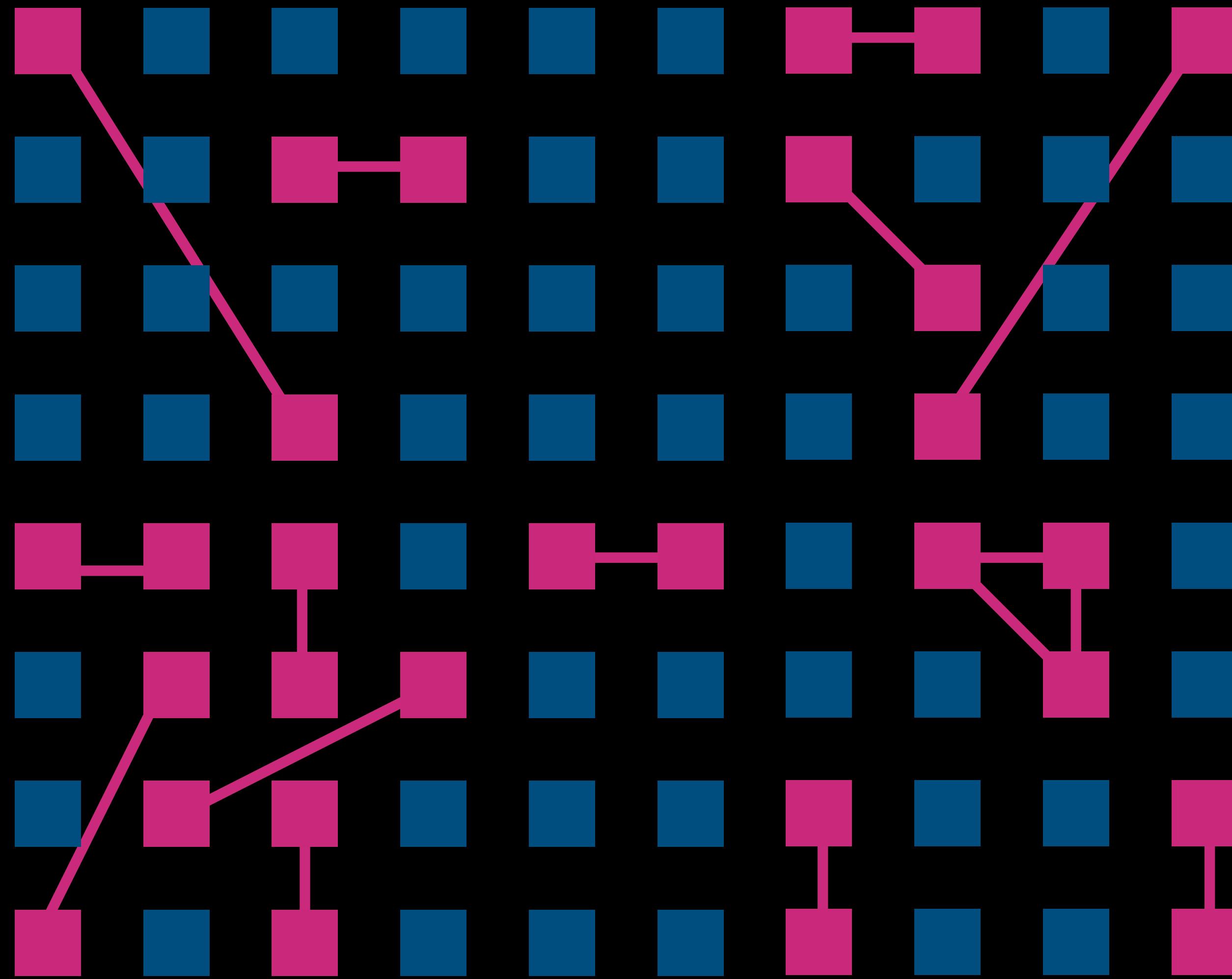
Problem: Find meshing that releases maximum number of pages

Run in `free()` slowpath
At most once every 100 ms
Treat each size class independently

Problem: Find meshing that releases the maximum number of pages



MinCliqueCover



MinCliqueCover

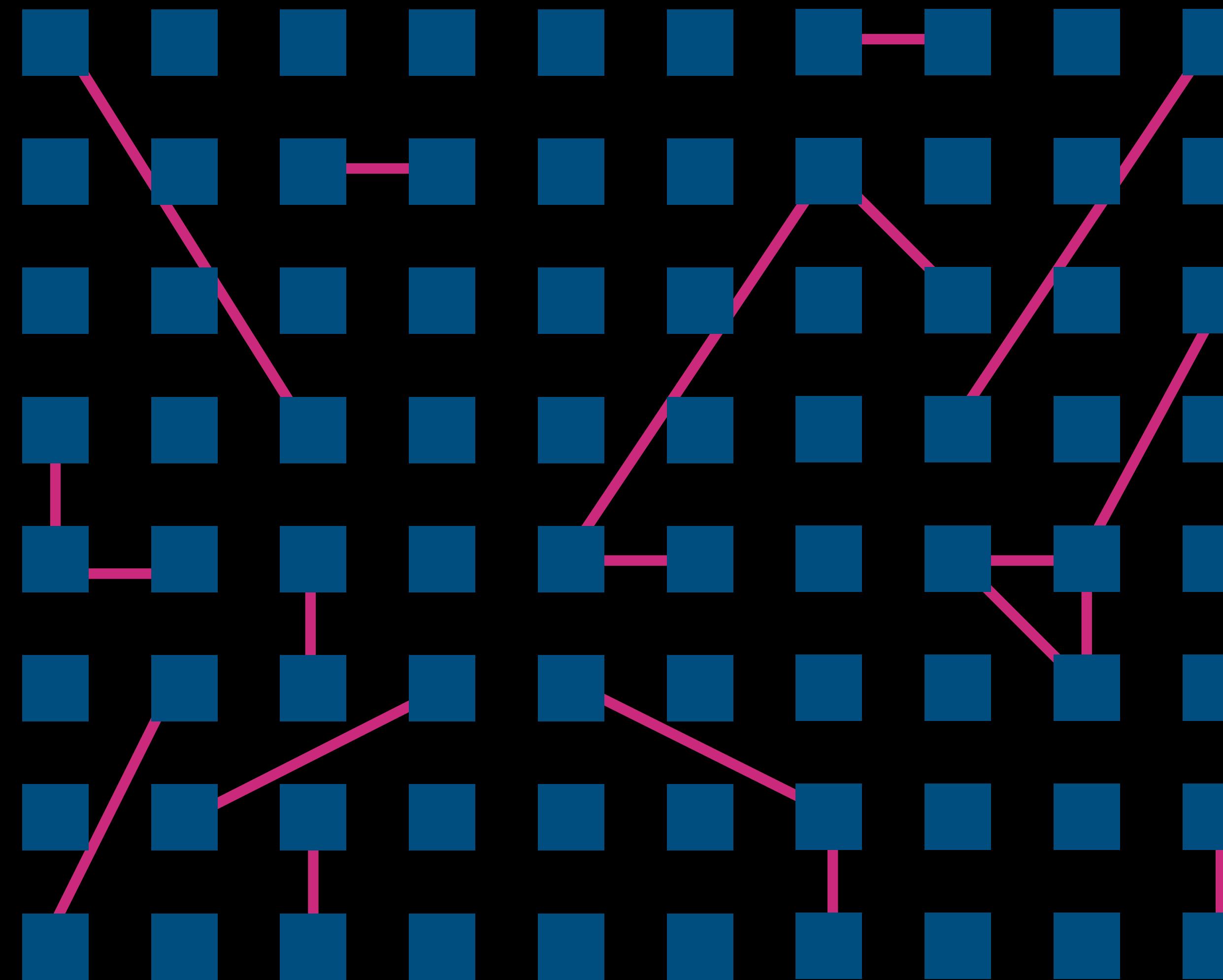
(NP-Complete)

MinCliqueCover

(NP-Complete)

BUT! Randomness ensures we can get
away with solving simpler graph problem
(Matching)

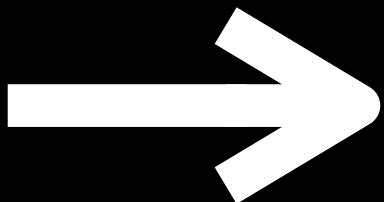
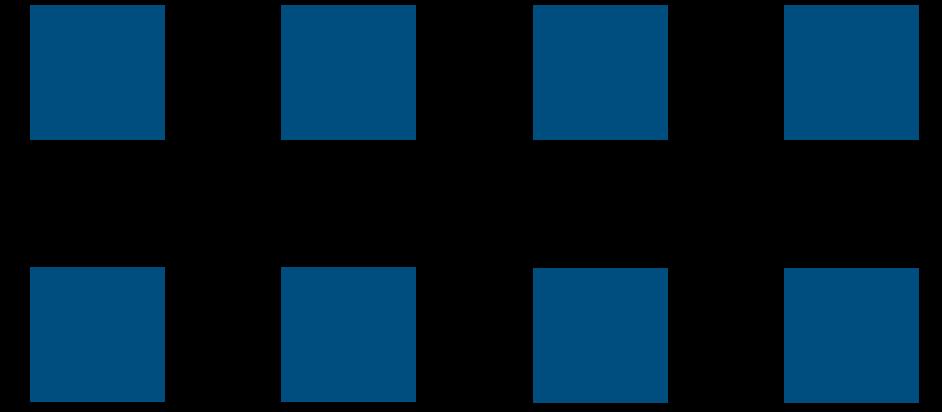
Wrinkle: building this graph would require RAM + time



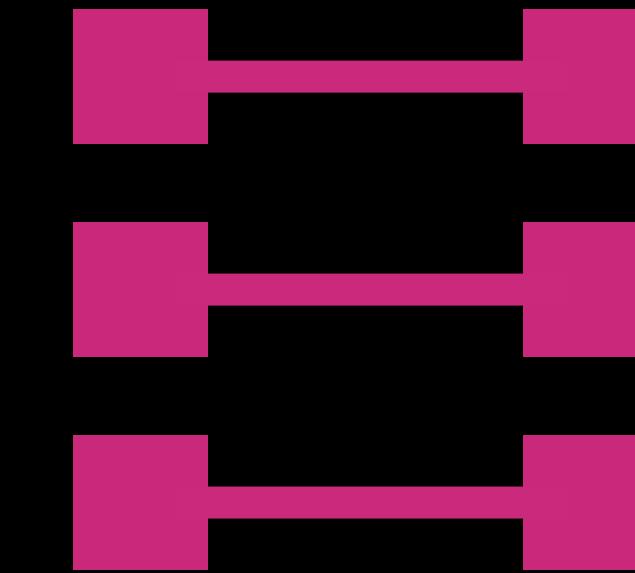
SplitMesher: approximates Matching without materializing meshing graph

SplitMesher: approximates Matching without materializing meshing graph

Set of partially full pages



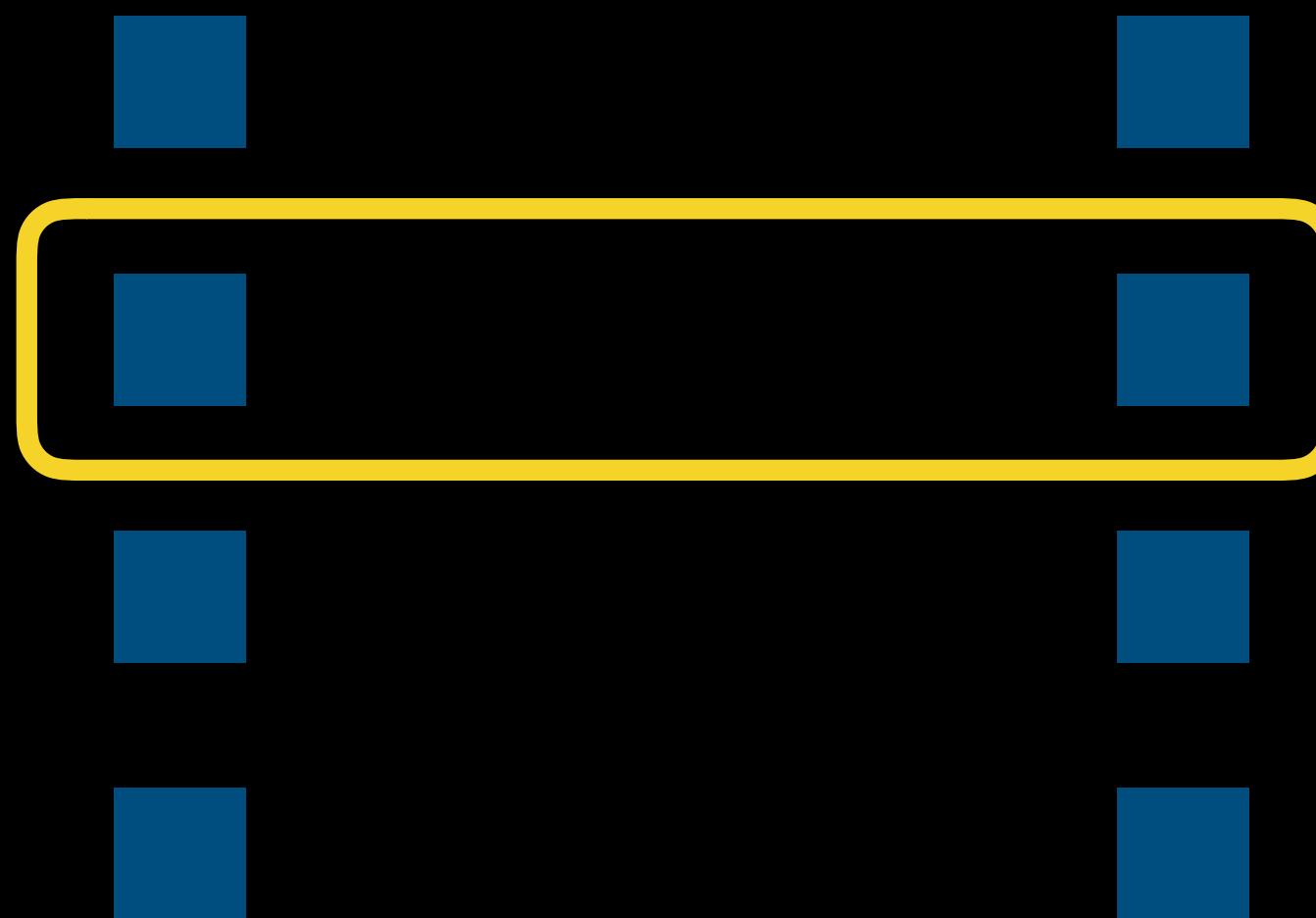
Pairs of meshable pages



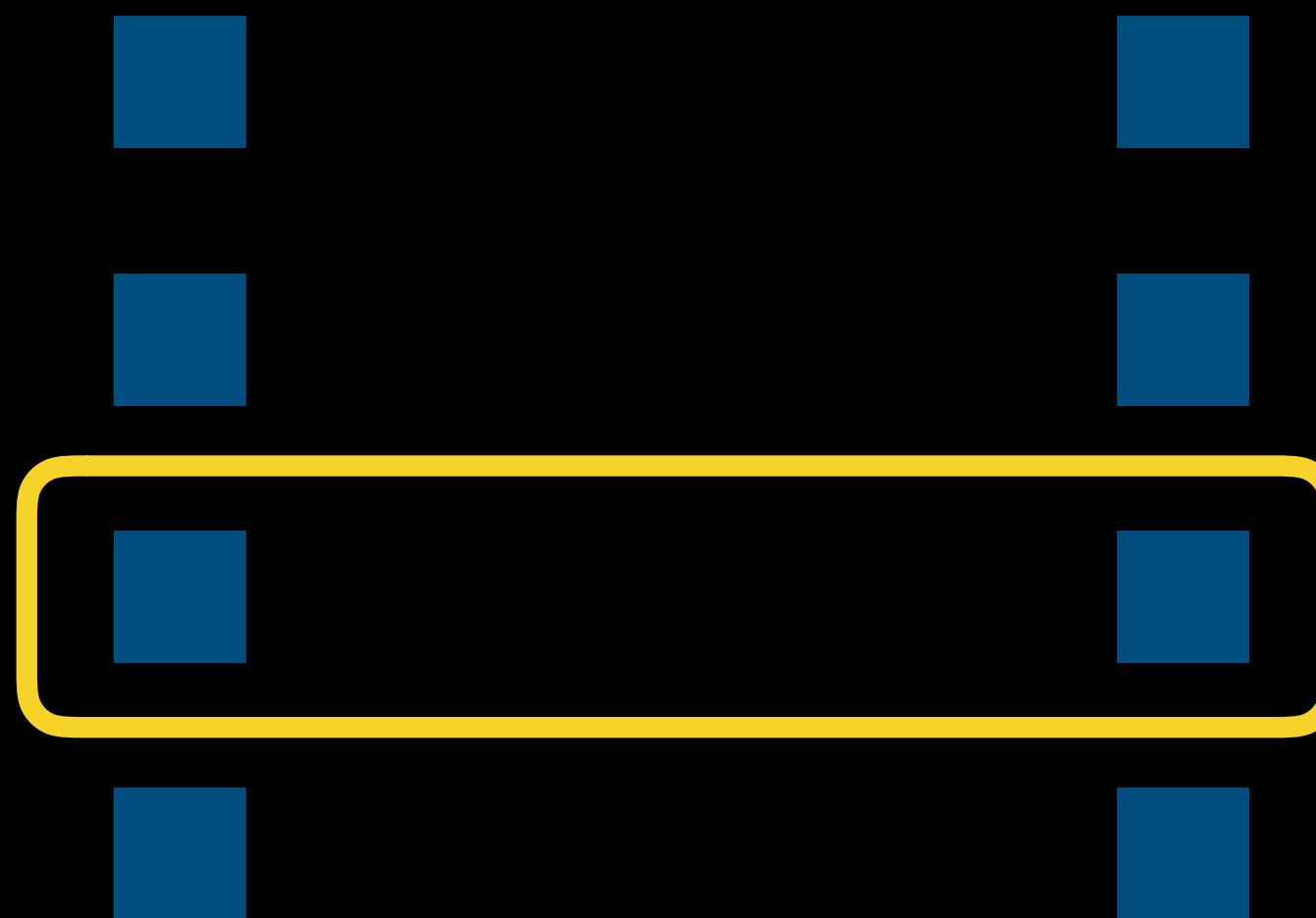
Iterate, comparing $a[i]$ to $b[i]$



Iterate, comparing $a[i]$ to $b[i]$



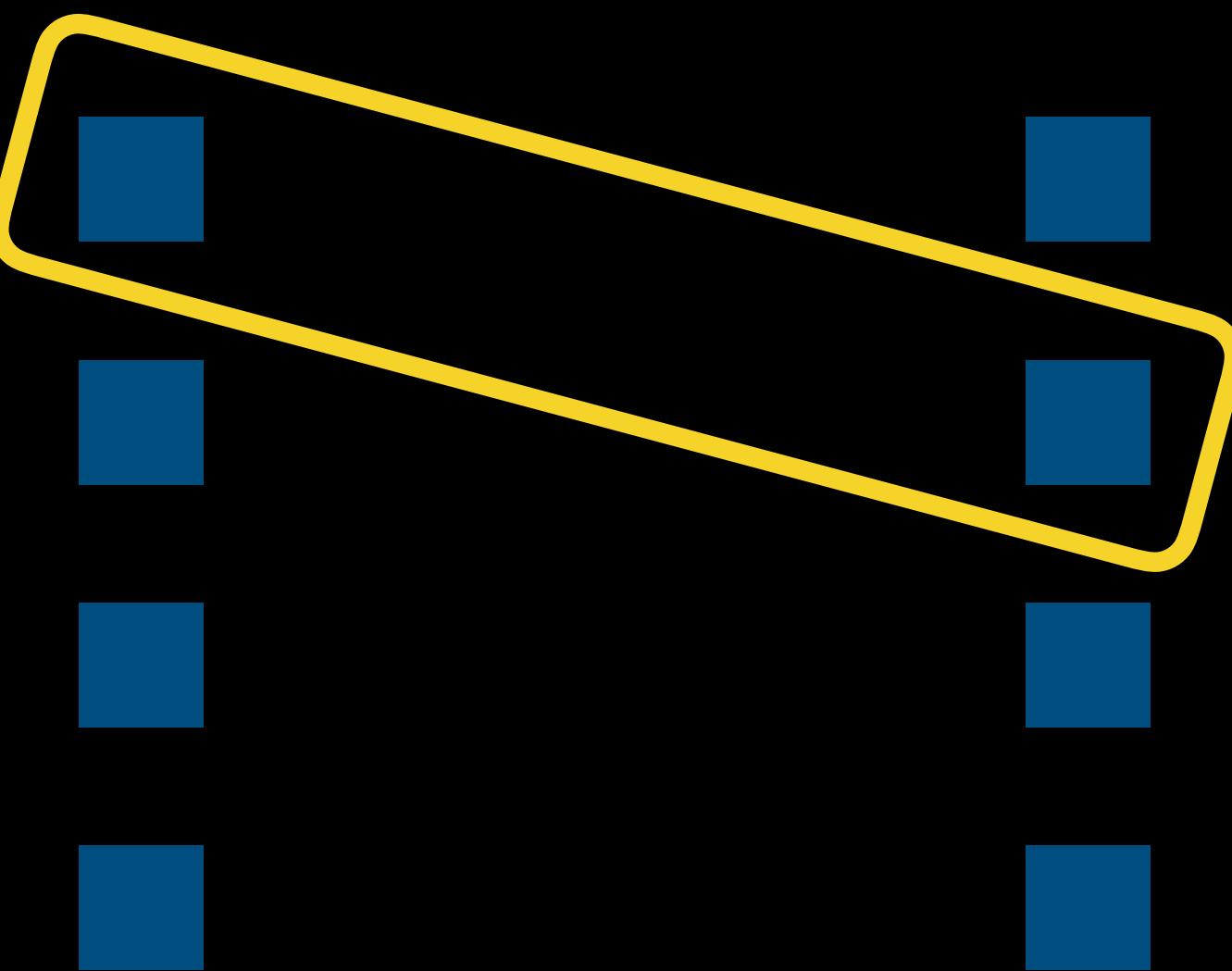
Iterate, comparing $a[i]$ to $b[i]$



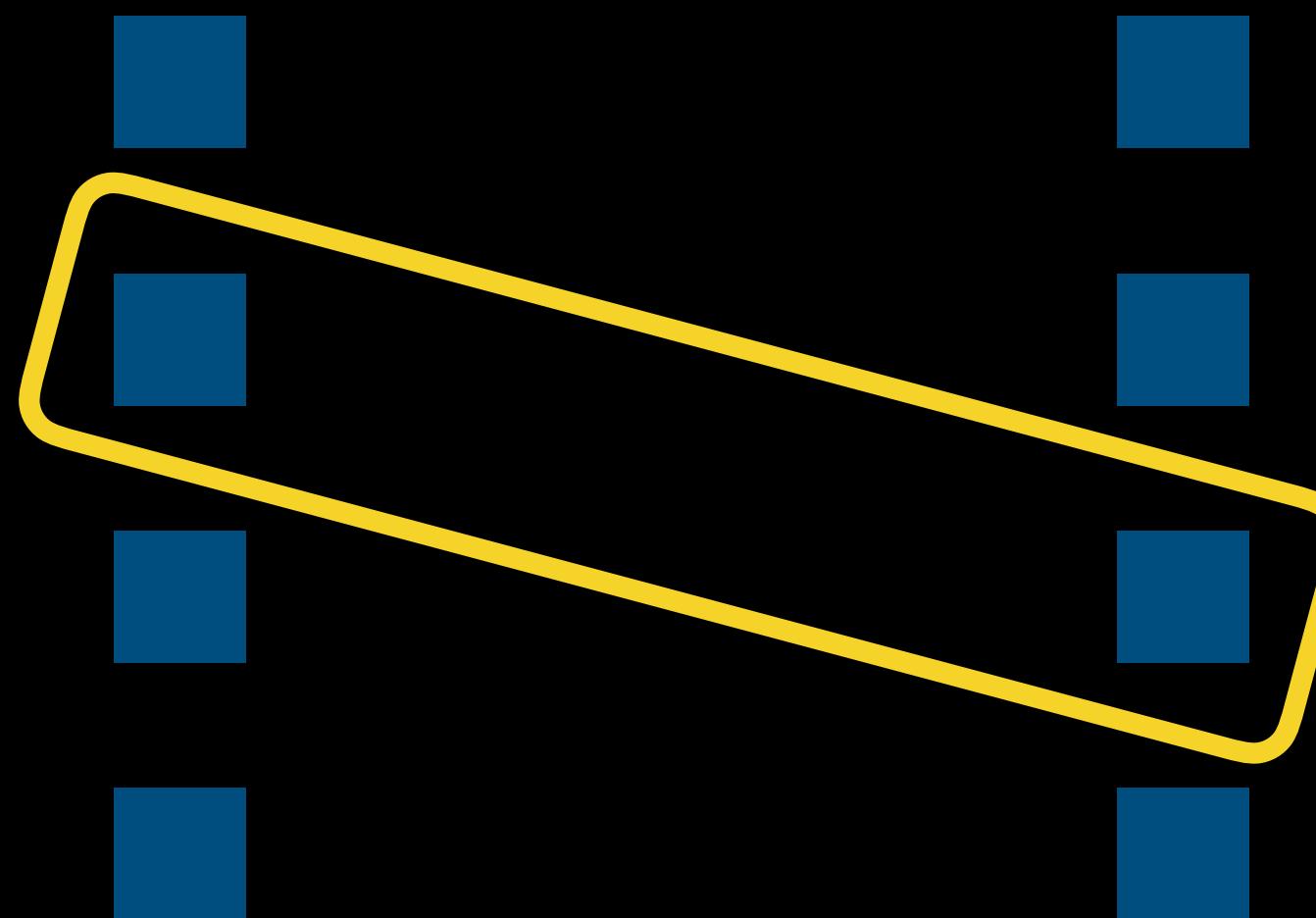
Iterate, comparing $a[i]$ to $b[i]$

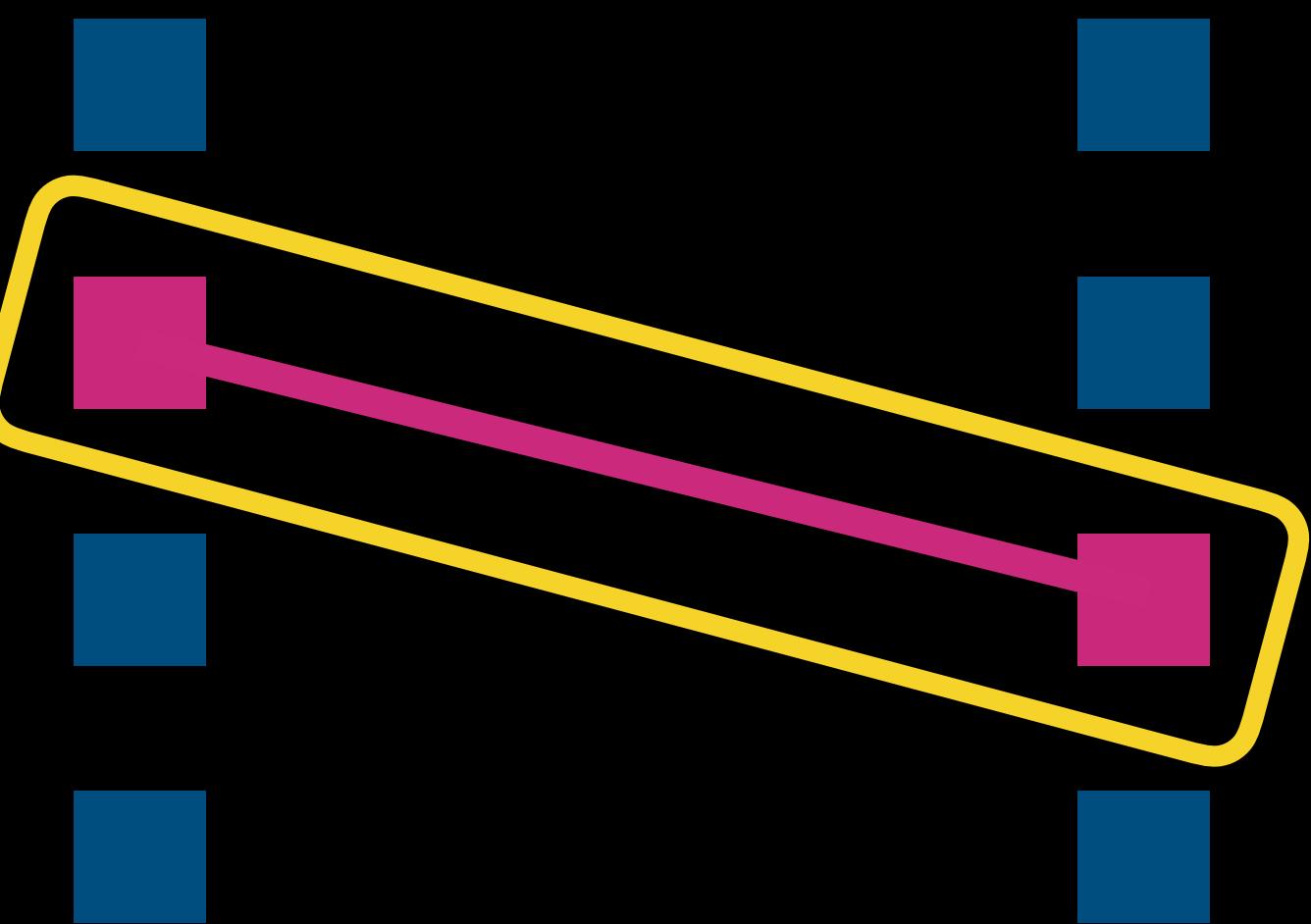


loop, comparing $a[i]$ to $b[(i+1)\%len]$

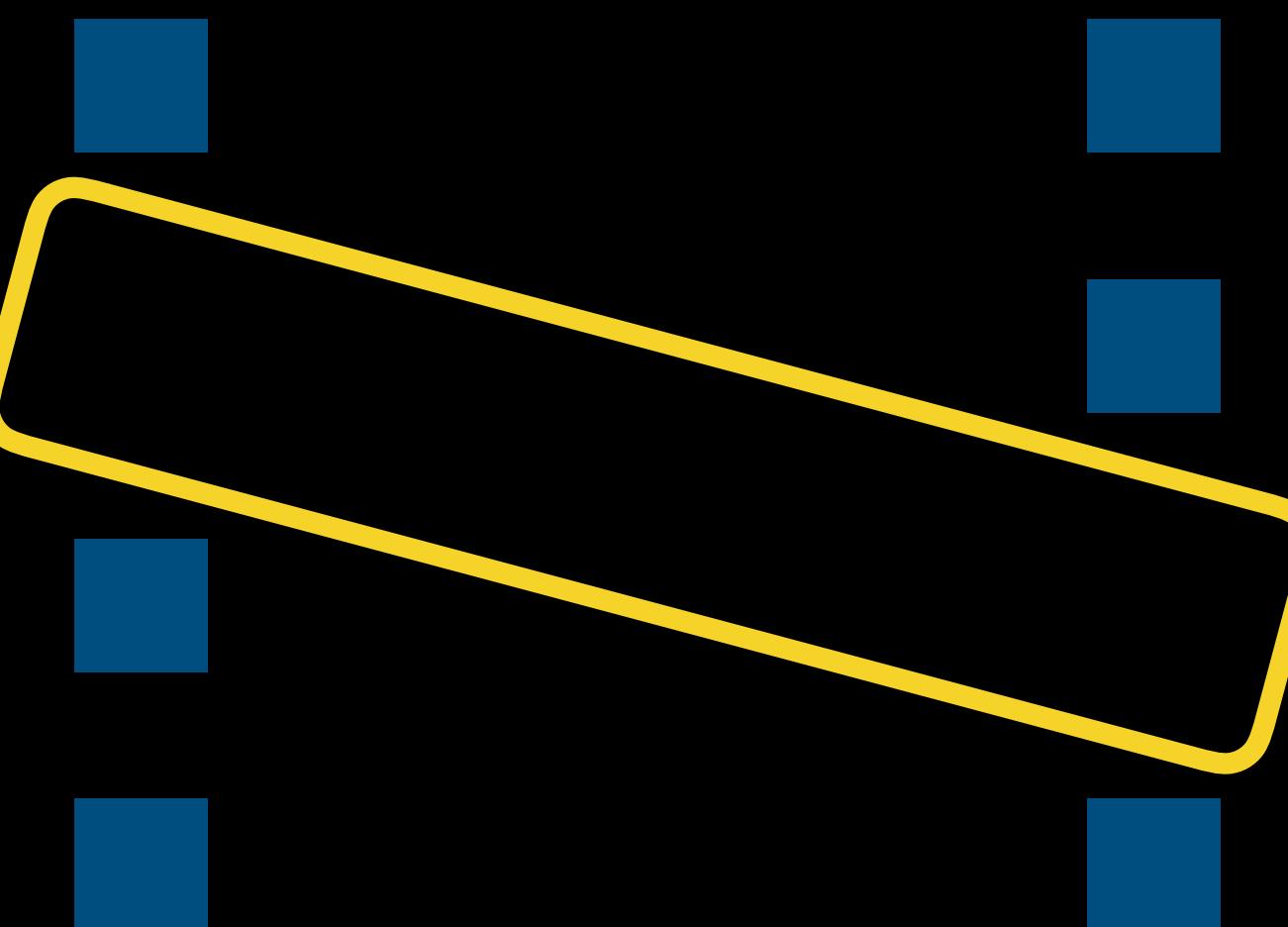


loop **fixed** number of times,
comparing $a[i]$ to $b[(i+1)\%len]$

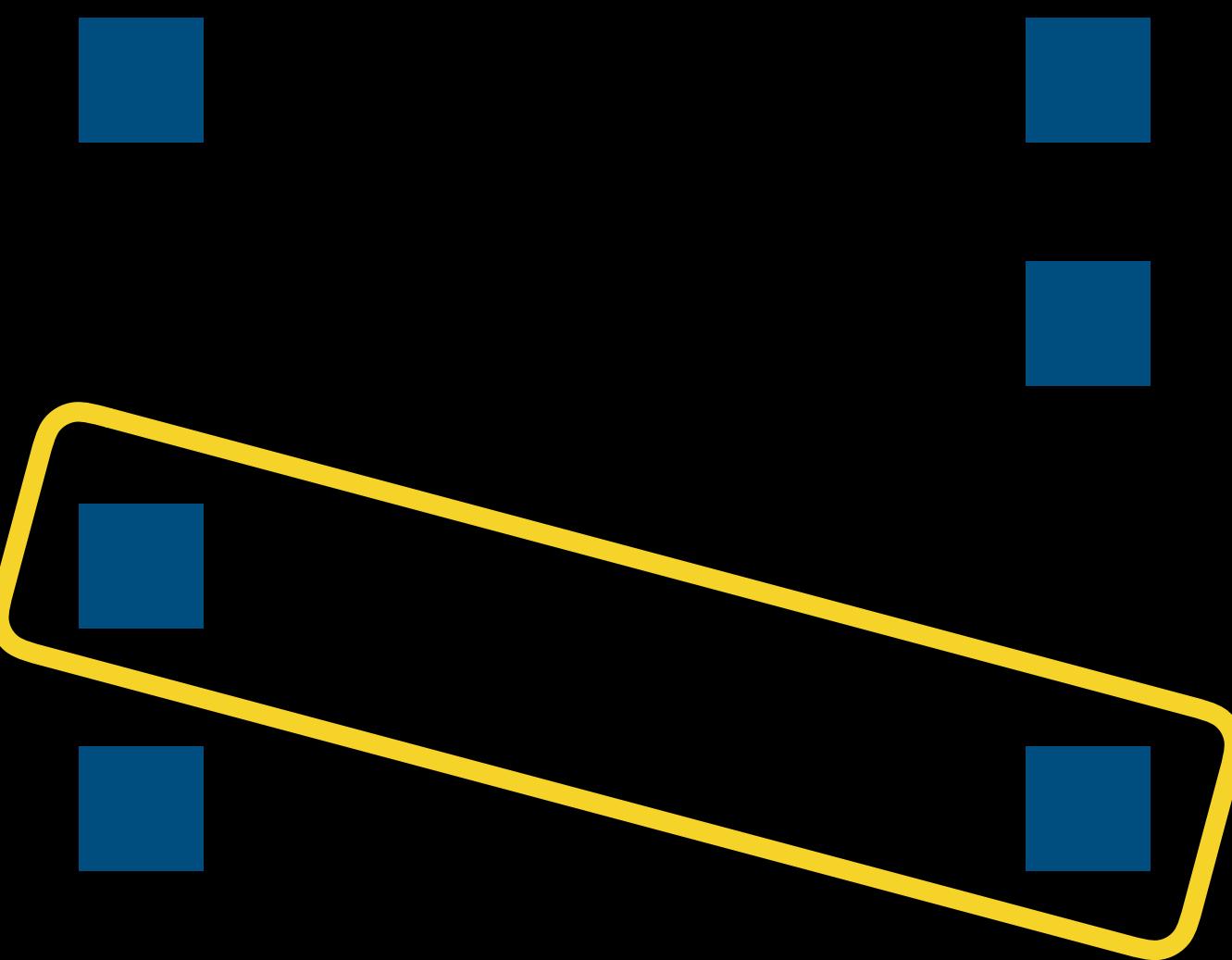




Remove found match



Continue



SplitMesher: approximates Matching without materializing meshing graph

SplitMesher: approximates Matching
without materializing meshing graph

$O(n/q)$ time

(q is the global probability of spans meshing)

SplitMesher: approximates Matching
without materializing meshing graph

$O(n/q)$ time

(q is the global probability of spans meshing)

$1/2^*$ approximation w.h.p.

MESH

implementation details

MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

fast framework for building custom allocators!

MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

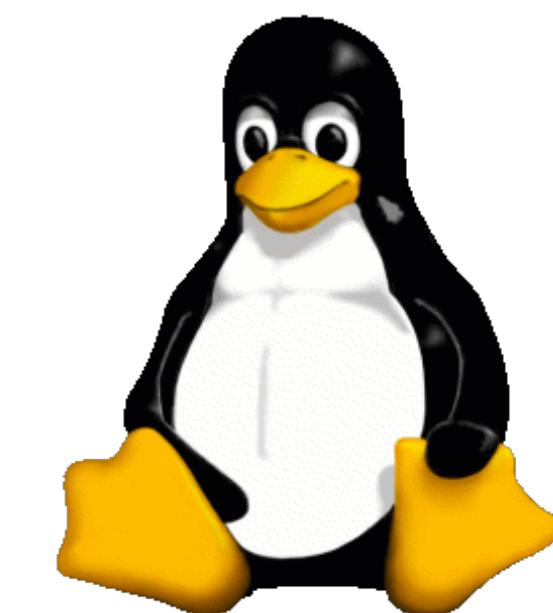
< 5K SLOC, C++

MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

< 5K SLOC, C++

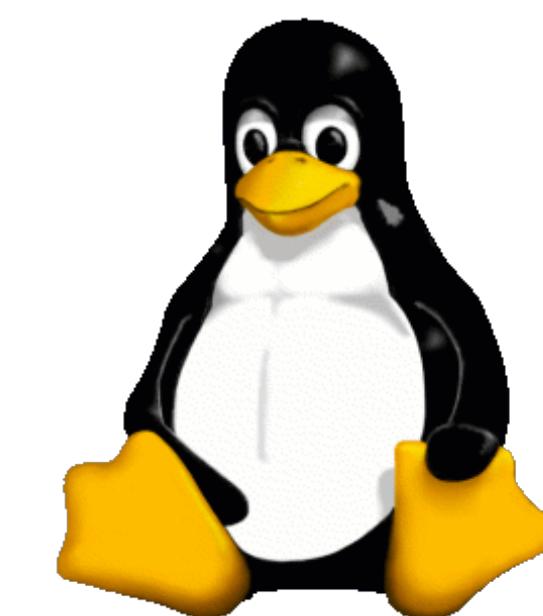


MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

< 5K SLOC, C++

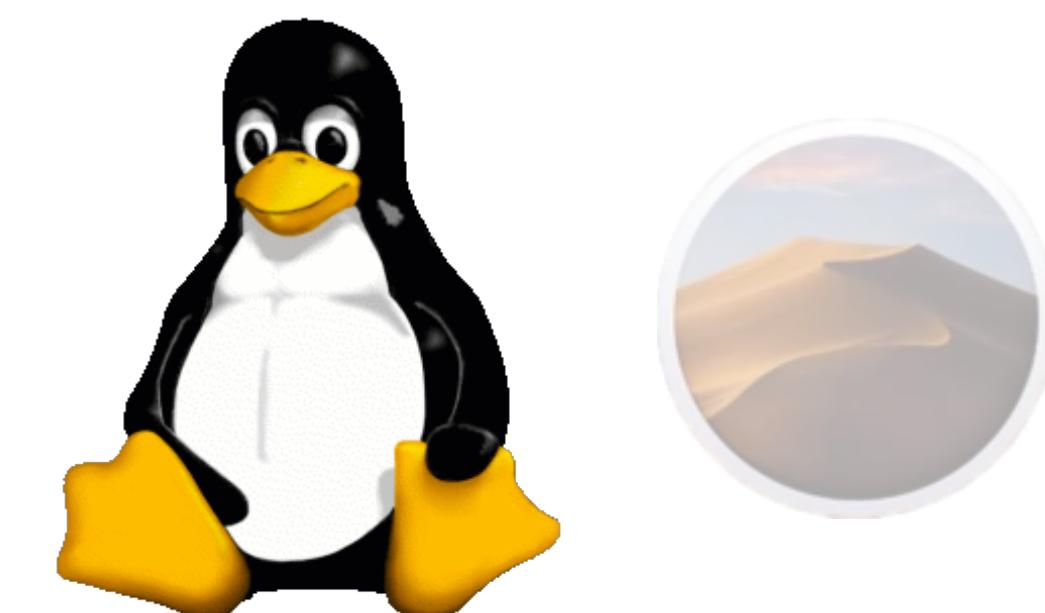


MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

< 5K SLOC, C++

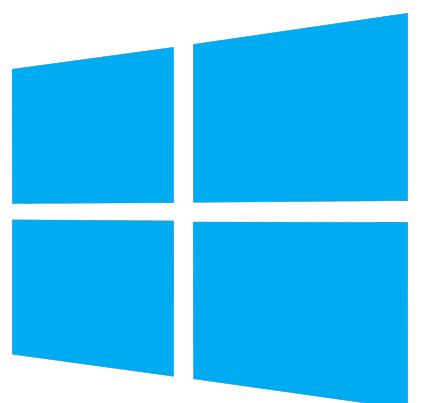
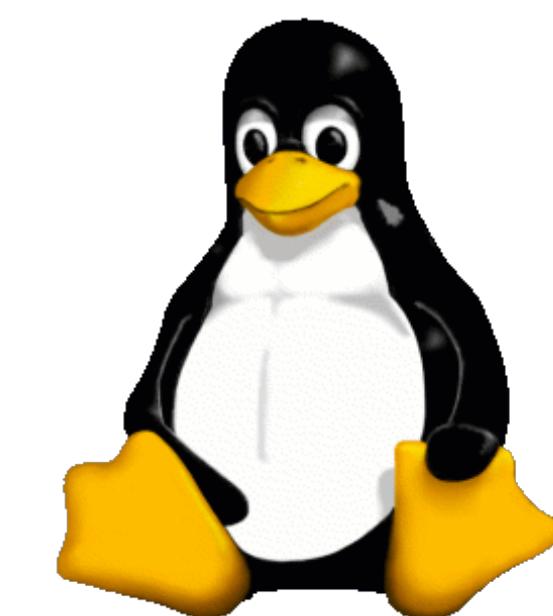


MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

< 5K SLOC, C++

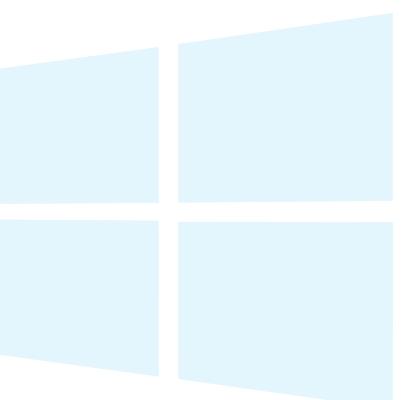
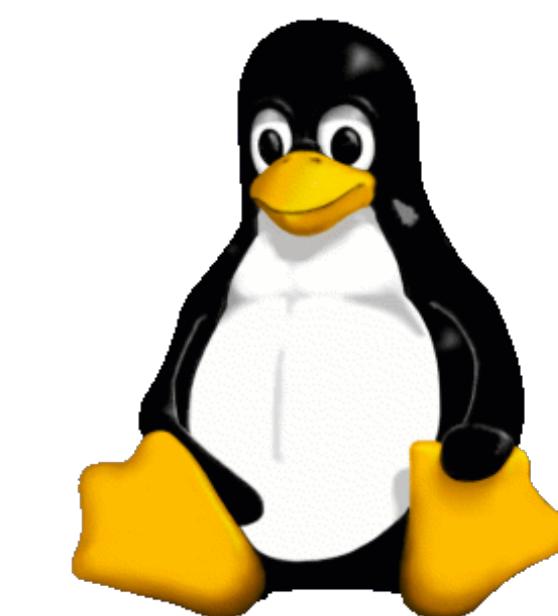


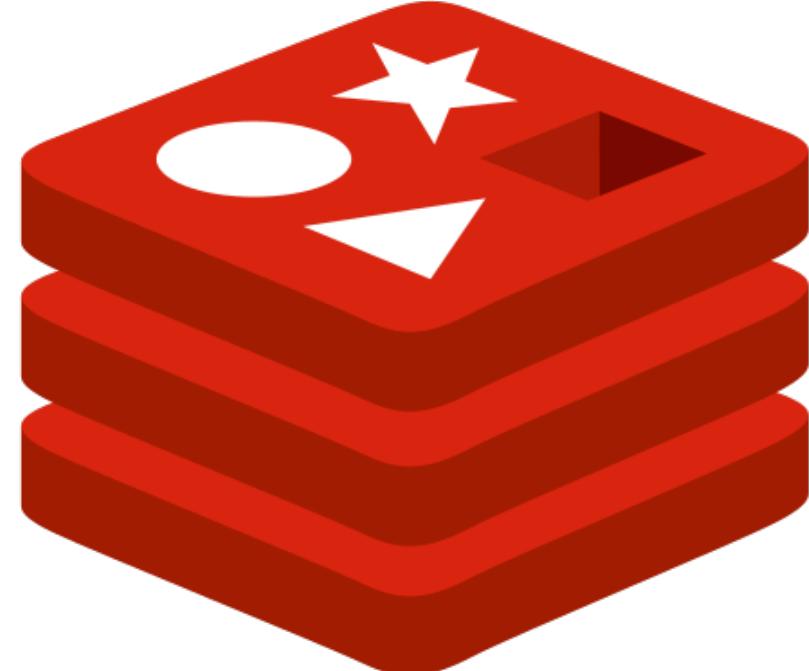
MESH implementation details built using Heap Layers

github.com/emeryberger/Heap-Layers

(used to build Hoard, DieHard, DieHarder...)

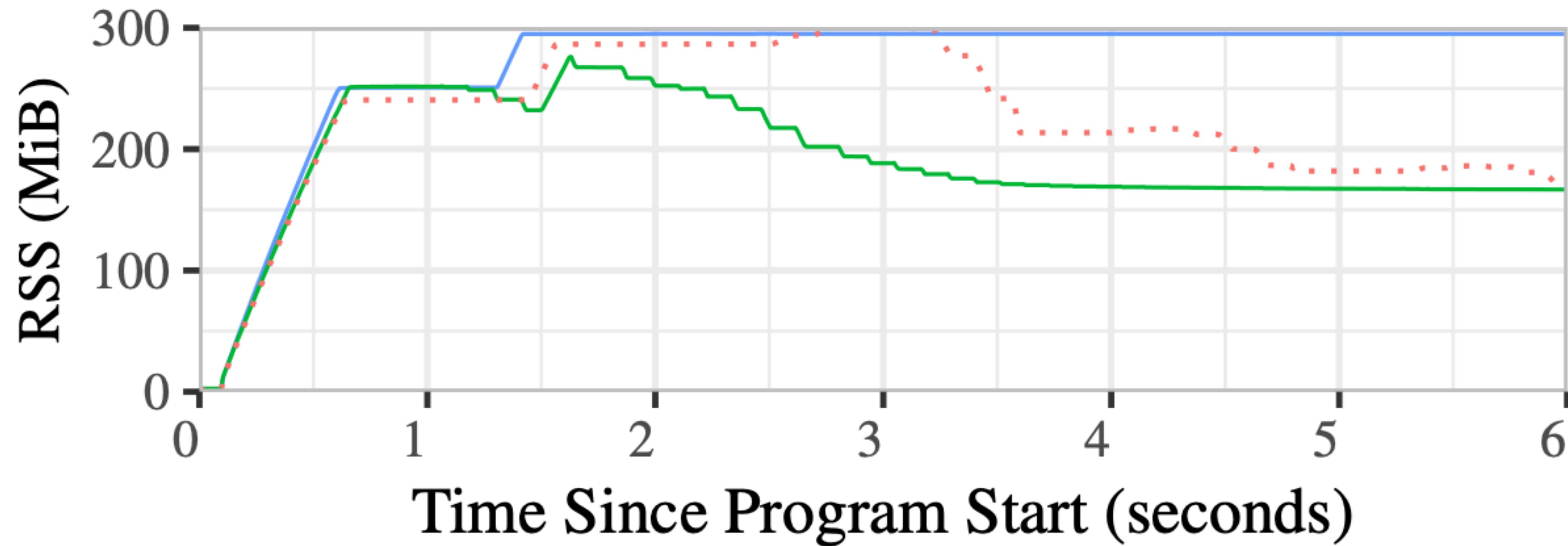
< 5K SLOC, C++

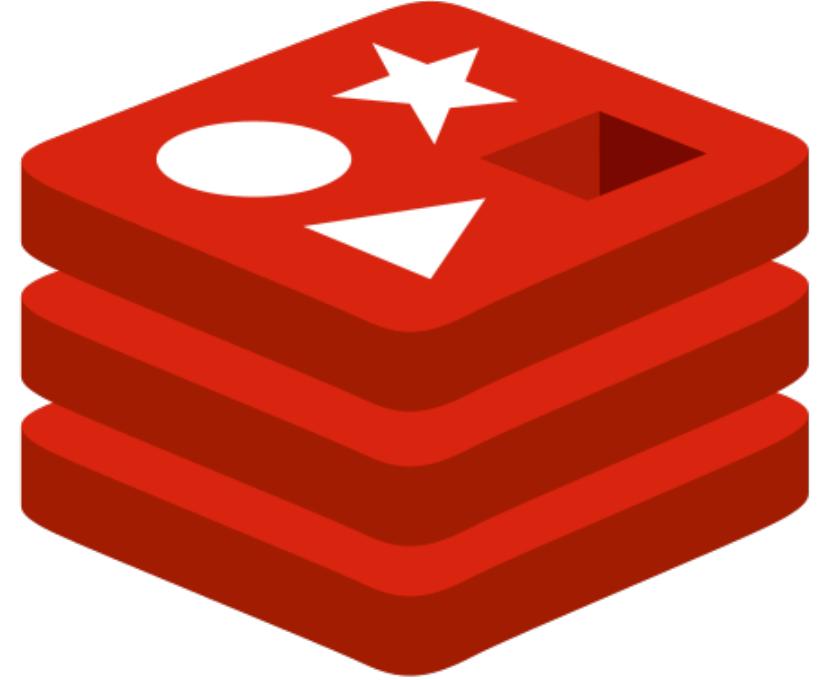




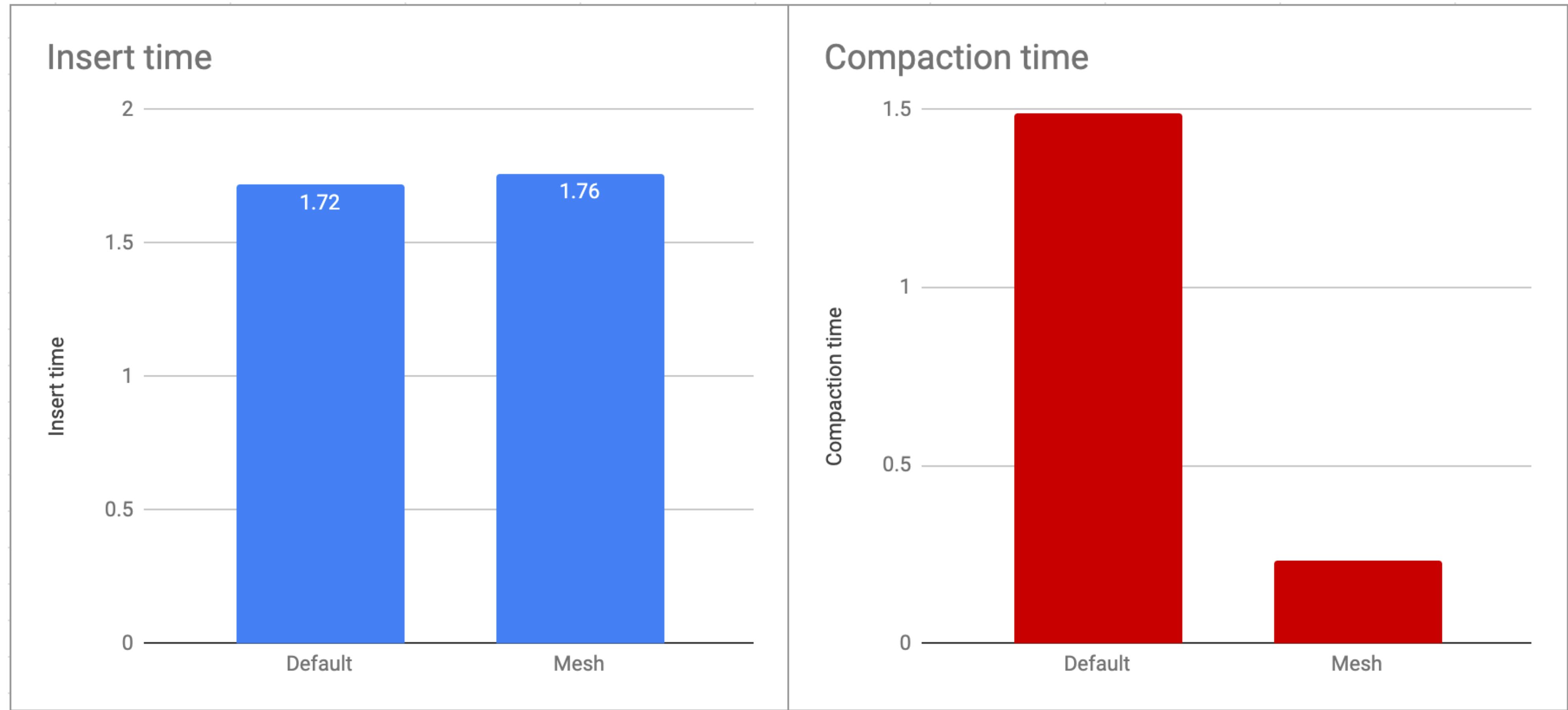
redis

- jemalloc + activatedefrag
- Mesh
- no compaction



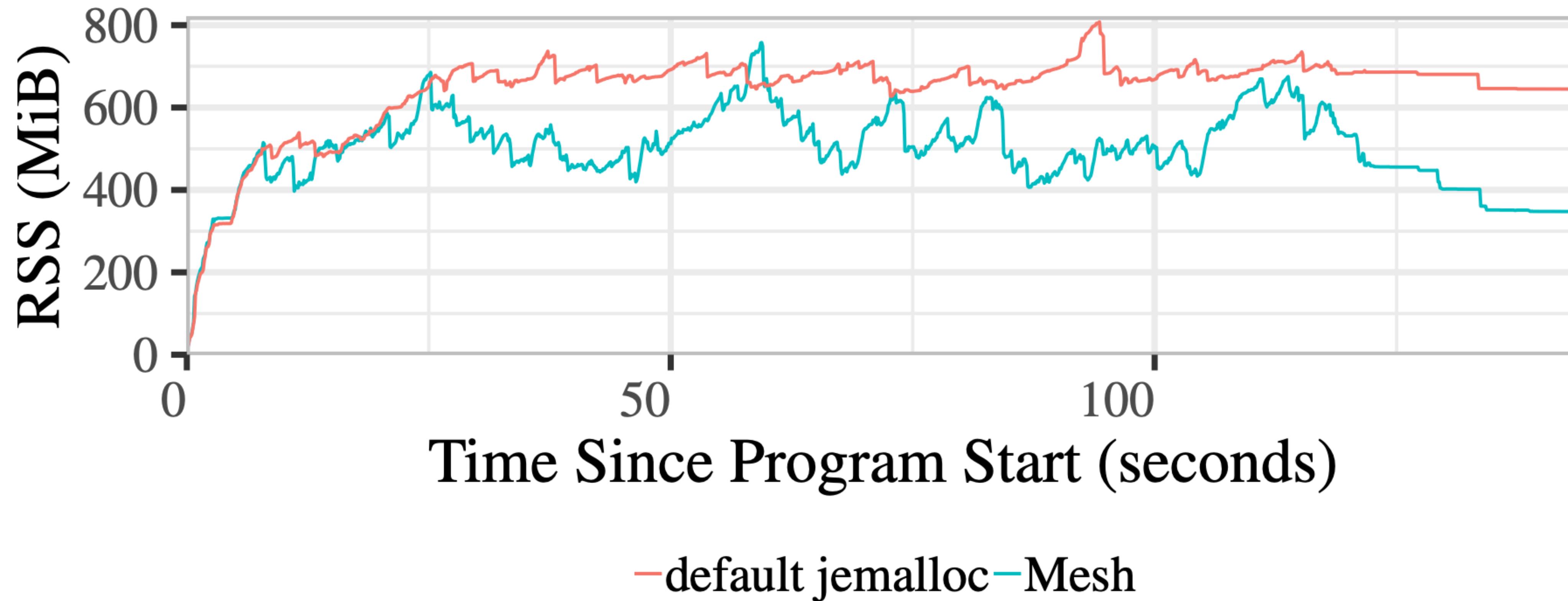


redis + MESH

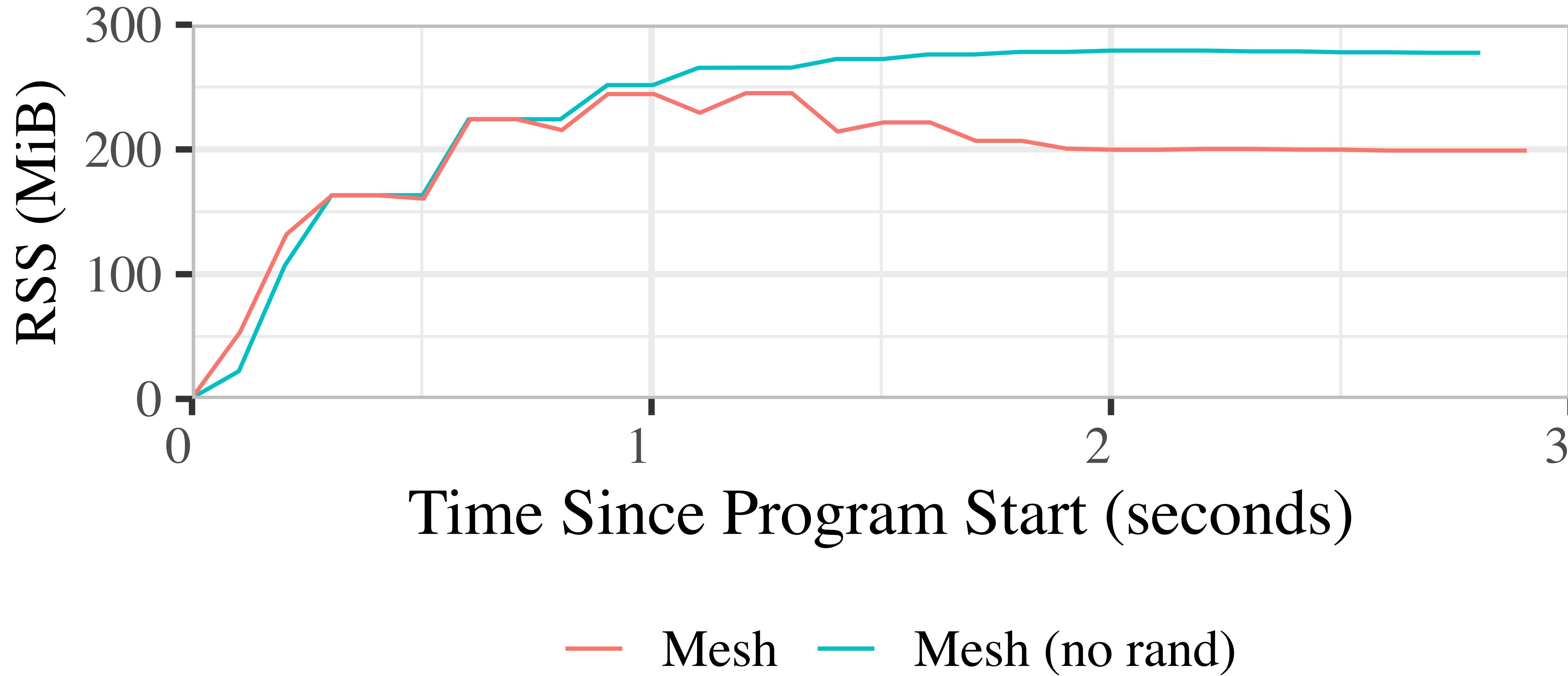


17% heap size reduction

< 1% performance overhead



Ruby Compaction for Free



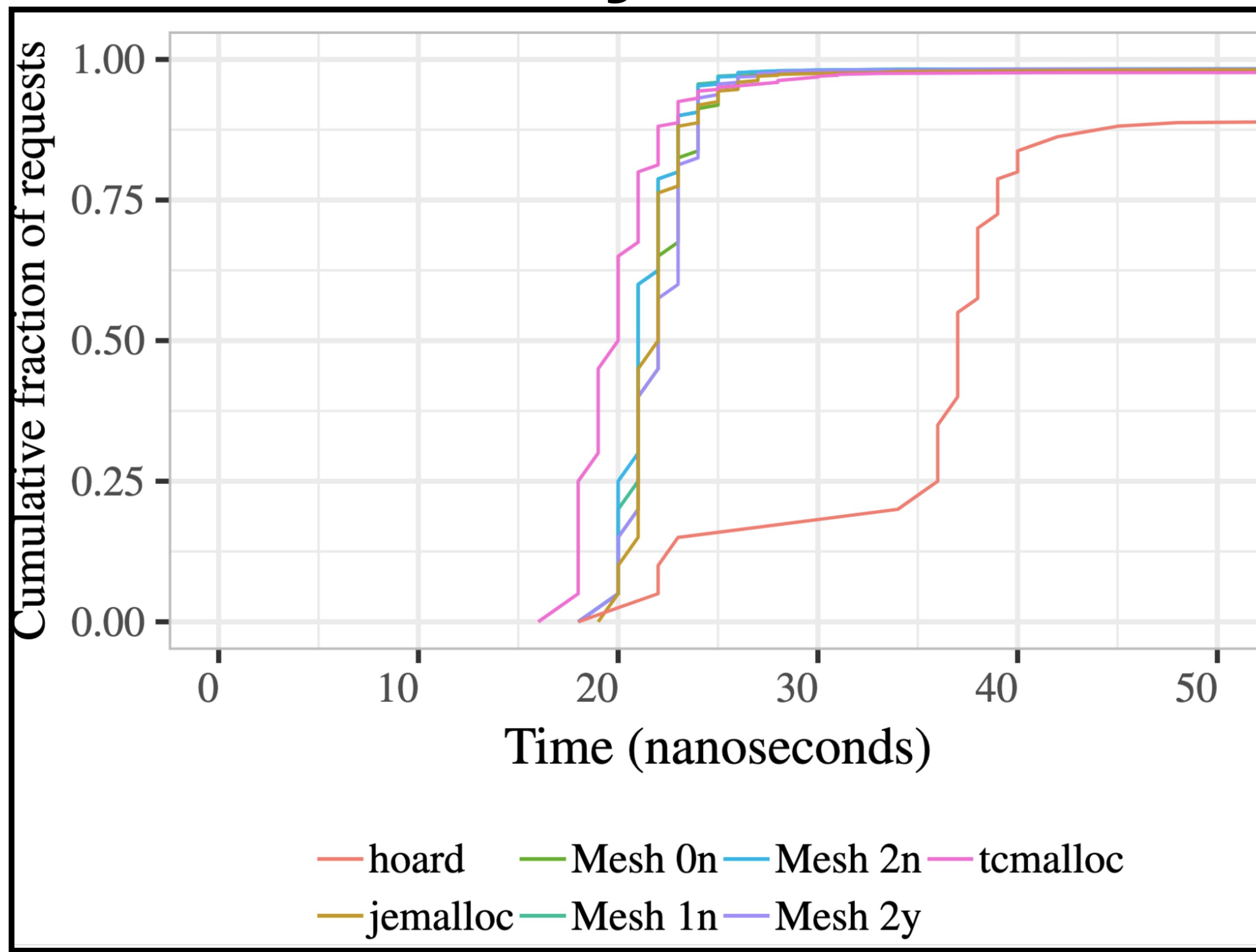
<http://LIBMESH.org>

Compaction without relocation!



Browsix, Coz, DieHard, Hoard...

Latency: malloc



Latency: free

